

SML

CSE 307 – Principles of Programming Languages

Stony Brook University

<http://www.cs.stonybrook.edu/~cse307>

Functional Programming

- *Function evaluation* is the basic concept for a programming paradigm that has been implemented in *functional programming languages*
- The language ML (“Meta Language”) was originally introduced in 1977 as part of a theorem proving system, and was intended for describing and implementing proof strategies in the Logic for Computable Functions (LCF) theorem prover (whose language, *pplambda*, a combination of the first-order predicate calculus and the simply typed polymorphic lambda calculus, had ML as its metalanguage)
- Standard ML of New Jersey (SML) is an implementation of ML
 - The basic mode of computation in SML is the use of the definition and application of functions

Install Standard ML

- Download from:
 - <http://www.smlnj.org>
- Start Standard ML:
 - Type **sml** from the shell (run command line in Windows)
- Exit Standard ML:
 - **Ctrl-Z** under Windows
 - **Ctrl-D** under Unix/Mac

Standard ML

- The basic cycle of SML activity has three parts:
 - read input from the user
 - evaluate it
 - print the computed value (or an error message)

First SML example

- SML prompt:

-

- Simple example:

- 3;

```
val it = 3 : int
```

- The first line contains the SML prompt, followed by *an expression* typed in by the user and ended by a *semicolon*
- The second line is SML's response, indicating the *value* of the input expression and its *type*

Interacting with SML

- SML has a number of built-in operators and data types.
- it provides the standard arithmetic operators

```
- 3+2;
```

```
val it = 5 : int
```

- The boolean values **true** and **false** are available, as are logical operators such as: **not** (negation), **andalso** (conjunction), and **orelse** (disjunction)

```
- not(true);
```

```
val it = false : bool
```

```
- true andalso false;
```

```
val it = false : bool
```

Types in SML

- As part of the evaluation process, SML determines the type of the output value using methods of *type inference*.
- Simple types include **int**, **real**, **bool**, and **string**
- One can also associate identifiers with values

```
- val five = 3+2;  
val five = 5 : int
```

and thereby establish a new value binding

```
- five;  
val it = 5 : int
```

Function Definitions in SML

- The general form of a function definition in SML is:

```
fun <identifier> (<parameters>) =  
    <expression>;
```

- For example,

```
- fun double(x) = 2*x;
```

```
val double = fn : int -> int
```

declares **double** as a function from integers to integers, i.e., of type **int** \rightarrow **int**

- Apply a function to an argument of the wrong type results in an error message:

```
- double(2.0);
```

```
Error: operator and operand don't agree ...
```


Function Definitions in SML

- The user may also **explicitly** indicate types:

```
- fun max(x:int,y:int,z:int):int =  
    if ((x>y) andalso (x>z)) then x  
    else (if (y>z) then y else z);
```

```
val max = fn : int * int * int -> int
```

```
- max(3,2,2);
```

```
val it = 3 : int
```

Recursive Definitions

- The use of recursive definitions is a main characteristic of functional programming languages, and these languages encourage the use of recursion over iterative constructs such as while loops:

```
- fun factorial(x) = if x=0 then 1  
  else x*factorial(x-1);
```

```
val factorial = fn : int -> int
```

- The definition is used by SML to evaluate applications of the function to specific arguments:

```
- factorial(5);
```

```
val it = 120 : int
```

```
- factorial(10);
```

```
val it = 3628800 : int
```

Example: Greatest Common Divisor

- The greatest common divisor (gcd) of two positive integers can be defined recursively based on the following observations:

$$\text{gcd}(n, n) = n,$$

$$\text{gcd}(m, n) = \text{gcd}(n, m), \text{ if } m < n, \text{ and}$$

$$\text{gcd}(m, n) = \text{gcd}(m - n, n), \text{ if } m > n.$$

- These identities suggest the following recursive definition:

```
- fun gcd(m,n):int = if m=n then n
  else if m>n then gcd(m-n,n)
  else gcd(m,n-m);
```

```
val gcd = fn : int * int -> int
```

```
- gcd(12,30);      - gcd(1,20);      - gcd(125,56345);
```

```
val it = 6 : int   val it = 1 : int   val it = 5 : int
```

Basic operators on the integers

<i>op</i>	:	<i>type</i>	<i>form</i>	<i>precedence</i>
+	:	$\text{int} \times \text{int} \rightarrow \text{int}$	infix	6
-	:	$\text{int} \times \text{int} \rightarrow \text{int}$	infix	6
*	:	$\text{int} \times \text{int} \rightarrow \text{int}$	infix	7
div	:	$\text{int} \times \text{int} \rightarrow \text{int}$	infix	7
mod	:	$\text{int} \times \text{int} \rightarrow \text{int}$	infix	7
=	:	$\text{int} \times \text{int} \rightarrow \text{bool}^*$	infix	4
<>	:	$\text{int} \times \text{int} \rightarrow \text{bool}^*$	infix	4
<	:	$\text{int} \times \text{int} \rightarrow \text{bool}$	infix	4
<=	:	$\text{int} \times \text{int} \rightarrow \text{bool}$	infix	4
>	:	$\text{int} \times \text{int} \rightarrow \text{bool}$	infix	4
>=	:	$\text{int} \times \text{int} \rightarrow \text{bool}$	infix	4
~	:	$\text{int} \rightarrow \text{int}$	prefix	
abs	:	$\text{int} \rightarrow \text{int}$	prefix	

- The infix operators associate to the left
- The operands are always all evaluated

Basic operators on the reals

<i>op</i>	:	<i>type</i>	<i>form</i>	<i>precedence</i>
+	:	real \times real \rightarrow real	infix	6
-	:	real \times real \rightarrow real	infix	6
*	:	real \times real \rightarrow real	infix	7
/	:	real \times real \rightarrow real	infix	7
=	:	real \times real \rightarrow bool *	infix	4
<>	:	real \times real \rightarrow bool *	infix	4
<	:	real \times real \rightarrow bool	infix	4
<=	:	real \times real \rightarrow bool	infix	4
>	:	real \times real \rightarrow bool	infix	4
>=	:	real \times real \rightarrow bool	infix	4
~	:	real \rightarrow real	prefix	
abs	:	real \rightarrow real	prefix	
Math.sqrt	:	real \rightarrow real	prefix	
Math.In	:	real \rightarrow real	prefix	

unary operator – is represented by ~

Type conversions

<i>op</i>	:	<i>type</i>
real	:	int \rightarrow real
ceil	:	real \rightarrow int
floor	:	real \rightarrow int
round	:	real \rightarrow int
trunc	:	real \rightarrow int

```
- real(2) + 3.5 ;  
val it = 5.5 : real  
- ceil(23.65) ;  
val it = 24 : int  
- ceil(~23.65) ;  
val it = ~23 : int  
- floor(23.65) ;  
val it = 23 : int
```

More recursive functions

```
- fun exp(b,n) = if n=0 then 1.0  
    else b * exp(b,n-1);  
val exp = fn : real * int -> real
```

```
- exp(2.0,10);  
val it = 1024.0 : real
```

Tuples in SML

- In SML tuples are finite sequences of arbitrary but fixed length, where different components need not be of the same type

```
- (1, "two");
```

```
val it = (1,"two") : int * string
```

```
- val t1 = (1,2,3);
```

```
val t1 = (1,2,3) : int * int * int
```

```
- val t2 = (4, (5.0,6));
```

```
val t2 = (4, (5.0,6)) : int * (real * int)
```

- The components of a tuple can be accessed by applying the built-in functions `#i`, where `i` is a positive number

```
- #1(t1);
```

```
val it = 1 : int
```

```
- #2(t2);
```

```
val it = (5.0,6) : real * int
```

If a function `#i` is applied to a tuple with fewer than `i` components, an error results.

Tuples in SML

- Functions using tuples should completely define the type of tuples, otherwise SML cannot detect the type, e.g., nth argument

```
- fun firstThird(Tuple:'a * 'b * 'c):'a * 'c =  
    (#1(Tuple), #3(Tuple));
```

```
val firstThird = fn : 'a * 'b * 'c -> 'a * 'c
```

```
- firstThird((1,"two",3));
```

```
val it = (1,3) : int * int
```

- Without types, we would get an error:

```
- fun firstThird(Tuple) = (#1(Tuple), #3(Tuple));
```

```
stdIn: Error: unresolved flex record (need to know the  
names of ALL the fields in this context)
```

Polymorphic functions

```
- fun id x = x;
```

```
val id = fn : 'a -> 'a
```

```
- (id 1, id "two");
```

```
val it = (1,"two") : int * string
```

```
- fun fst(x,y) = x;
```

```
val fst = fn : 'a * 'b -> 'a
```

```
- fun snd(x,y) = y;
```

```
val snd = fn : 'a * 'b -> 'b
```

```
- fun switch(x,y) = (y,x);
```

```
val switch = fn : 'a * 'b -> 'b * 'a
```

Polymorphic functions

- ' **a** means "any type", while ' ' **a** means "any type that can be compared for equality" (see the **concat** function later which compares a polymorphic variable list with **[]**)
- There will be a "Warning: calling polyEqual" that means that you're comparing two values with polymorphic type for equality
 - Why does this produce a warning? Because it's less efficient than comparing two values of known types for equality
 - How do you get rid of the warning? By changing your function to only work with a specific type instead of any type
 - Should you do that or care about the warning? Probably not. In most cases having a function that can work for any type is more important than having the most efficient code possible, so you should just ignore the warning.

Lists in SML

- A list in SML is a finite sequence of objects, all of the same type:

- `[1,2,3];`

```
val it = [1,2,3] : int list
```

- `[true,false,true];`

```
val it = [true,false,true] : bool list
```

- `[[1,2,3],[4,5],[6]];`

```
val it = [[1,2,3],[4,5],[6]] :  
          int list list
```

- The last example is **a list of lists of integers**

Lists in SML

- All objects in a list must be of the same type:

- `[1, [2]];`

Error: operator and operand don't agree

- An empty list is denoted by one of the following expressions:

- `[];`

`val it = [] : 'a list`

- `nil;`

`val it = [] : 'a list`

- Note that the type is described in terms of a type variable `'a`. Instantiating the type variable, by types such as `int`, results in (different) empty lists of corresponding types

Operations on Lists

- SML provides various functions for manipulating lists
 - The function **hd** returns the first element of its argument list

```
- hd[1,2,3];
```

```
val it = 1 : int
```

```
- hd[[1,2],[3]];
```

```
val it = [1,2] : int list
```

Applying this function to the empty list will result in an error.

- The function **tl** removes the first element of its argument lists, and returns the remaining list

```
- tl[1,2,3];
```

```
val it = [2,3] : int list
```

```
- tl[[1,2],[3]];
```

```
val it = [[3]] : int list list
```

- The application of this function to the empty list will also result in an error

Operations on Lists

- Lists can be constructed by the (binary) function `::` (read *cons*) that adds its first argument to the front of the second argument.

```
- 5 :: [];
```

```
val it = [5] : int list
```

```
- 1 :: [2,3];
```

```
val it = [1,2,3] : int list
```

```
- [1,2] :: [[3],[4,5,6,7]];
```

```
val it = [[1,2],[3],[4,5,6,7]] : int list list
```

- **IMPORTANT:** The arguments must be of the right type (such that the result is a list of elements of the same type):

```
- [1] :: [2,3];
```

```
Error: operator and operand don't agree
```

Operations on Lists

- Lists can also be compared for equality:

```
- [1,2,3]=[1,2,3];
```

```
val it = true : bool
```

```
- [1,2]=[2,1];
```

```
val it = false : bool
```

```
- tl[1] = [];
```

```
val it = true : bool
```


Defining List Functions

- Recursion is particularly useful for defining functions that process lists
- For example, consider the problem of defining an SML function that takes as arguments two lists of the same type and returns the concatenated list.
- In defining such list functions, it is helpful to keep in mind that a list is either
 - an empty list `[]` or
 - of the form `x :: y`

Concatenation

- In designing a function for concatenating two lists \mathbf{x} and \mathbf{y} we thus distinguish two cases, depending on the form of \mathbf{x} :
 - If \mathbf{x} is an empty list $[\]$, then concatenating \mathbf{x} with \mathbf{y} yields just \mathbf{y} .
 - If \mathbf{x} is of the form $\mathbf{x1} :: \mathbf{x2}$, then concatenating \mathbf{x} with \mathbf{y} is a list of the form $\mathbf{x1} :: \mathbf{z}$, where \mathbf{z} is the result of concatenating $\mathbf{x2}$ with \mathbf{y} .
 - We can be more specific by observing that
$$\mathbf{x} = \mathbf{x1} :: \mathbf{x2} = \text{hd}(\mathbf{x}) :: \text{tl}(\mathbf{x})$$

Concatenation

```
- fun concat(x,y) = if x=[] then y  
  else hd(x)::concat(tl(x),y);
```

```
val concat = fn : 'a list * 'a list -> 'a list
```

- Applying the function yields the expected results:

```
- concat([1,2],[3,4,5]);
```

```
val it = [1,2,3,4,5] : int list
```

```
- concat([], [1,2]);
```

```
val it = [1,2] : int list
```

```
- concat([1,2], []);
```

```
val it = [1,2] : int list
```

Length

- The following function computes the length of its argument list:

```
- fun length(L) = if (L=nil) then 0  
                  else 1+length(tl(L));
```

```
val length = fn : 'a list -> int
```

```
- length[1,2,3];
```

```
val it = 3 : int
```

```
- length[[5],[4],[3],[2,1]];
```

```
val it = 4 : int
```

```
- length[];
```

```
val it = 0 : int
```

doubleall

- The following function doubles all the elements in its argument list (of integers):

```
- fun doubleall(L) =  
    if L=[] then []  
    else (2*hd(L))::doubleall(tl(L));  
val doubleall = fn : int list -> int list  
  
- doubleall([1,3,5,7]);  
val it = [2,6,10,14] : int list
```

Reversing a List

```
- fun reverse(L) =  
    if L = nil then nil  
    else concat(reverse(tl(L)), [hd(L)]);  
val reverse = fn : ''a list -> ''a list
```

```
- reverse [1,2,3];
```

calls

```
- concat(reverse([2,3]), [1])
```

```
- concat([3,2], [1]);
```

```
val it = [3,2,1] : int list
```

Reversing a List

- Concatenation of lists, for which we gave a recursive definition, is actually a built-in operator in SML, denoted by the symbol @
 - We can use this operator in reversing:

```
- fun reverse(L) =  
    if L = nil then nil  
    else reverse(tl(L)) @ [hd(L)];  
val reverse = fn : 'a list -> 'a list  
- reverse [1,2,3];  
val it = [3,2,1] : int list
```

Reversing a List

```
- fun reverse(L) =  
    if L = nil then nil  
    else concat(reverse(tl(L)), [hd(L)]);
```

This method is not efficient: $O(n^2)$

$$\begin{aligned}T(N) &= T(N-1) + (N-1) = \\ &= T(N-2) + (N-2) + (N-1) = \\ &= 1 + 2 + 3 + \dots + N-1 = N * (N-1) / 2\end{aligned}$$

Reversing a List

- This way (using an accumulator) is better: $O(n)$

```
- fun reverse_helper(L, L2) =  
  if L = nil then L2  
  else reverse_helper(tl(L), hd(L) :: L2);  
- fun reverse(L) = reverse_helper(L, []);  
- reverse [1,2,3];  
- reverse_helper([1,2,3], []);  
- reverse_helper([2,3], [1]);  
- reverse_helper([3], [2,1]);  
- reverse_helper([], [3,2,1]);
```

```
[3,2,1]
```

Removing List Elements

- The following function **removes all occurrences** of its first argument from its second argument list

```
- fun remove(x,L) = if (L=[]) then []  
    else if x=hd(L) then remove(x,tl(L))  
    else hd(L)::remove(x,tl(L));  
val remove = fn : 'a * 'a list -> 'a list
```

```
- remove(1,[5,3,1]);  
val it = [5,3] : int list
```

```
- remove(2,[4,2,4,2,4,2,2]);  
val it = [4,4,4] : int list
```

Removing Duplicates

- The remove function can be used in the definition of another function that **removes all duplicate occurrences** of elements from its argument list:

```
- fun removedupl(L) =  
  if (L=[]) then []  
  else hd(L)::removedupl(remove(hd(L),tl(L)));  
val removedupl = fn : 'a list -> 'a list
```

```
- removedupl([3,2,4,6,4,3,2,3,4,3,2,1]);  
val it = [3,2,4,6,1] : int list
```

Definition by Patterns

- In SML functions can also be defined via patterns.
 - The general form of such definitions is:

```
fun <identifier>(<pattern1>) = <expression1>  
| <identifier>(<pattern2>) = <expression2>  
| ...  
| <identifier>(<patternK>) = <expressionK>;
```

where the identifiers, which name the function, are all the same, all patterns are of the same type, and all expressions are of the same type.

- Example:

```
- fun reverse(nil) = nil  
  | reverse(x::xs) = reverse(xs) @ [x];  
val reverse = fn : 'a list -> 'a list
```

The patterns are inspected in order and the first match determines the value of the function.

Sets with lists in SML

```
fun member(X,L) =  
  if L=[] then false  
  else if X=hd(L) then true  
  else member(X,tl(L));
```

OR with patterns:

```
fun member(X,[]) = false  
  | member(X,Y::Ys) =  
    if (X=Y) then true  
    else member(X,Ys);  
member(1,[1,2]); (* true *)  
member(1,[2,1]); (* true *)  
member(1,[2,3]); (* false *)
```

Sets UNION

```
fun union(L1,L2) =  
  if L1=[] then L2  
  else if member(hd(L1),L2)  
        then union(tl(L1),L2)  
        else hd(L1)::union(tl(L1),L2);
```

```
union([1,5,7,9],[2,3,5,10]);
```

```
(* [1,7,9,2,3,5,10] *)
```

```
union([], [1,2]); (* [1,2] *)
```

```
union([1,2], []); (* [1,2] *)
```

Sets UNION with patterns

```
fun union([],L2) = L2
  | union(X::Xs,L2) =
    if member(X,L2) then union(Xs,L2)
    else X::union(Xs,L2);
```

Sets Intersection (\cap)

```
fun intersection(L1,L2) =  
  if L1=[] then []  
  else if member(hd(L1),L2)  
  then hd(L1)::intersection(tl(L1),L2)  
  else intersection(tl(L1),L2);  
  
intersection([1,5,7,9],[2,3,5,10]);  
(* [5] *)
```


Sets \cap with patterns

```
fun intersection([], L2) = []  
  | intersection(L1, []) = []  
  | intersection(X::Xs, L2) =  
    if member(X, L2)  
    then X::intersection(Xs, L2)  
    else intersection(Xs, L2);
```

Sets subset

```
fun subset(L1,L2) = if L1=[] then true
  else if L2=[] then false
  else if member(hd(L1),L2)
    then subset(tl(L1),L2)
  else false;
```

```
subset([1,5,7,9],[2,3,5,10]); (* false *)
subset([5],[2,3,5,10]);      (* true  *)
```

Sets subset patterns

```
fun subset([],L2) = true
  | subset(L1,[]) = if(L1=[])
                    then true
                    else false
  | subset(X::Xs,L2) =
    if member(X,L2)
      then subset(Xs,L2)
      else false;
```

Sets equal

```
fun setEqual(L1, L2) =  
    subset(L1, L2) andalso subset(L2, L1);
```

```
setEqual([1, 5, 7], [7, 5, 1, 2]); (* false *)
```

```
setEqual([1, 5, 7], [7, 5, 1]); (* true *)
```

Set difference patterns

```
fun minus([],L2) = []  
  | minus(X::Xs,L2) =  
    if member(X,L2)  
      then minus(Xs,L2)  
      else X::minus(Xs,L2);
```

```
minus([1,5,7,9],[2,3,5,10]);  
(* [1,7,9] *)
```

Sets Cartesian product

```
fun product_one(X, []) = []  
  | product_one(X, Y :: Ys) =  
    (X, Y) :: product_one(X, Ys) ;  
product_one(1, [2, 3]) ;  
(* [(1, 2), (1, 3)] *)  
fun product([], L2) = []  
  | product(X :: Xs, L2) =  
    union(product_one(X, L2),  
          product(Xs, L2)) ;  
product([1, 5, 7, 9], [2, 3, 5, 10]) ;  
(* [(1, 2), (1, 3), (1, 5), (1, 10), (5, 2),  
    (5, 3), (5, 5), (5, 10), (7, 2), (7, 3), ...] *)
```

Sets Powerset

```
fun insert_all(E,L) =
  if L=[] then []
  else (E::hd(L)) :: insert_all(E,tl(L));
insert_all(1, [[]], [2], [3], [2,3]);
(* [ [1], [1,2], [1,3], [1,2,3] ] *)
fun powerSet(L) =
  if L=[] then [[]]
  else powerSet(tl(L)) @
    insert_all(hd(L), powerSet(tl(L)));
powerSet([]);
powerSet([1,2,3]);
powerSet([2,3]);
```

Higher-Order Functions

- In functional programming languages functions (called *first-class functions*) can be used as parameters or return value in definitions of other (called *higher-order*) functions
 - The following function, **map**, applies its first argument (a function) to all elements in its second argument (a list of suitable type):

```
- fun map(f,L) = if (L=[]) then []  
  else f(hd(L)) :: (map(f,tl(L)));
```

```
val map = fn : ('a -> 'b) * 'a list -> 'b list
```

- We may apply **map** with any function as argument:

```
- fun square(x) = (x:int)*x;
```

```
val square = fn : int -> int
```

```
- map(square,[2,3,4]);
```

```
val it = [4,9,16] : int list
```


Higher-Order Functions

- Anonymous functions:

```
- map (fn x=>x+1, [1,2,3,4,5]);
```

```
val it = [2,3,4,5,6] : int list
```

```
- fun incr(list) = map (fn x=>x+1, list);
```

```
val incr = fn : int list -> int list
```

```
- incr[1,2,3,4,5];
```

```
val it = [2,3,4,5,6] : int list
```

McCarthy's 91 function

- McCarthy's 91 function:

```
- fun mc91(n) = if n>100 then n-10  
  else mc91(mc91(n+11));
```

```
val mc91 = fn : int -> int
```

```
- map mc91 [101, 100, 99, 98, 97, 96];
```

```
val it = [91,91,91,91,91,91] : int list
```

Filter

- Filter: keep in a list only the values that satisfy some logical condition/boolean function:

```
- fun filter(f,l) =  
    if l=[] then []  
    else if f(hd l)  
        then (hd l)::(filter (f, tl l))  
        else filter(f, tl l);  
val filter = fn : ('a -> bool) * 'a list -> 'a list  
  
- filter((fn x => x>0), [~1,0,1,2,3,~2,4]);  
val it = [1,2,3,4] : int list
```

Permutations

```
- fun myInterleave(x, []) = [[x]]
  | myInterleave(x, h::t) =
      (x::h::t)::(
          map((fn l => h::l), myInterleave(x, t)));

- myInterleave(1, []);
val it = [[1]] : int list list

- myInterleave(1, [3]);
val it = [[1,3],[3,1]] : int list list

- myInterleave(1, [2,3]);
val it = [[1,2,3],[2,1,3],[2,3,1]] : int list list
```

Permutations

```
- fun appendAll(nil) = nil  
| appendAll(z::zs) = z @ (appendAll(zs));
```

flattens the list

```
- appendAll([[1,2],[2,1]]);  
val it = [[1,2],[2,1]] : int list list
```

```
- fun permute(nil) = [[]]  
| permute(h::t) = appendAll(  
    map((fn l => myInterleave(h,l)), permute(t)));
```

```
- permute([1,2,3]);  
val it = [[1,2,3],[2,1,3],[2,3,1],[1,3,2],[3,1,2],  
          [3,2,1]] : int list list
```

Currying = partial application

- fun f a b c = a+b+c;

OR

- fun f(a) (b) (c) = a+b+c;

val f = fn : int -> int -> int -> int

val f = fn : int -> (int -> (int -> int))

- val inc1 = f(1);

val inc1 = fn : int -> int -> int

val inc1 = fn : int -> (int -> int)

- val inc12 = inc1(2);

val inc12 = fn : int -> int

- inc12(3);

val it = 6 : int

Currying and *Lazy evaluation*

```
- fun mult x y = if x = 0 then 0 else x * y;
```

Eager evaluation: reduce as much as possible before applying the function

```
mult (1-1) (3 div 0)
```

```
-> (fn x => (fn y => if x = 0 then 0 else x * y)) (1-1) (3 div 0)
```

```
-> (fn x => (fn y => if x = 0 then 0 else x * y)) 0 (3 div 0)
```

```
-> (fn y => if 0 = 0 then 0 else 0 * y) (3 div 0)
```

```
-> (fn y => if 0 = 0 then 0 else 0 * y) error
```

```
-> error
```

Lazy evaluation:

```
mult (1-1) (3 div 0)
```

```
-> (fn x => (fn y => if x = 0 then 0 else x * y)) (1-1) (3 div 0)
```

```
-> (fn y => if (1-1) = 0 then 0 else (1-1) * y) (3 div 0)
```

```
-> if (1-1) = 0 then 0 else (1-1) * (3 div 0)
```

```
-> if 0 = 0 then 0 else (1-1) * (3 div 0)
```

```
-> 0
```

Currying and *Lazy evaluation*

- Argument evaluation as late as possible (possibly never)
- Evaluation only when indispensable for a reduction
- Each argument is evaluated at most once
- Lazy evaluation in Standard ML for the primitives: **if then else** , **andalso** , **orelse** , and pattern matching
- Property: If the eager evaluation of expression **e** gives **n1** and the lazy evaluation of **e** gives **n2** then **n1 = n2**
- Lazy evaluation gives a result more often

Sum sequence

```
- fun sum f n =  
    if n = 0 then 0  
    else f(n) + sum f (n-1);  
val sum = fn : (int → int) → int → int
```

```
- sum (fn x => x * x) 3 ;  
val it = 14 : int  
    because
```

$$f(3) + f(2) + f(1) + f(0) = 9 + 4 + 1 + 0 = 14$$

Composition

- Composition is another example of a higher-order function:

```
- fun comp (f, g) (x) = f (g (x)) ;
```

```
val comp = fn : ('a -> 'b) * ('c -> 'a) -> 'c -> 'b
```

```
- val f = comp(Math.sin, Math.cos) ;
```

```
val f = fn : real -> real
```

SAME WITH:

```
- val g = Math.sin o Math.cos ;
```

(* Composition "o" is predefined *)

```
val g = fn : real -> real
```

```
- f(0.25) ;
```

```
val it = 0.824270418114 : real
```

```
- g(0.25) ;
```

```
val it = 0.824270418114 : real
```

Find

- Pick only the first element of a list that satisfies a given predicate:

```
- fun myFind pred nil = raise Fail "No such element"
```

```
  | myFind pred (x::xs) =
```

```
    if pred x then x
```

```
    else myFind pred xs;
```

```
val myFind = fn : ('a -> bool) -> 'a list -> 'a
```

```
- myFind (fn x => x > 0.0) [~1.2, ~3.4, 5.6, 7.8];
```

```
val it = 5.6 : real
```

Reduce (aka. foldr)

- We can generalize the notion of recursion over lists as follows: all recursions have a base case, an iterative case, and a way of combining results:

```
- fun reduce f b nil = b
  | reduce f b (x::xs) = f(x, reduce f b xs);
```

```
- fun sumList aList = reduce (op +) 0 aList;
val sumList = fn : int list -> int
```

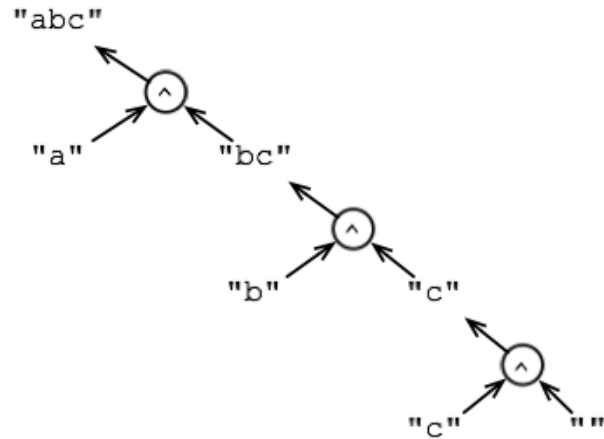
```
- sumList [1, 2, 3];
val it = 6 : int
```

foldl

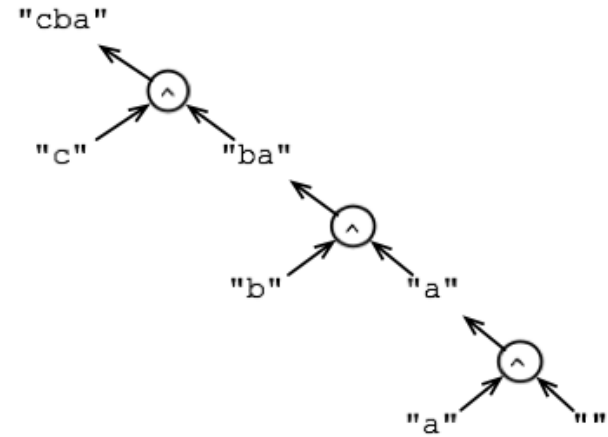
- `fun foldl(f: 'a*'b->'b, acc: 'b, l: 'a list): 'b =
 if l=[] then acc
 else foldl(f, f(hd(l),acc), tl(l));`
- `fun sum(l:int list):int =
 foldl((fn (x,acc) => acc+x),0,l);`
- `sum[1, 2, 3];`
`val it = 6 : int`
- it walks the list from left to right

foldl vs. reduce (foldr)

```
foldr (op ^) "" ["a", "b", "c"]  
type: ('a * 'b -> 'b) -> 'b -> 'a list -> 'b
```

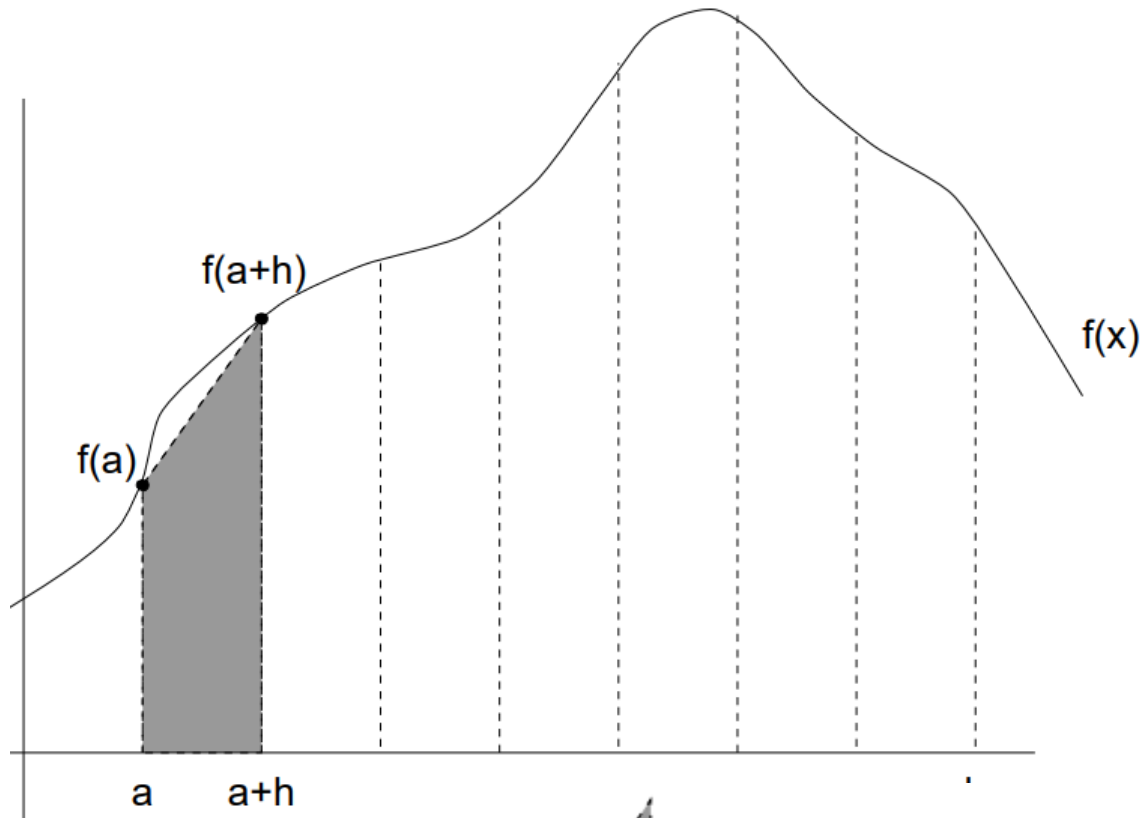


```
foldl (op ^) "" ["a", "b", "c"]  
type: ('a * 'b -> 'b) -> 'b -> 'a list -> 'b
```




Numerical integration

- Computation of $\int_a^b f(x) dx$ by the trapezoidal rule:



n intervals

$$h = (b - a) / n$$


$$= h * (f(a) + f(a+h)) / 2$$

Numerical integration

```
- fun integrate (f,a,b,n) =  
    if n <= 0 orelse b <= a then 0.0  
    else ((b-a) / real n) * ( f(a) + f(a+h) ) / 2.0 +  
        integrate (f,a+((b-a) / real n),b,n-1);  
val integrate = fn : (real → real) * real * real * int  
    → real
```

```
- fun cube x:real = x * x * x ;  
val cube = fn : real -> real  
- integrate ( cube , 0.0 , 2.0 , 10 ) ;  
val it = 4.04 : real
```


Collect like in Java streams

```
- fun collect(b, combine, accept, nil) = accept(b)
  | collect(b, combine, accept, x::xs) =
    collect(combine(b,x), combine, accept, xs);

- fun average(aList) = collect((0,0),
  (fn ((total,count),x) => (total+x,count+1)),
  (fn (total,count) => real(total)/real(count)),
  aList);

- average [1, 2, 4];
val it = 2.3333333333333 : real
```

Mutually recursive function definitions

```
- fun odd(n) = if n=0 then false
                else even(n-1)
```

and

```
    even(n) = if n=0 then true
                else odd(n-1);
```

```
val odd = fn : int -> bool
val even = fn : int -> bool
```

```
- even(1);
```

```
val it = false : bool
```

```
- odd(1);
```

```
val it = true : bool
```

Sorting

- *Merge-Sort:*

- To sort a list L:

- first split L into two disjoint sublists (of about equal size),
- then (recursively) sort the sublists, and
- finally merge the (now sorted) sublists

- It requires suitable functions for

- splitting a list into two sublists AND
- merging two sorted lists into one sorted list

Splitting

- We split a list by applying two functions, **take** and **skip**, which extract alternate elements; respectively, the elements at odd-numbered positions and the elements at even-numbered positions
- The definitions of the two functions mutually depend on each other, and hence provide an example of mutual recursion, as indicated by the SML-keyword **and**:

```
- fun take(L) =  
    if L = nil then nil  
    else hd(L)::skip(tl(L))  
  
and  
  
    skip(L) =  
    if L=nil then nil  
    else take(tl(L));  
  
val take = fn : 'a list -> 'a list  
val skip = fn : 'a list -> 'a list  
  
- take[1,2,3,4,5,6,7];  
val it = [1,3,5,7] : int list  
  
- skip[1,2,3,4,5,6,7];  
val it = [2,4,6] : int list
```

Merging

- Merge pattern definition:

```
- fun merge([],M) = M
```

```
| merge(L,[]) = L
```

```
| merge(x::xl,y::yl) =
```

```
    if (x:int)<y then x::merge(xl,y::yl)
```

```
    else y::merge(x::xl,yl);
```

```
val merge = fn : int list * int list -> int list
```

```
- merge([1,5,7,9],[2,3,6,8,10]);
```

```
val it = [1,2,3,5,6,7,8,9,10] : int list
```

```
- merge([],[1,2]);
```

```
val it = [1,2] : int list
```

```
- merge([1,2],[]);
```

```
val it = [1,2] : int list
```

Merge Sort

```
- fun sort(L) =  
  if L=[] then []  
  else if tl(L)=[] then L  
  else merge(sort(take(L)), sort(skip(L)));  
  
val sort = fn : int list -> int list
```

Local declarations

```
- fun fraction (n,d) =  
  let val k = gcd (n,d)  
  in  
    ( n div k , d div k )  
  end;
```

- The identifier **k** is local to the expression after **in**
- Its binding exists only during the evaluation of this expression
- All other declarations of **k** are hidden during the evaluation of this expression

Sorting with comparison

- How to sort a list of elements of type α ?
 - We need the **comparison function/operator** for elements of type α !

```
- fun sort order [ ] = [ ]  
  | sort order [x] = [x]  
  | sort order xs =  
    let fun merge [ ] M = M  
        | merge L [ ] = L  
        | merge (L as x::xs) (M as y::ys) =  
            if order(x,y) then x::merge xs M  
            else y::merge L ys  
    in merge (sort order ys) (sort order xs) end;  
- sort (op >) [5.1, 3.4, 7.4, 0.3, 4.0] ;  
val it = [7.4,5.1,4.0,3.4,0.3] : real list
```


Sorting with comparison

```
- fun split_helper(L: 'a list, Acc:'a list * 'a list)
    : 'a list * 'a list =
  if L=[] then Acc
  else split_helper(tl(L), (#2(Acc), (hd(L)) :: #1(Acc)));

- fun split(L) = split_helper(L, ([], []));
- split([1,2,3,4,5,6]);
  split([1,2,3,4,5,6])
  split_helper([1,2,3,4,5,6], ([], []))
  split_helper([2,3,4,5,6], ([], [1]))
  split_helper([3,4,5,6], ([1], [2]))
  split_helper([4,5,6], ([2], [3,1]))
  split_helper([5,6], ([3,1], [4,2]))
  split_helper([6], ([4,2], [5,3,1]))
  split_helper([], ([5,3,1], [6,4,2]))
  ([5,3,1], [6,4,2])
```

Quicksort

- C.A.R. Hoare, in 1962: Average-case running time: $\Theta(n \log n)$

```
- fun sort [ ] = [ ]  
| sort (x::xs) =  
  let val (S,B) = partition (x,xs)  
      in (sort S) @ (x :: (sort B))  
  end;
```

Double recursion and no tail-recursion

```
- fun partition (p,[ ]) = ([ ],[ ])  
| partition (p,x::xs) =  
  let val (S,B) = partition (p,xs)  
      in if x < p then (x::S,B) else (S,x::B)  
  end
```

Nested recursion

For $m, n \geq 0$:

`acker(0, m) = m+1`

`acker(n, 0) = acker(n-1, 1) for n > 0`

`acker(n, m) = acker(n-1, acker(n, m-1)) for n, m > 0`

- `fun acker 0 m = m+1`

| `acker n 0 = acker (n-1) 1`

| `acker n m = acker (n-1) (acker n (m-1));`

It is guaranteed to end because of *lexicographic order*:

`(n', m') < (n, m) iff n' < n or (n'=n and m' < m)`

Nested recursion

- *Knuth's up-arrow operator* \uparrow^n (invented by Donald Knuth):

$$a \uparrow^1 b = a^b$$

$$a \uparrow^n b = a \uparrow^{n-1} (b \uparrow^{n-1} a) \text{ for } n > 1$$

```
- fun opKnuth 1 a b = Math.pow (a,b)
  | opKnuth n a b = opKnuth (n-1) a
                    (opKnuth (n-1) b b);
```

```
- opKnuth 2 3.0 3.0 ;
```

```
val it = 7.62559748499E12 : real
```

```
- opKnuth 3 3.0 3.0 ;
```

```
! Uncaught exception: Overflow;
```

- *Graham's number* (also called the “largest” number):

```
- opKnuth 63 3.0 3.0,
```

Recursion on a generalized problem

- It is impossible to determine whether n is prime via the reply to the question “is $n - 1$ prime”?
 - It seems impossible to directly construct a recursive program
 - We thus need to find another function that is more general than prime, in the sense that prime is a particular case of this function
 - for which a recursive program can be constructed

```
- fun ndivisors n low up = low > up orelse  
(n mod low) <> 0 andalso ndivisors n (low+1) up;
```

```
- fun prime n = if n <= 0  
  then error "prime: non-positive argument"  
  else if n = 1 then false  
  else ndivisors n 2 floor(Math.sqrt(real n));
```

- The discovery of divisors requires imagination and creativity

Tail recursion

```
- fun length [ ] = 0
| length (x::xs) = 1 + length xs;
```

- The recursive call of **length** is nested in an expression: during the evaluation, all the terms of the sum are **stored**, hence the memory consumption for expressions & bindings is proportional to the length of the list!

```
length [5,8,4,3]
-> 1 + length [8,4,3]
-> 1 + (1 + length [4,3])
-> 1 + (1 + (1 + length [3]))
-> 1 + (1 + (1 + (1 + length [ ])))
-----
-> 1 + (1 + (1 + (1 + 0)))
-> 1 + (1 + (1 + 1))
-> 1 + (1 + 2)
-> 1 + 3
-> 4
```

Tail recursion

```
- fun lengthAux [ ] acc = acc
| lengthAux (x::xs) acc = lengthAux xs (acc+1);
- fun length L = lengthAux L 0;
- length [5,8,4,3];
  -> lengthAux [5,8,4,3] 0
  -> lengthAux [8,4,3] (0+1)
  -> lengthAux [8,4,3] 1
  -> lengthAux [4,3] (1+1)
  -> lengthAux [4,3] 2
  -> lengthAux [3] (2+1)
  -> lengthAux [3] 3
  -> lengthAux [ ] (3+1)
  -> lengthAux [ ] 4
  -> 4
```

- **Tail recursion:** recursion is the outermost operation
 - **Space complexity:** constant memory consumption for expressions & bindings (SML can use the **same stack frame/activation record**)
 - Time complexity: (still) one traversal of the list

Tail recursion

```
- fun factAux 0 acc = acc
  | factAux n acc = factAux (n-1) (n*acc);
- fun fact n =
  if n < 0 then error "fact: negative argument"
  else factAux n 1;

- fact (3) ;
-> factAux (3, 1)
-> factAux (2, 3)
-> factAux (1, 6)
-> factAux (0, 6)
6
```


Records

- Records are structured data types of heterogeneous elements that are labeled

- `{x=2, y=3};`

- The order does not matter:

- `{make="Toyota", model="Corolla", year=2017, color="silver"}`

- = `{model="Corolla", make="Toyota", color="silver", year=2017};`

```
val it = true : bool
```

- `fun full_name{first:string, last:string, age:int, balance:real}:string =`

- `first ^ " " ^ last;`

- `(* ^ is the string concatenation operator *)`

```
val full_name=fn:{age:int, balance:real, first:string, last:string} -> string
```

string and char

```
- "a";
```

```
val it = "a" : string
```

```
- #"a";
```

```
val it = #"a" : char
```

```
- explode("ab");
```

```
val it = [#"a",#"b"] : char list
```

```
- implode([#"a",#"b"]);
```

```
val it = "ab" : string
```

```
- "abc" ^ "def" = "abcdef";
```

```
val it = true : bool
```

```
- size ("abcd");
```

```
val it = 4 : int
```

string and char

```
- String.sub("abcde",2);  
val it = #"c" : char  
  
- substring("abcdefghij",3,4);  
val it = "defg" : string  
  
- concat ["AB"," ","CD"];  
val it = "AB CD" : string  
  
- str("#x");  
val it = "x" : string
```

Functional programming in SML

- Covered fundamental elements:
 - Evaluation by reduction of expressions
 - Recursion
 - Polymorphism via type variables
 - Strong typing
 - Type inference
 - Pattern matching
 - Higher-order functions
 - Tail recursion

Beyond functional programming

- *Relational programming* (aka *logic programming*)
 - For which triples does the **append** relation hold?
 - ?- **append** ([1,2], [3], **X**).
 - Yes**
 - X** = [1,2,3]
 - ?- **append** ([1,2], **X**, [1,2,3]).
 - X** = [3]
 - ?- **append** (**X**, **Y**, [1,2,3]).
 - X** = [], **Y** = [1,2,3];
 - X** = [1], **Y** = [2,3];
 - ...
 - X** = [1,2,3], **Y** = [];
 - **No differentiation between arguments and results!**

Beyond functional programming

- *Backtracking* mechanism to enumerate all the possibilities
- *Unification* mechanism, as a generalization of pattern matching
- Power of the logic paradigm / relational framework

Beyond functional programming

- *Constraint Processing:*

- Constraint Satisfaction Problems (CSPs)

- Variables: X_1, X_2, \dots, X_n

- Domains of the variables: D_1, D_2, \dots, D_n

- Constraints on the variables: examples: $3 \cdot X_1 + 4 \cdot X_2 \leq X_4$

- What is a solution?

- An assignment to each variable of a value from its domain, such that all the constraints are **satisfied**

- **Objectives:**

- Find a solution

- Find all the solutions

- Find an optimal solution, according to some cost expression on the variables

Beyond functional programming

- The n-Queens Problem:

- How to place n queens on an $n \times n$ chessboard such that no queen is threatened?
- Variables: X_1, X_2, \dots, X_n (one variable for each column)
- Domains of the variables: $D_i = \{1, 2, \dots, n\}$ (the rows)
- Constraints on the variables:
 - No two queens are in the same column: this is impossible by the choice of the variables!
 - No two queens are in the same row: $X_i \neq X_j$, for each $i \neq j$
 - No two queens are in the same diagonal: $|X_i - X_j| \neq |i - j|$, for each $i \neq j$
 - Number of candidate solutions: n^n

- Exhaustive Enumeration

- **Generation** of possible values of the variables.
- **Test** of the constraints.

- Optimization:

- Where to place a queen in column k such that it is compatible with r_{k+1}, \dots, r_n ?
- Eliminate possible locations as we place queens

Beyond functional programming

- Applications:
 - Scheduling
 - Planning
 - Transport
 - Logistics
 - Games
 - Puzzles
- Complexity
 - Generally these problems are NP-complete with exponential complexity

The program of Young McML

```
fun tartan_column(i,j,n) =  
  if j=n+1 then "\n"  
  else if (i+j) mod 2=1 then  
    concat(["* ",tartan_column(i,j+1,n)])  
  else concat(["+ ",tartan_column(i,j+1,n)]);  
fun tartan_row(i,n) =  
  if i=n+1 then ""  
  else concat([tartan_column(i,1,n),  
    tartan_row(i+1,n)]);  
fun tartan(n) = tartan_row(1,n);  
print(tartan(30));
```