## Authentication

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## Authentication

The process of verifying someone's identity or role User, device, service, request, ...

What is identity?
Which characteristics uniquely identify an entity?
Authentication is a critical service
Enables communicating parties to verify the identity of their peers
Many other security mechanisms rely on it
Two main types
Human to computer
Computer to computer

## Credentials

## Evidence used to prove an identity

User Authentication: credentials supplied by a person
Something you know
Something you have
Something you are
Computer authentication: cryptography, secret tokens, location, ...
Computers (in contrast to humans) can "remember" large secrets (keys or tokens) and perform complex cryptographic operations
Location: evidence that an entity is at a specific place (IP, subnet, switch port, ...)
Authentication can be delegated
The verifying entity relies on a trusted third party to establish authentication $\rightarrow$ Identity and Access Management (IAM) services (e.g., Okta, Duo, OneLogin)

## Something You Know: Password-based Authentication

Passwords, passphrases, pins, key-phrases, access codes, ...
Good passwords are easy to remember and hard to guess
Easy to remember $\rightarrow$ easy to guess
Hard to guess $\rightarrow$ hard to remember
Bad ideas: date of birth, SSN, zip code, favorite team name, ...
Password space (bits) depends on:
Password length
Character set
Better way to think about strong passwords: long passphrases
Can be combined with custom variations, symbols, numbers, capitalization, ...


THROUGH 20 YEARS OF EFFORT, WE'VE SUCCESSFULLY TRAINED EVERYONE TO USE PASSWORDS THIAT ARE HARD FOR HUMANS TO REMEMBER, BUT EASY FOR COMPUTERS TO GUESS.
"online_throttling_100_per_hour": "3600000000036.000199840144435",
"online_no_throttling_10_per_second": "10000000000.1"
"online_no_throttling_10_per_second": "10000000000.1",
"offline_slow_hashing_1e4_per_second": "10000000.0001",
\}, "of
"crack_times_display": \{
"online_throttling_100_per_hour": "centuries",
"online_no_throttling_10_per_second": "centuries",
"offline_slow_hashing_1e4_per_second": "4 months" "offline_fast_hashing_1e10_per_second": "10 seconds"
\},
$\}^{\}}$

| Deductions | Type | Rate | Count | Bonus |
| :--- | :--- | :--- | :--- | :--- |
| Letters only | Flat | $-n$ | 0 | 0 |
| Numbers only | Flat | $-n$ | 0 | 0 |
| Repeat Characters (case insensitive) | Comp | - | 2 | -1 |
| Consecutive uppercase letters | Flat | $-\left(n^{\star} 2\right)$ | 0 | 0 |
| Consecutive lowercase letters | Flat | $-\left(n^{\star} 2\right)$ | 3 | -6 |
| Consecutive numbers | Flat | $-\left(n^{\star} 2\right)$ | 0 | 0 |
| Sequential letters (3+) | Flat | $-\left(n^{\star} 3\right)$ | 0 | 0 |
| Sequential numbers (3+) | Flat | $-\left(n^{\star} 3\right)$ | 0 | 0 |
| Sequential symbols (3+) | Flat | $-\left(n^{\star} 3\right)$ | 0 | 0 |

            "online_throttling_100_per_hour": "7697230882272000427282.147568",
            "online_throttling_100_per_hour
            "online_no_throttling_10_per_second": "21381196895200000000
            "offline_fast_hashing_1e10_per_second": "21381196895.2"
    \},
"crack_times_display": \{
"online_throttling_100_per_hour": "centuries",
"online_no_throttling_10_per_second": "centuries",
"offline_slow_hashing_1e4_per_second": "centuries",
"offline_fast hashing_1e10_per_second": "centuries"
\},

| Additions | Type | Rate | Count | Bonus |
| :--- | :--- | :--- | :--- | :--- |
| Number of characters | Flat | $+\left(n^{\star} 4\right)$ | 28 | +112 |
| Uppercase letters | Cond/lncr | $+\left((l e n-n)^{\star} 2\right)$ | 0 | 0 |
| Lowercase Letters | Cond/lncr | $+\left((l e n-n)^{\star} 2\right)$ | 25 | +6 |
| Numbers | Cond | $+\left(n^{\star} 4\right)$ | 0 | 0 |
| Symbols | Flat | $+\left(n^{\star} 6\right)$ | 0 | 0 |
| Middle numbers or symbols | Flat | $+\left(n^{\star} 2\right)$ | 0 | 0 |
| Requirements | Flat | $+\left(n^{\star} 2\right)$ | 2 | 0 |


| Deductions | Type | Rate | Count | Bonus |
| :--- | :--- | :--- | :--- | :--- |
| Letters only | Flat | $-n$ | 28 | -28 |
| Numbers only | Flat | $-n$ | 0 | 0 |
| Repeat Characters (case insensitive) | Comp | - | 20 | -2 |
| Consecutive uppercase letters | Flat | $-\left(n^{\star} 2\right)$ | 0 | 0 |
| Consecutive lowercase letters | Flat | $-\left(n^{*} 2\right)$ | 24 | -48 |
| Consecutive numbers | Flat | $-\left(n^{\star} 2\right)$ | 0 | 0 |
| Sequential letters $(3+)$ | Flat | $-\left(n^{*} 3\right)$ | 0 | 0 |
| Sequential numbers $(3+)$ | Flat | $-\left(n^{*} 3\right)$ | 0 | 0 |
| Sequential symbols $(3+)$ | Flat | $-\left(n^{*} 3\right)$ | 0 | 0 |

## Password Policies (often have the opposite effect)

Password rules (miss the point)
"At least one special character," "Minimum/Maximum length of 8/12 characters," "Must contain at least one number," "Must contain at least one capital letter"
Makes passwords hard to remember! $\rightarrow$ encourages password reuse
Better: encourage long passphrases, and evaluate strength on-the-fly
Periodic password changing (does more harm than good)
"You haven't changed your password in the last 90 days"
Probably too late anyway if password has already been stolen
Makes remembering passwords harder $\rightarrow$ more password resets
Hinders the use of password managers (!)
What users do: password $1 \rightarrow$ password $2 \rightarrow$ password $3 \rightarrow \ldots$

If the chosen secret is found in the list, the CSP or verifier SHALL advise the subs


#### Abstract

provide the reason for rejection, and SHALL require the subscriber to choose a di


 Verifiers SHOULD offer guidance to the subscriber, such as a password-strength memorized secret. This is particularly important following the rejection of a memo modification of listed (and likely very weak) memorized secrets [Blacklists].
## Digital Identity Guidelines

Authentication and Lifecycle Management

Verifiers SHALL implement a rate-limiting mechanism that effectively limits the number of failed authentication attempts that can be made on the subscriber's account as described in Section 5.2.2.

Verifiers SHOULD NOT impose other composition rules (e.g., requiring mixtures of different character types or prohibiting consecutively repeated characters) for memorized secrets. Verifiers SHOULD NOT require memorized secrets to be changed arbitrarily (e.g., periodically). However, verifiers SHALL force a change if there is evidence of compromise of the authenticator.

Verifiers SHOULD permit claimants to use "paste" functionality when entering a memorized secret. This facilitates the use of password managers, which are widely used and in many cases increase the likelihood that users will choose stronger memorized secrets.

In order to assist the claimant in successfully entering a memorized secret, the verifier SHOULD offer an option to display the secret - rather than a series of dots or asterisks - until it is entered. This allows the claimant to verify their entry if they are in a location where their screen is unlikely to be observed. The verifier MAY also permit the user's device to display individual entered characters for a short time after each character is typed to verify correct entry. This is particularly applicable on mobile devices.

The verifier SHALL use approved encryption and an authenticated protected channel when requesting memorized secrets in order to provide resistance to eavesdropping and MitM attacks.

Verifiers SHALL store memorized secrets in a form that is resistant to offline attacks. Memorized secrets SHALL be salted and hashed using a

## Attacking Passwords

Offline cracking
Online guessing
Brute force attacks

Eavesdropping
Capturing

## Password Storage

## Storing passwords as plaintext is disastrous

Better way: store a cryptographic hash of the password

## Even better: store the hash of a "salted" version of the password

Defend against dictionary attacks: prevent precomputation of hash values (wordlists of popular passwords, rainbow tables, ...)
Unique salt per user (no need to be secret): even if two users happen to have the same password, their hash values will be different $\rightarrow$ need to be cracked separately
Salting does not make brute-force guessing a given password harder!

| Username | Salt | Password hash |
| :--- | :--- | :--- |
| Bobbie | 4238 | h(4238, \$uperman) |
| Tony | 2918 | h(2918, 63\%TaeFF) |
| Mitsos | 6902 | h(6902, zour1da) |
| Mark | 1694 | h(1694, Rockybrook\#1) |

Password databases are still getting leaked...

## Password Cracking

Exhaustive search $\rightarrow$ infeasible for large password spaces
Dictionary attacks (words, real user passwords from previous leaks, ...)
Variations, common patterns, structure rules
Prepend/append symbols/numbers/dates, weird capitalization, I33tspeak, visually similar characters, intended misspellings, ...

Target-specific information
DOB, family names, favorite team, pets, hobbies, anniversaries, language, slang, ...
Easy to acquire from social networking services and other public sites
Particularly effective against "security questions"
Advanced techniques
Probabilistic context-free grammars, Markov models, ...
hashcat
advanced
password
recovery
hashcat Forums Wiki Tools Events

## Example hashes

If you get a "line length exception" error in hashcat, it is often because the hash mode that you have requested does not match the hash. To verify, you can test your commands against example hashes.

Unless otherwise noted, the password for all example hashes is hashcat
Generic hash types

| Hash- <br> Mode | Hash-Name | Example |
| :--- | :--- | :--- |
| 0 | MD5 | 8743b52063cd84097a65d1633f5c74f5 |
| 10 | md5(\$pass.\$salt) | 01dfae6e5d4d90d9892622325959afbe:7050461 |
| 20 | md5(\$salt.\$pass) | f0fda58630310a6dd91a7d8f0a4ceda2:4225637426 |
| 30 | md5(utf16le(\$pass).\$salt) | b31d032cfdcf47a399990a71e43c5d2a:144816 |
| 40 | md5(\$salt.utf16le(\$pass)) | d63d0e21fdc05f618d55ef306c54af82:13288442151473 |
| 50 | HMAC-MD5 (key $=$ \$pass) | fc741db0a2968c39d9c2a5cc75b05370:1234 |
| 60 | HMAC-MD5 (key $=$ \$salt) | bfd280436f45fa38eaacac3b00518f29:1234 |
| 100 | SHA1 | b89eaac7e61417341b710b727768294d0e6a277b |
| 110 | sha1(\$pass.\$salt) | 2fc5a684737ce1bf7b3b239df432416e0dd07357:2014 |
| 120 | sha1(\$salt.\$pass) | cac35ec206d868b7d7cb0b55f31d9425b075082b:5363620024 |
| 130 | sha1(utf16le(\$pass).\$salt) | c57f6ac1b71f45a07dbd91a59fa47c23abcd87c2:631225 |
| 140 | sha1(\$salt.utf16le(\$pass)) | 5db61e4cd8776c7969cfd62456da639a4c87683a:8763434884872 |
| 150 | HMAC-SHA1 (key $=\$$ \$pass) | c898896f3f70f61bc3fb19bef222aa8600e5ea717:1234 |
| 160 | HMAC-SHA1 (key $=\$$ salt) | d89c92b4400b15c39e462a8caa939ab40c3aeeea:1234 |
| 200 | MySQL323 | 7196759210defdc0 |
| 300 | MySQL4.1/MySQL5 | fcf7c1b8749cf99d88e5f34271d636178fb5d130 |

## 50 Most-used (Worse) Passwords

| 123456 | 1234567 | 123 | ashley | evite |
| :--- | :--- | :--- | :--- | :--- |
| 123456789 | qwerty | omgpop | 987654321 | $123 a b c$ |
| picture1 | abc123 | 123321 | unknown | 123qwe |
| password | Million2 | 654321 | zxcvbnm | sunshine |
| 12345678 | 000000 | qwertyuiop | 112233 | 121212 |
| 111111 | 1234 | qwer123456 | chatbooks | dragon |
| 123123 | iloveyou | $123456 a$ | 20100728 | $1 q 2 w 3 e 4 r$ |
| 12345 | aaron431 | a123456 | 123123123 | 5201314 |
| 1234567890 | password1 | 666666 | princess | 159753 |
| senha | qqww1122 | asdfghjkl | jacket025 | 0123456789 |

Distribution of 4-digit
sequences within RockYou passwords


## Wordlists

| ce\#ebc.dk | 4637324 | gea8mw4yz | fujinshan | masich | gothpunksk8er | 20081010 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| goddess5 | bugger825 | kukumbike | counter | pengaiwei | rftaeo48 | leelou44 |
| 20071002 | marmaris | 260888 | N8mr0n | coalesce | 8d7R0K | 8UfjeGb0 |
| 271075711 | jinjin111 | jordi10 | 520057 | 56402768 | 5172032 | 200358808 |
| zs3cu7za | 170383gp | lexusis | adc123 | thesis | aics07 | dellede |
| scoopn | 3484427 | kj011a039 | bmaster | aabbcc894 | 34mariah | liang123. |
| frygas1411 | fl33321 | c84bwlrb | qbjh04zg | marion\&maxime | dongqinwei | captainettekt |
| SL123456sl | zwqrfg | priyanka05 | ueldaa79 | 614850 | samarica | kwiki-mart |
| 12345687ee123 | 67070857 | loveneverdies | EMANUELLI | ydz220105 | cap1014 | mdovydas |
| xuexi2010 | 432106969 | u8Aqebj576 | yanjing | 584521584521 | 0167387943 | tigmys2001 |
| daigoro | 6856 | FGYfgy77 | assynt | txudecp | AE86Trueno | denial |
| 12345614 | 704870704870 | 659397 | 62157173 | 84410545 | 19700913 | 678ad5251 |
| DICK4080 | pv041886 | 327296 | 0704224950753 | pietro.chiara | mcsuap | woaiwuai |
| 567891234 | 20060814 | 74748585 | 6903293 | jman1514 | bu56mpbu | 1591591591212 |
| tilg80 | 512881535 | 19720919 | axaaxa | heryarma | danbee | hNbDGN |
| 6z08c861 | milanimilani | 050769585 | hilall | 39joinmam | passw<> | cardcap |
| :zark: | 472619 | nicopa | 30091983 | timelapse | money521 | 13985039393 |
| ravishsneha | dbyxw888 | 2232566 | 2510618981 | mwinkar | conan83 | 001104 |
| 150571611369 | 85717221 | bearss | soukuokpan | 251422 | nxfjpl | desare11 |
| 661189 | cc841215 | n0tpublic | tosecondlife | willrock | rateg143 | 412724198 |
| passme | ariana19321 | isitreal00 | p4os8m6q | YHrtfgDK | kojyihen | nibh1kab |
| trolovinasveta | bbbnnn | ashraf19760 | 015614117 | xys96exq | 058336257 | asferg |
| abdulkhaleque | ang34hehiu | 48144 | acw71790 | mercadotecnia | sarah4444 | hqb555 |
| 007816 | wj112358 | 22471015 | lsyljm2 | 8s5sBEx7 | 7363437 | xgames7 |
| xLDSX | Brenda85 | antyzhou115 | 2xgialdl | 0125040344 | freindship | muckerlee |
| Florida2011 | 786525pb | 0167005246 | gaybar9 | margitka | JytmvW0848 | choqui67 |
| 037037 | shi461988 | ec13kag | 88203009 | omaopa | sb inbau | 12130911 |
| WestC0untry | pingu | 226226226226 | MKltyh87 | dfTi6nh | 30907891 | lierwei120 |
| hitsugaiya | yeybozip | 6767537/33 | quiggle | 1314520521 | 0515043111 | skytdvn |
| 955998126 | 71477nak | mimilebrock | 2063775206 | pixma760 | 1973@ati | milena1995 |
| $3 n 3 r m a x$ | stokurew | gueis8850 | fr3iH3it | pearpear | wlxgjf | kambala11 |

## LEAKED LISTS

Complete left lists from public leaks

| ID | Name | Last Update | Num of Hashes | Progress | Left Hashes | Found |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 6505 | H4v3 1 b33n pwn3d (SHA1) | 02.10.2017-02:03:24 | 320'294'464 | $319^{\prime} 837$ '535 (99.86\%) | Get | Get |
| 5638 | P4y4sUGym (MD5) | 02.10.2017-02:04:19 | 241'266 | 221 '152 (91.66\%) | Get | Get |
| 4920 | L1nk3d1n (SHA1) | 02.10.2017-03:24:58 | 61'829'262 | $60^{\prime} 1477^{\prime} 825$ (97.28\%) | Get | Get |
| 3282 | 4mzr3v13w7r4d3r.c0m (MYSQL5) | 02.10.2017-03:25:32 | 41'823 | $39^{\prime} 166$ (93.65\%) | Get | Get |
| 3186 | X5pl17 (SHA1) | 02.10.2017-03:32:38 | 2'227'254 | $2^{\prime} 162^{\prime} 101$ (97.07\%) | Get | Get |
| 2499 | Hashkiller 32-hex left total | 02.10.2017-11:48:14 | 9'976'651 | 1'723'709 (17.28\%) | Get | Get |
| 2498 | Hashkiller 40-hex left total | 02.10.2017-13:22:34 | $1^{\prime} 739^{\prime} 204$ | $350 ' 788$ (20.17\%) | Get | Get |
| 1619 | 4m4t3urc0mmuni7y.com | 02.10.2017-13:33:26 | 197'302 | $57^{\prime} 407$ (29.1\%) | Get | Get |
| 1535 | b73r.c0m (MD5) | 02.10.2017-13:34:43 | $63^{\prime} 070$ | $32^{\prime} 543$ (51.6\%) | Get | Get |
| 1427 | 4v17r0n.fr | 02.10.2017-13:34:43 | $2^{\prime} 405$ | 2'334 (97.05\%) | Get | Get |
| 1366 | v0d4f0n3 (MD5(\$pass."s+(_a*)" ) | 02.10.2017-13:34:44 | 322 | 307 (95.34\%) | Get | Get |

755
pwned websites

13,044,161,748
pwned accounts

115,769
pastes
paste accounts

## Largest breaches



772,904,991 Collection \#1 accounts
$763,117,241$ Verifications.io accounts
711,477,622 Onliner Spambot accounts
622,161,052 Data Enrichment Exposure From PDL Customer accounts

593,427,119 Exploit.In accounts
509,458,528 Facebook accounts
457,962,538 Anti Public Combo List accounts
393,430,309 River City Media Spam List accounts
tayspace
359,420,698 MySpace accounts
268,765,495 Wattpad accounts

## Recently added breaches

49,102,176 Alleged AT\&T accounts
$3,262,980$ ClickASnap accounts
552,094 Flipkart accounts
3,517,679 Habib's accounts
2,451,197 APK.TW accounts
3,805,265 Online Trade (Онлайн Трейд) accounts

21,994 WoTLabs accounts
27,123 Mr. Green Gaming accounts
19,972,829 Cutout.Pro accounts
243,462 Tangerine accounts

## Password Hashing Functions

Hash functions are very fast to evaluate $\rightarrow$ facilitate fast password cracking
Solution: slow down the guessing process (password "stretching")
Benefit: cracking becomes very inefficient (e.g., 10-100ms per check)
Drawback: increased cost for the server if it must authenticate many users
Make heavy use of available resources
Fast enough computation to validate honest users, but render password guessing infeasible
Adaptable: flexible cost (time/memory complexity) parameters

## Bcrypt [Provos and Mazières, 1999]

Cost-parameterized, modified version of the Blowfish encryption algorithm
Tunable cost parameter (exponential number of loop iterations)

## Alternatives: Scrypt (memory-hard), PBKDF2 (PKCS standard)

## Online Guessing

Similar strategy to offline guessing, but rate-limited
Connect, try a few passwords, get disconnected, repeat...

## Prerequisite: know a valid user name

Credential stuffing: try username + password combinations from previous breaches

## Many failed attempts can lead to a system reaction

Introduce delay before accepting future attempts (exponential backoff)
Shut off completely (e.g., ATM capturing/disabling the card after 3 tries)
Ask user to solve a CAPTCHA

## Very common against publicly accessible SSH, VPN, RDP, and other servers

Main reason people move sshd to a non-default port
Fail2Ban: block IP after many failed attempts $\rightarrow$ attackers may now be able to lock you out Better: disable password authentication altogether and use a key pair $\rightarrow$ cumbersome if having to log in from several devices or others' computers


```
LOGIN: mitch LOGIN: carol
PASSWORD: FooBar!-7 INVALID LOGIN NAME
SUCCESSFUL LOGIN
(a)
```

LOGIN: carol INVALID LOGIN NAME LOGIN:
(b)

LOGIN: carol PASSWORD: Idunno INVALID LOGIN LOGIN:
(c)

```
(a) Successful login
(b) Login rejected after name is entered
(c) Login rejected after name and password are typed \(\rightarrow\) less information makes guessing harder
```


## RouterPasswords.com

Welcome to the internets largets and most updated default router passwords database,

| Select Router Manufacturer: |  | Before guessing, try the default first... |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| CISCO |  |  |  |  |  |
| Find Passwo |  |  |  |  |  |
| Manufacturer | Model | Protocol | Username | Password |  |
| CISCO | CACHE ENGINE | CONSOLE | admin | diamond |  |
| CISCO | CONFIGMAKER | cmaker |  | cmaker |  |
| CISCO | CNR Rev. ALL | CNR GUI | admin | changeme |  |
| CISCO | NETRANGER/SECURE IDS | MULTI | netrangr | attack |  |
| CISCO | BBSM Rev. 5.0 AND 5.1 | TELNET OR NAMED PIPES | bbsd-client | changeme2 |  |
| CISCO | BBSD MSDE CLIENT Rev. $\text { 5.0 AND } 5.1$ | TELNET OR NAMED | bbsd-client | NULL | $\checkmark$ |

## Eavesdropping and Replay

## Physical world

Post-it notes, notebooks, ...
Lift fingerprints (e.g., Apple Touch ID)
Network
Sniffing (LAN, WiFi, ...)
Man-in-the-Middle attacks
Defenses
Encryption
One-time password schemes

## Kerberos Network Authentication Protocol

Most widely used (non-web) single sign-on system
Originally developed at MIT, now used in Unix, Windows, ...

## Long-lived vs. session keys

Use long-lived key for authentication and negotiating session keys
Use "fresh," ephemeral session keys for encrypted communication, MACs, ...
Prevent replay, cryptanalysis, old compromised keys
Authenticate users to services: using their password as the initial key, without having to retype it for every interaction

A Key Distribution Center (KDC) acts as a trusted third party for key distribution
Online authentication: variant of Needham-Schroeder protocol
Assumes a non-trusted network: prevents eavesdropping
Assumes that the Kerberos server and user workstations are secure...
Use cases: workstation login, remote share access, printers, ...

## Password Capture

Hardware bugs/keyloggers
Software keyloggers/malware
Shoulder surfing
Cameras (e.g., ATM skimmers)


KeyGrabber

Social engineering



## Press Ctrl-Alt-Delete to begir.

Requiring this key combination at startup helps keep computer secure. For more information, click Help.

(a) Correct login screen
(b) Phony login screen

## Something You Have: Authentication Tokens

One-time passcode tokens
Time-based or counter-based
Various other authentication tokens


Store certificates, encryption keys, challenge-response, ...
Smartcards (contact or contactless)
Identification, authentication, data storage, limited processing
Magnetic stripe cards, EMV (chip-n-pin credit cards), SIM cards, RFID tags, ...
USB/BLE/NFC tokens, mobile phones, watches, ...
Can be used as authentication devices


## Something You Are: Biometrics

Fingerprint reader


Face recognition
Depth sensing, infrared cameras, ...
Liveness detection (pulse, thermal) to foil simple picture attack
Retina/iris scanner
Voice recognition $\rightarrow$ broken


Related concept: continuous authentication
Keystroke timing, usage patterns, ...
"The probability that a random person the population [sic] could look at your iPhone X and unlock it using Face ID is approximately 1 in 1,000,000 (versus 1 in 50,000 for Touch ID).
For additional protection, Face ID allows only five unsuccessful match attempts before a passcode is required to obtain access to your iPhone.
The probability of a false match is different for twins and siblings that look like you as well as among children under the age of 13, because their distinct facial features may not have fully developed. If you're concerned about this, we recommend using a passcode to authenticate."

## How I Broke Into a Bank Account With an AIGenerated Voice

Banks in the U.S. and Europe tout voice ID as a secure way to log into your account. I proved it's possible to trick such systems with free or cheap AI-generated voices.

By Joseph Cox

T
he bank thought it was talking to me; the AI-generated voice certainly sounded the same.

## Multi-factor Authentication

## Must provide several separate credentials of different types

Most common: two-factor authentication (2FA)
Example: Password + hardware token/SMS message/authenticator app, ...
Example: ATM card + PIN
Motivation: a captured/cracked password is now not enough to compromise a victim's account $\rightarrow$ not always true

Man-in-the-Middle: set up fake banking website, relay password to real website, let the user deal with the second factor...
Man-in-the-Browser: hijack/manipulate an established web session after authentication has been completed (malware, e.g., banking trojans)
Dual infection: compromise both PC and mobile device (rare)
More importantly: the most commonly used 2nd factor (SMS) is the least secure

## SMS Is Not a Secure 2nd Factor

(but still better than no 2nd factor)

## Social engineering

Call victim's mobile operator and hijack the phone number

## Verify

Remember this computer for 30 days. SIM swaping, message/call forwarding, ...

## Message interception

Rogue cell towers: IMSI catchers, StingRays,...
Some phones even display text messages on the lock screen (!)
SS7 attacks
The protocol used for inter-provider signaling is severely outdated and vulnerable Allows attackers to spoof change requests to users' phone numbers and intercept calls or text messages


## A Hacker Got All My Texts for \$16

A gaping flaw in SMS lets hackers take over phone numbers in minutes by simply paying a company to reroute text messages.

By Joseph Cox

March 15, 2021, 1:10pm
f ShareTweet Snap

I hadn't been SIM swapped, where hackers trick or bribe telecom employees to port a target's phone number to their own SIM card. Instead, the hacker used a service by a company called Sakari, which helps businesses do SMS marketing and mass messaging, to reroute my messages to him. This overlooked attack vector shows not only how unregulated commercial SMS tools are but also how there are gaping holes in our telecommunications infrastructure, with a hacker sometimes just having to pinky swear they

The US indicts a Chicago man who allegedly led a SIM-swap gang;
ain Martin / forbes:
Blockchain Capital co-founder Bart Stephens sues a hacker who stole 6.3M in crypto via a SIM-swap attack; $\mathrm{FBI}: \$ 72 \mathrm{M}$ was stolen via SIM
 ug 21, 2023, 1:06 PM-In coneet
Emma Roth / The verge: members stole millions and posed as other people in Apple, I-Mobile,
AT\&T, and Verizon stores - Schene allegedy yargeed ppple, ATIE, veizon, and T-


The US SEC says the January 9 hack of it X a account was via a SIM
swap attack to reset its password; it had disabled 2FA in July 2023 over account access issues - The US. Securities and Exchange Commission sid

CISA releases a report detailing Lapsuss's key techniques, calls for passwordess logins, and asks the FTC and the FCC for stricter SIM


sloomberg
Gary Miller /The Citizen Latik

How members of the Community, a group of teenage SIM swappers
who met on the forum OGUsers, stole millions in crypto in 2018 who met on the forum OGUUsers, stole millions in crypto in 2018
before turning on each other
 Lug 5, 2023,9:45 AM - In conter
David Canellis $/$ he Next Web:
Eederal judge refuses to dismiss $\$ 224 \mathrm{M}$ lawsuit againf ATST for allegedy 2 etting a customer be SIM-swapped twice, leading to the


Indy Greenbery / Wired:
While many foreign phone carriers are sharing real-time SIM swa
data with banks to stop financial fraud, US carriers are dragging the
 Apr $27,2019,12: 00$ PM-In conlex
poseph Cox /VICE:
A look at so-called Russian, encrypted, or "white" SIMs, used by
criminals to spoof phone numbers, add voice manipulation to calls in teal-time, and more - Criminals ses secalleded Russian, enonyputed or or white csus so
 Lug 12, 2020, $11: 15 \mathrm{AM}$ - In con
. vuinerabilities in signaling protocols used by

Oct $29,2023,2: 35 \mathrm{PM}$-In contex

Europol, working with US, UK, and others, says 10 people have been
Europol, working with US, UK, and others, says 10 people have been
arrested for allegedly stealing $\$ 100 \mathrm{M}$ in cryptocurrency from celebr



A collegegestudent whe who solele $\$ 5 M$ in in cryptocurrency via SIM hijacking
gets 10 years in prison and is the first person in the US to be sentenced
3 eb 3, 2019, 7:55 PM-II conext
Brian Kecebs/ Krebs on Security:



Researchers: AT\&T, T-Mobile, Tracfone, US Mobile, and Verizon use vulnerable procedures for customer support that put users at risk of SIM
swapping attacks - Researchess find that 17 of 140 maijo online senices are evineabble to jn 13, 2020, 3:20 PM - In conter
 feb 10, 2021, 2:45 PM - In conter


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& \text { Canada arrests a teenager for allegedly stealing } \$ 36.5 \mathrm{M} \text { in crypto fre }
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\begin{aligned}
& \text { Canada arrests a teenager for allegedly stealing } \$ 36.5 \mathrm{M} \text { in crypto frr } \\
& \text { a US victim using SIM swapping, the largest reported single-person }
\end{aligned}
$$

[^0]
## SMS as 2nd Factor vs. SMS for Account Recovery

Despite its shortcomings, SMS as a 2nd factor is better than nothing
Data point (Google): prevented 100\% of 3.3B automated password stuffing attacks, $96 \%$ of 12 M bulk phishing, and even $76 \%$ of $<10 \mathrm{k}$ targeted attacks seen over a year

Unfortunately, the convenience of phone numbers has led many services to overload SMS as the sole authentication factor

SMS-based onboarding
SMS-based authentication (login with phone number)
SMS-based password reset/account recovery
These are disastrous: a simple SIM-swap attack can take over an account without knowing the password

Password reset via email is much more secure

## Better Alternative: Authenticator Apps

Time-based one-time password (TOTP)
Six/eight digit code provided after password validation HMAC of a shared secret key and the current time
The key is negotiated during registration
Requires "rough" client-server synchronization

dauthy Google Code constantly changes in 30 -second intervals

User-friendly alternative: push notification (e.g., Duo Push) MFA "fatigue" attacks: flood a user with push notifications

More importantly: Phishing is still possible!
The attacker just needs to proxy the captured credentials in real time (rather than collecting them for later use)

## MFA fatigue attacks：Users tricked into allowing device access due to overload of push notifications

Jessica Haworth 16 February 2022 at 15：40 UTC
02 March 2023
Updated： 18 February 2022 at 14：24 UTC

## 노（ in

Social engineering technique confuses victims to gain entry to their accounts
Malicious hackers are targeting Office 365 users with a spare of＇MFA fatigue attacks＇，bombarding victims with 2FA push notifications to trick them into authenticating their login attempts．

This is according to researchers from GoSecure，who have warned that there is an increase in attacks that are exploiting human behavior to gain access to devices．
Multi－factor authentication（MFA）fatigue is the name given to a technique used by adversaries to flood a user＇s authentication app with push notifications in the hope they will accept and therefore enable an attacker to gain entry to an account or device．

## Uber：Lapsus\＄Targeted External Contractor With MFA Bombing Attack

The ride－sharing giant says a member of the notorious Lapsus\＄hacking group started the attack by compromising an external contractor＇s credentials，as researchers parse the incident for takeaways．

Jai Vijayan，Contributing Writer September 19， 2022

Uber has attributed last week＇s massive breach at U Lapsus\＄hacking group and released additional deta Researchers say the incident has highlighted the ris trusting too much in multifactor authentication（MFA risk around cloud－service adoption．

The attacker then repeatedly tried to log in to the Uber account using the illegally obtained credentials，prompting a two－factor login approval request each time．After the contractor initially blocked those requests，the attacker contacted the target on WhatsApp posing as tech support，telling the person to accept the MFA prompt－thus allowing the attacker to log in．

[^1]
## Evilginx2 https://github.com/kgretzky/evilginx2

Man-in-the-middle attack framework for phishing login credentials along with session cookies

Bypasses 2-factor authentication
No need for HTML templates: just a web proxy
Victim's traffic is forwarded to the real website
TLS termination at the proxy (e.g., using a LetsEncrypt certificate)


## Google

## Sign in

with your Google Account

Email or phone

Forgot email?

Not your computer? Use Guest mode to sign in privately. Learn more

Create account
NEXT

English (United States)
Help
Privacy
Terms

## Even Better Alternative: U2F Tokens (AKA Security Keys)

Universal Second Factor (U2F)
FIDO (Fast IDentity Online) alliance: Google, Yubico, ...
Supported by all popular browsers and many online services
A different key pair is generated for each origin during registration
Origin = <protocol, hostname, port>
Private key stored re-generated on device Public key sent to server

Additions to the authentication flow:
Origin (URI): prevents phishing
TLS Channel ID (optional): prevents MitM


done!


## Key Generation

Storing a private key + metadata per service would require a lot of storage Alternative: store only a master symmetric key

Generated on-device upon first startup, and never leaves the YubiKey in any form

## Registration

YubiKey generates a random key pair per credential
YubiKey encrypts the private key + metadata with the master key $\rightarrow$ key handle
Key handle + public key sent to server

## Authentication

The server presents the key handle to the YubiKey, along with a challenge
YubiKey decrypts the key handle and reveals the private key (authenticated encryption: ensures integrity, and that the credential is used with the correct AppID)
YubiKey signs the challenge with the private key to complete the authentication

## U2F tokens

## Benefits

Easy: just tap the button (no typing)
Works out of the box (no drivers to install)
USB, NFC, Bluetooth communication
No shared secret between client and server
Origin checking $\rightarrow$ prevents phishing!

## Drawbacks

Can be lost $\rightarrow$ need a fallback (backup codes, 2nd U2F token, authenticator app, ...)
Cumbersome: have to pull keychain out and plug token in (or have an always pugged-in token, in which case though it can be stolen along with the device)
Cost (\$10-\$70)

## －

## Google＇s strongest security helps keep your private information safe．

The Advanced Protection Program safeguards users with high visibility and sensitive information from targeted online attacks．New protections are automatically added to defend against today＇s wide
range of threats．

自 G Advanced Protection


Get security keys
First，you need 2 security keys，one of them for backup．Your security
key will be used in addition to your password to sign in to your accoun
You can use keys that you already own or buy new ones．Learn more
Ship to：United States ＊


Make sure to
your devices．

Buy now

Because you use a physical key instead of the six-digit code, security keys strengthen the two-factor authentication process and help prevent your second authentication factor from being intercepted or requested by an attacker.

> You're responsible for maintaining access to your security keys. If you lose all of your trusted devices and security keys, you could be locked out of your account permanently.

Learn more about two-factor authentication >

## What's required for Security Keys for Apple ID

- At least two FIDO ${ }^{\circledR}$ Certified* security keys that work with the Apple devices that you use on a regular basis.
- iOS 16.3, iPadOS 16.3, or macOS Ventura 13.2, or later on all of the devices where you're signed in with your Apple ID.
- Two-factor authentication set up for your Apple ID.
- A modern web browser. If you can't use your security key to sign in on the web, update your browser to the latest version or try another browser.

2FA Recap - What threats does it prevent?
SMS: useful against two main threats
Credential stuffing (people tend to reuse passwords across different services)
Leaked passwords (post-it, hardware keyloggers, cameras, shoulder surfing, ...)
Introduces new security/privacy issues: SIM swapping, SMS account recovery, SMS spam...

## Authenticator Apps/Push Auth: much better alternative than SMS

Protects against the same threats without relying on phone numbers

## U2F: additional protection against phishing

Modern phishing toolkits bypass SMS/Authenticator/Push 2FA through MitM
Humans fall for typosquatting, but U2F's origin check doesn't

## None of the above protect against session hijacking and Man-in-the-Browser

Game over anyway if the host is compromised after the user has successfully logged in

## Password Managers

Have become indispensable
Encourage the use of complex/non-memorable passwords
Obviate the need for password reuse: unique passwords per site/service
Protection against phishing: auto-fill won't work for incorrect domains
As long as users don't copy/paste passwords out of the password manager (!)
Various options: third-party applications, OS-level, in-browser
Password synchronization across devices
Can the service provider access all my passwords or not?
Preferable option: passwords should be encrypted locally with a master password never visible to the cloud service

Single point of failure (!)
$\square$

## -ars TECHNICA

## LastPass says employee's home computer was hacked and corporate vault taken

Already smarting from a breach that stole customer vaults, LastPass has more bad news.

DAN GOODIN - 2/27/2023, 8:01 PM

Already smarting from a breach that put partially encrypted login data into a threat actor's hands, LastPass on Monday said that the same attacker hacked an employee's home computer and obtained a decrypted vault available to only a handful of company developers.


Although an initial intrusion into LastPass ended on August 12, officials with the leading password manager said the threat actor "was actively engaged in a new series of reconnaissance, enumeration, and exfiltration activity" from August 12 to August 26 . In the process, the unknown threat actor was able to steal valid credentials from a senior DevOps engineer and access the contents of a LastPass data vault. Among other things, the vault gave access to a shared cloud-storage environment that contained the encryption keys for customer vault backups stored in Amazon S3 buckets.

## Another bombshell drops

## Single Sign-on/Social Login

Use a central authentication service for multiple sites Pros

Convenience: fewer passwords to remember
Easier development: outsource user registration/management
Rich experience through social features

## Cons

Same credentials for multiple sites: single point of failure
Third-parties gain access to users' profiles
Provider can track users

## WebAuthn

## W3C Web Authentication standard (FIDO2): Successor of FIDO U2F

## Use cases

Low friction and phishing-resistant 2FA (in conjunction with a password)
Passwordless, biometrics-based re-authorization
2FA without a password (passwordless login)
Authenticators: devices that can generate private/public key pairs and gather consent (simple tap, fingerprint read, ...)

Roaming Authenticators:
USB/BLE/NFC security keys
Platform Authentications:
Built-in fingerprint readers, cameras, ...


Relying party



## Passkeys

Completely replace passwords with cryptographic key pairs
Server only keeps a user's public key
Based on WebAuthn: rely on biometric identification (Face ID, Windows Hello, ...)
Key enabler: identity providers (Apple, Google, ...) who also sell devices
The device becomes an authenticator: what if it gets lost? $\rightarrow$ recovery through vendor Users have more than one device $\rightarrow$ seamless syncing


SAFETY \＆SECURITY

## Passwordless by default：Make the switch to passkeys

Oct 10， 2023Christiaan Brand Group Product Manager


## Passkeys.directory

Passkeys.directory is a community-driven index of websites, apps, and services
that offer signing in with passkeys.


## Multi-factor vs. Multi-step

Factor: something you know/have/are
Step: user-specific action
Type password, tap fingerprint reader, press security key, look at camera, ...
Example: U2F flow with passwords
Type password + tap security key $\rightarrow$ two factors, two steps
Example: FIDO2 passwordless flow
Tap biometric security key $\rightarrow$ two factors, one step
Phone Face ID $\rightarrow$ two factors, one step

## Recap: Crypto-based Authentication

Rely on a cryptographic key to prove a user's identity
User performs a requested cryptographic operation on a value (challenge) that the verifier supplies

Usually based on knowledge of a key (shared secret key or private key)
Can use symmetric (e.g., Kerberos) or public key (e.g., U2F, passkeys) schemes
How can we trust a key? Why is it authentic?
Need to establish a level of trust
Different approaches: TOFU, PKI, Web of Trust

## Trust on First Use (aka Key Continuity)

## Use case: SSH

Performs mutual authentication

## Server always authenticates the client

password, key pair, ...
Client almost always authenticates the server - except the first time!
First connection: server presents its public key
No other option for the user but to accept it: MitM opportunity
Subsequent connections: client remembers server's key, and triggers an alert on key mismatch
Pragmatic solution, but shifts the burden to users
Users must determine the validity of the presented key
Accepting a key change without verifying the new key offers no protection against MitM (unfortunately, that's what most users do)

@ WARNING: REMOTE HOST IDENTIFICATION HAS CHANGED! @
@@@@@@@@@cccccccccccccccccccccccccccccccccccccccccccccccccc
IT IS POSSIBLE THAT SOMEONE IS DOING SOMETHING NASTY!
Someone could be eavesdropping on you right now (man-in-the-middle attack)!
It is also possible that the RSA host key has just been changed.
The fingerprint for the RSA key sent by the remote host is
df:c8:52:aa:cd:e3:da:8c:ec:50:46:db:4d:21:d9:c7.
Please contact your system administrator.
Add correct host key in /root/.ssh/known hosts to get rid of this message.
Offending key in /root/.ssh/known hosts: $\overline{1}$
RSA host key for 192.168.2.5 has changed and you have requested strict checking. Host key verification failed.

## TODAY

Messages you send to this chat and calls are now secured with end-to-end encryption. Tap for more info.

This is a normal message in a normal conversation.

Now Alice is going to reinstall. 12:31 PM //

Alice's security code changed. Tap for more info.

As soon as Alice reinstalled, I saw the notice above. Impressive.

Now Alice has uninstalled, and this message is being transmitted before Alice reinstalls.

12:34 PM
Alice's security code changed. Tap for more infoType a message


```
56890}599295 61701 15415
38897 13310 80072 75067
50646 41640 61012 94324
```

Scan the code on your contact's phone, or ask them to scan your code, to verify that your messages and calls to them are end-to-end encrypted. You can also compare the number above to verify. This is optional. Learn more.

## Certificates

## How can we distribute "trusted" public keys?

Public directory $\rightarrow$ risk of forgery and tampering, scalability issues
More practical solution: "certified" public keys
A certificate is a digitally signed message that contains an identity and a public key

Makes an association between a user/entity and a private key
Valid until a certain period
Most common format: X. 509
Why trust a certificate?


Because it is signed by an "authority"
Requiring a signature by a third party prevents straightforward tampering

## Public Key Infrastructures (PKI)

Facilitate the authentication and distribution of public keys with the respective identities of entities

People, organizations, devices, applications, ...
Set of roles, policies, hardware, software, and procedures to create, mange, distribute, use, store, and revoke digital certificates and manage public key encryption

An issuer signs certificates for subjects: "Trust anchor"
Methods of certification
Certificate authorities (hierarchical structure - root of trust)
Web of trust (decentralized, peer-to-peer structure)

## Certificate Authorities

Trusted third-parties responsible for certifying public keys
Most CAs are tree-structured
A public key for any website in the world will be accepted without a browser warning if it has been certified by a trusted CA

Why should we trust an authority?
How do we know the public key of the Certificate Authority?
CA's public key (trust anchor) must somehow be provided out of band
Operating systems and browsers are pre-configured with tens/hundreds of trusted root certificates (more on this in the TLS lecture)

Single point of failure: CAs can be compromised!


A Dutch certificate authority that suffered a major hack attack this summer has been unable to recover from the blow and filed for bankruptcy this week.

## Web of Trust (mainly used in PGP for encrypted email - future lecture)

## Entirely decentralized authentication

No need to buy certs from CAs: users create their own certificates
Users validate other users' certificates, forming a "web of trust"
No trusted authorities: trust is established through friends (yay! key signing parties!)

## Main problems

Privacy issues: social graph metadata
Bootstrapping: new users are not readily trusted by others
When opinions vary, "stronger set" wins: impersonation through collusion/compromised keys
Scalability: challenging to create a WoT for the whole world

## WoT: Finding Public Keys

Public PGP key servers
pgp.mit.edu
keyserver.pgp.com
Cache certificates from received emails
Integration with user management systems (LDAP, IAM/IDP)
Ad-hoc approaches
List public key on home page
Print on business card
Exchange through another medium on a case-by-case basis
Association with social profiles/identities
keybase.io

## Online Social "Tracking"



## Keybase.io

In essence, a directory associating public keys with names
Identity established through public signatures
Identity proofs: "I am Joe on Keybase and MrJoe on Twitter"
Follower statements: "I am Joe on Keybase and I just looked at Chris's identity"
Key ownership: "I am Joe on Keybase and here's my public key"
Revocations: "I take back what I said earlier"
Keybase identity $=$ sum of public identities
Twitter, Facebook, Github, Reddit, domain ownership, ...
michalis @polychronakis • 28 Aug 2014
Verifying myself: I am mikepo on Keybase.io. NpbEbc8BJOrT4k70TcmM2oA4G24IXVNt89R /

An attacker has to compromise all connected identities
The more connected identities, the harder to impersonate a user

## Best Practices

## Use long passphrases instead of passwords

Never reuse the same password on different services

## Use two-factor authentication

Avoid SMS if possible! Use an authenticator app or even better U2F (or passkeys)
Remove phone number from account after authenticator/U2F setup
Store your backup codes/backup key in a safe location
Use a password manager
Pick non-memorable passwords and avoid copy/pasting them
Password auto-fill helps against phishing! (auto-fill will fail if the domain is wrong)
Use SSH keys instead of passwords


[^0]:    Europol, alongside Italian and Spanish police, arrest 106 people

    ccused of working for the Italian Mafia and laundering over $£ 10 \mathrm{M}$ nade through cybercrimes - European poicee accused sereari people of SIM | mepping phishing and hacking in |
    | :--- |
    | Sep 20, 2021, $11: 00$ AM - I contex |

[^1]:    In an update on Monday，Uber laid out the attribution：＂We believe that this attacker（or attackers）are affiliated with a hacking group called Lapsus\＄，which has been increasingly active over the last year or so．＂Uber＇s announcement pointed to other companies that had been targeted by the notorious gang via similar techniques，including Cisco，Microsoft，Nvidia，Okta， and Samsung，

