### **Tabled Resolution**

CSE 595 – Semantic Web

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http://www3.cs.stonybrook.edu/~pfodor/courses/cse595.html

### Recap: OLD Resolution

- Prolog follows *OLD resolution* = SLD with left-to-right literal selection.
- Prolog searches for OLD proofs by expanding the resolution tree depth first.
  - This depth-first expansion is close to how procedural programs are evaluated:
    - Consider a goal  $G_1$ ,  $G_2$ ,...,  $G_n$  as a "procedure stack" with  $G_1$ , the selected literal on top.
    - Call **G**<sub>1</sub>.
    - If and when  $G_1$  returns, continue with the rest of the computation: call  $G_2$ , and upon its return call  $G_3$ , etc. until nothing is left
      - Note:  $\mathbf{G_2}$  is "opened up" only when  $\mathbf{G_1}$  returns, not after executing only some part of  $\mathbf{G_1}$ .

### **OLD** Resolution

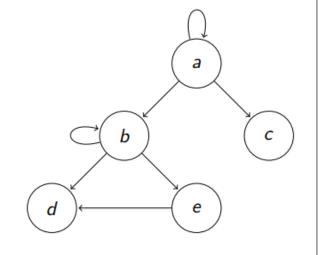
• Depth-first expansion, however, contributes to the *incompleteness* of Prolog's evaluation, which may **not terminate** even when the least model is finite (see the next example!)

### Example:Reachability in Directed Graphs

- Determining <u>whether there is</u> a <u>path between two vertices</u> in a directed graph is an important and widespread problem
- For instance, consider checking whether (or not) a program accesses a shared resource before obtaining a lock:
  - A program itself can be considered as a graph with vertices representing program states!
    - A state may be characterized by the program counter value, and values of variables.
    - There are richer models for representing program evaluation, but a directed graph is most basic.
  - If we can go from state **s** by executing one <u>instruction</u> to **s'**, then we can place an edge from **s** to **s'**
  - The <u>reachability question</u> may be whether we <u>can reach from the start</u> state to a state accessing a shared resource, without going through a state that obtained a lock.

### Graph Reachability as a Logic Program

- A finite directed graph can be represented by a set of binary facts representing an "<u>edge</u>" relation
  - Predicate "q" on right is an example
- Reachability can then be written as a "transitive closure" over the edge relation
  - Observe the predicate "**r**" defined on right using two clauses:
    - The first clause: there is a *path* from **X** to **Y** if there is an edge from **X** to **Y**
    - The second clause: there is a *path* from **X** to **Y** if there an intermediate vertex **Z** such that:
      - there is an edge from **X** to **Z**, and
      - there is a path from **Z** to **Y**



```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

r(X,Y) := q(X,Y).

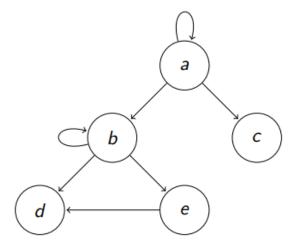
q(X,Z), r(Z,Y).

r(X,Y) :-

#### **Bottom-Up Evaluation**

- Note that the program on the left is a Datalog program: no function symbols
  - Its Herbrand Universe is **finite**, and its least model computation using the bottom up evaluation will **terminate**:

```
M_0 = \emptyset
M_1 = T_p(M_0) = M_0 \cup \{q(a,a),
 q(a,b), q(a,c), q(b,b),
 q(b,d), q(b,e), q(e,d)}
M_2 = T_D(M_1) = M_1 \cup \{r(a,a),
 r(a,b), r(a,c), r(b,b),
 r(b,d), r(b,e), r(e,d)}
M_3 = T_p(M_2) = M_2 \cup \{r(a,d),
r(a,e)}
M_4 = T_D(M_3) = M_3
                (c) Paul Fodor (CS Stony Brook) and Elsevier
```



```
q(a, a).

q(a, b).

q(a, c).

q(b, b).

q(b, d).

q(b, e).

q(e, d).

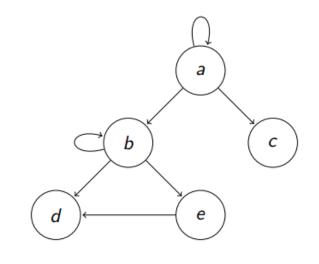
r(X,Y) :- q(X,Y).
```

q(X,Z), r(Z,Y).

r(X,Y) :-

### Bottom-Up Evaluation

```
M<sub>4</sub> = {q(a,a), q(a,b), q(a,c),
q(b,b), q(b,d), q(b,e),
q(e,d), r(a,a), r(a,b),
r(a,c), r(b,b), r(b,d),
r(b,e), r(e,d), r(a,d),
r(a,e)}
```



With care, using bottom-up evaluation allpairs reachability can be computed in
O(V·E) time for a graph with V vertices and E edges

```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
r(X,Y) := q(X,Y).

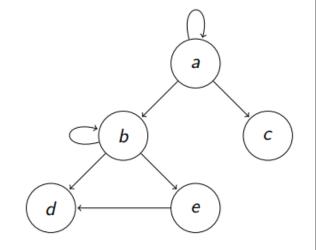
r(X,Y) :=

q(X,Z), r(Z,Y).
```

?- r(a,N).

Consider initial query "?- r(a,N)."

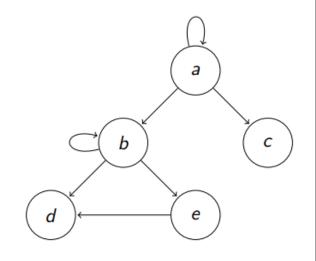
Let's construct the <u>SLD tree</u> for this query



```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```

Resolving this goal with the first clause of  $\mathbf{r}$ , we get a new goal "?-  $\mathbf{q}(\mathbf{a}, \mathbf{N})$ ."

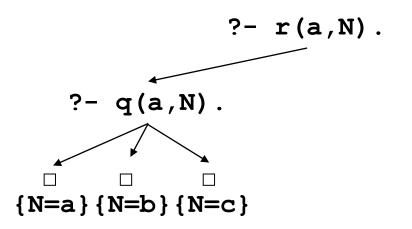


```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

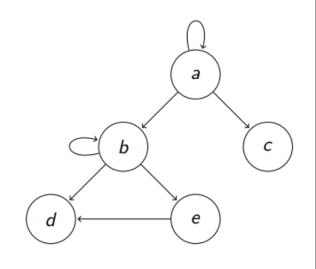
```
r(X,Y) := q(X,Y).

r(X,Y) :=

q(X,Z), r(Z,Y).
```



Resolving "?- q(a,N)." results in the empty goal, under three answer substitutions: a, b, and c. There are no more ways to resolve "?- q(a,N)."

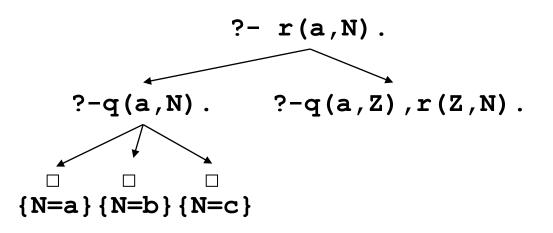


```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
r(X,Y) := q(X,Y).

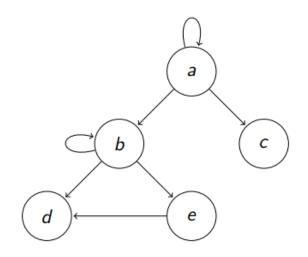
r(X,Y) :=

q(X,Z), r(Z,Y).
```



Resolving "?- r(a,N)." with the second clause of r results in goal
"?- q(a,Z),r(Z,N)."

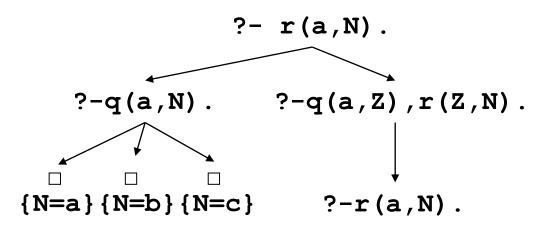
Note: We will use same variable names as in the program clause when possible, instead of mechanically inventing new variable names in every step.



```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).

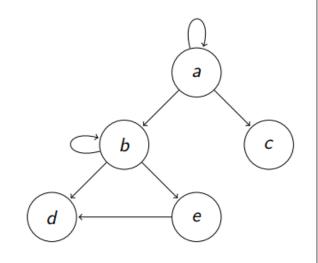
r(X,Y) :- q(X,Y).
r(X,Y) :-
```

q(X,Z), r(Z,Y).



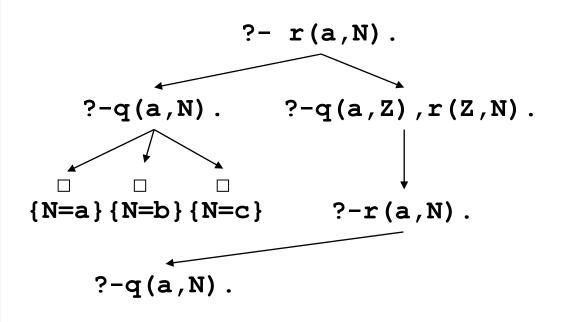
The selected literal is q(a,Z) which unifies with fact q(a,a) with Z=a. Thus we get the goal "?-r(a,N)."

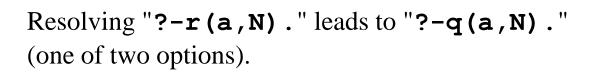
Note: This is the same goal that we had at the beginning.

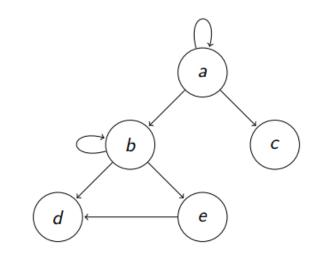


```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

$$r(X,Y) := q(X,Y).$$
  
 $r(X,Y) :=$   
 $q(X,Z), r(Z,Y).$ 





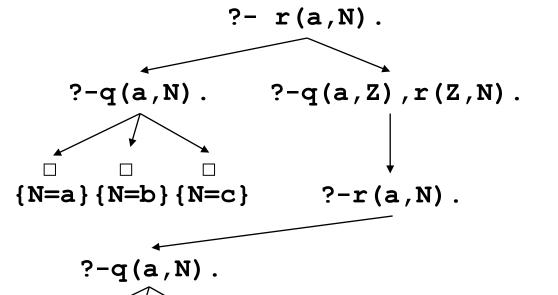


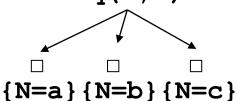
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
r(X,Y) := q(X,Y).

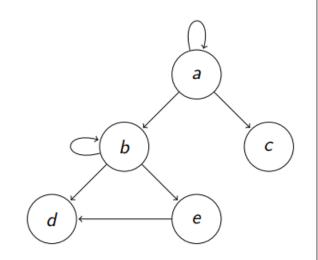
r(X,Y) :=

q(X,Z), r(Z,Y).
```





"?-q(a,N).", in turn, leads to empty goal with answers N=a, b, and c.

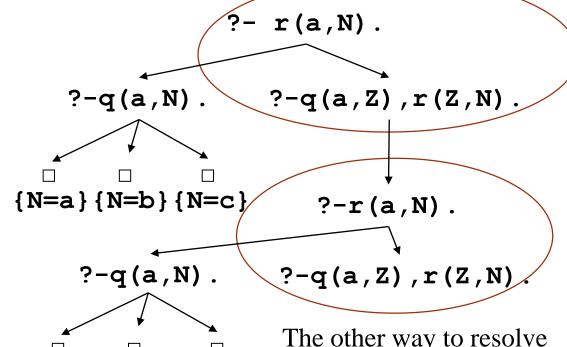


```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
r(X,Y) := q(X,Y).

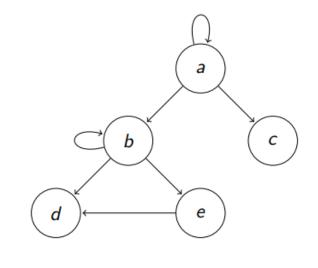
r(X,Y) :=

q(X,Z), r(Z,Y).
```



The other way to resolve "?-r(a,N)." is ...

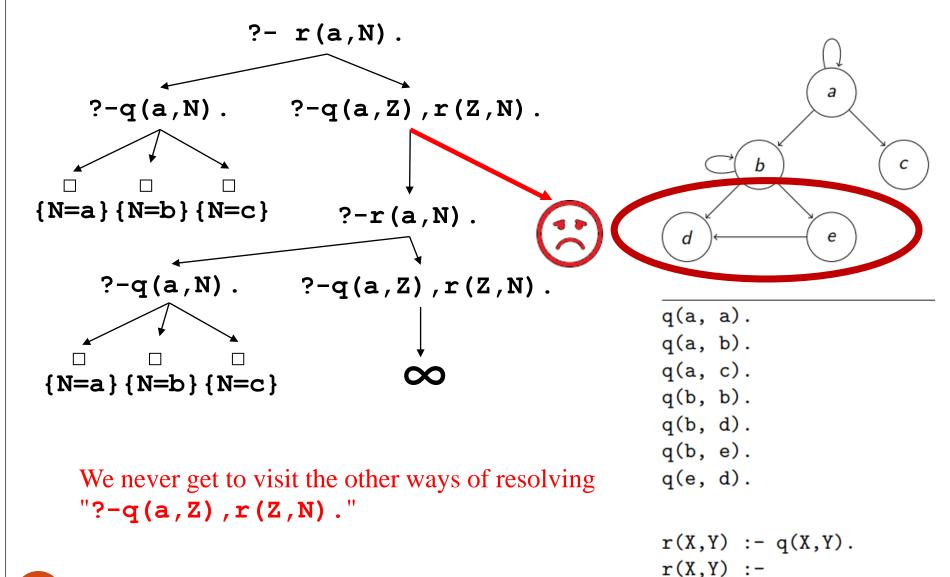
You get the drift: we are repeating work done before, and there is an infinite branch here.



```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

$$r(X,Y) := q(X,Y).$$
  
 $r(X,Y) :=$   
 $q(X,Z), r(Z,Y).$ 

 ${N=a}{N=b}{N=c}$ 



q(X,Z), r(Z,Y).

### Depth-First Expansion of the OLD tree

- If the underlying graph is acyclic, all branches in the OLD tree will be finite
- If the graph is cyclic, nothing to the right of an infinite branch is expanded.
- This renders the evaluation <u>incomplete</u>: goals for which there are OLD derivations, but they are not found.
- Moreover, the same answer may be returned multiple times (even <u>infinitely</u>!)
  - Even if the underlying graph is acyclic, this evaluation is not efficient
    - For query of the form **r**(**a**, **N**) we will return **N=b** for each path from "**a**" to "**b**".

### Depth-First Expansion of the OLD tree

- <u>Breadth-First</u> expansion does have the completeness property:
  - Every OLD derivation will be eventually constructed.
    - If something is a logical consequence, we will eventually confirm it in a finite number of levels
  - But we may not be able to conclude negative information
    - If something is not a logical consequence, we may never be able to identify it because we don't know when to stop?

### Depth-First Expansion of the OLD tree

- Moreover, Breadth-First expansion does not give a natural operational understanding:
  - If we view predicates as being defined by "procedures", then breadth first expansion steps through a procedure's evaluation, switching contexts at the end of each step.
  - As in procedural programming, <u>context switching</u> <u>is expensive</u> (in this case, we've to <u>switch</u> <u>substitutions</u>)

### Programming our way around the problem

• Is there a path from **X** to **Y** that does not visit any vertex already seen in **L**?

```
p(L, X, Y) :-
        q(X, Y),
        not member(Y, L).
p(L, X, Y) :-
        q(X, Z),
        not member(Z, L),
        p([Z|L], Z, Y).
```

• Now, start from **L=[]** to look for reachable vertices

```
r(X, Y) :-
p([], X, Y).
```

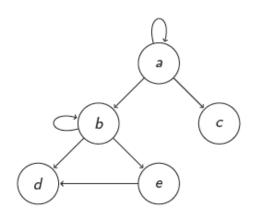
• In **L** we remember the path so far, and use this to avoid loops

### Programming our way around the problem

- We are assured termination for reachability queries
  - We stop if a node has been seen before on the same branch.
- Still, this is inefficient: may take exponential time
  - We re-execute queries on different branches of the SLD/OLD tree

### What is Tabled Resolution?

- Memoize calls and results to avoid repeated subcomputations.
  - Termination: Avoid performing computations that repeat infinitely often.
    - Complete for Datalog programs
  - Efficiency: Dynamically share common subexpressions.



```
?- r(a, N)
```

```
q(a, a).

q(a, b).

q(a, c).

q(b, b).

q(b, d).

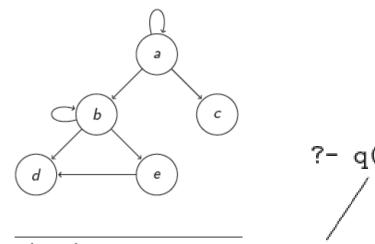
q(b, e).

q(e, d).

r(X,Y) :- q(X,Y).

r(X,Y) :-
```

q(X,Z), r(Z,Y).

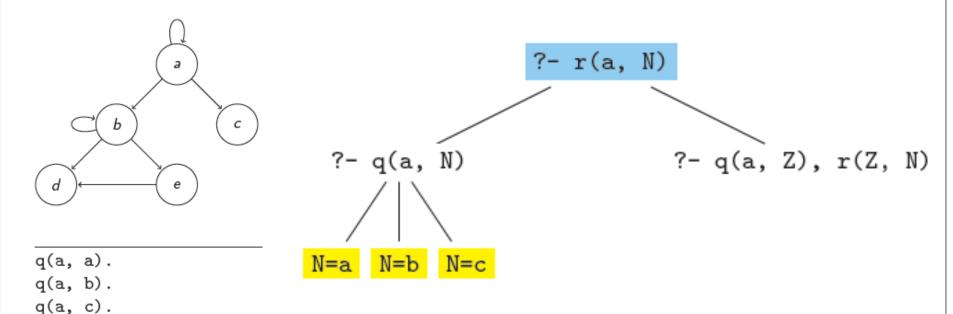


```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
r(X,Y) := q(X,Y).

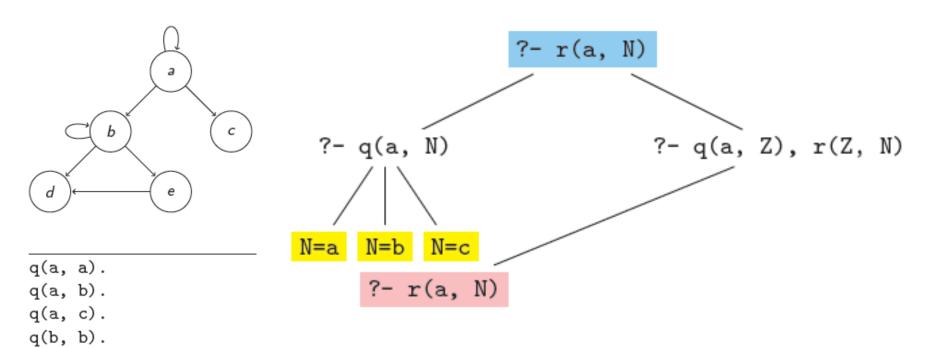
r(X,Y) := q(X,Z), r(Z,Y).
```

```
?- r(a, N)
?- q(a, N)
N=a N=b N=c
```



```
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```

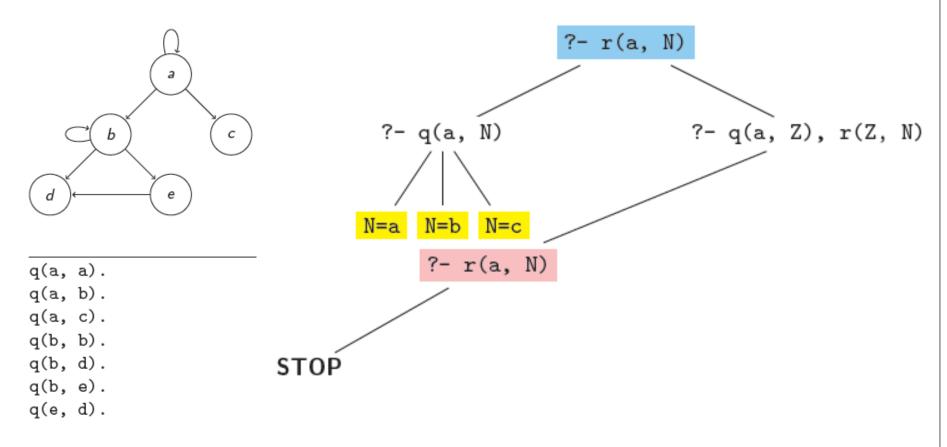
q(b, b). q(b, d). q(b, e). q(e, d).



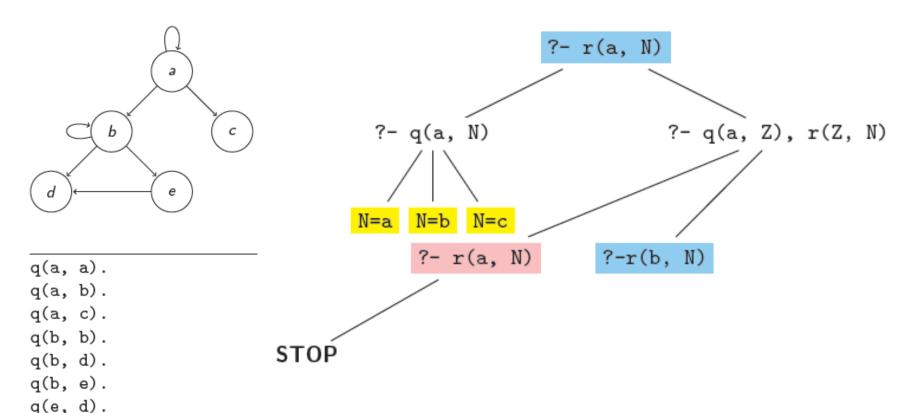
```
r(X,Y) := q(X,Y).

r(X,Y) := q(X,Z), r(Z,Y).
```

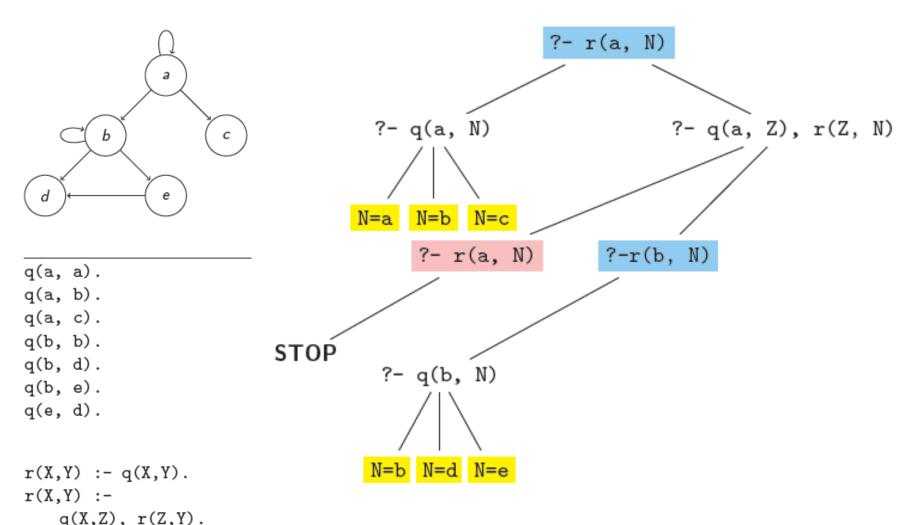
q(b, d). q(b, e). q(e, d).

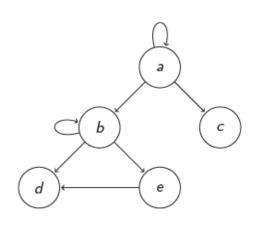


r(X,Y) := q(X,Y).r(X,Y) := q(X,Z), r(Z,Y).



```
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```

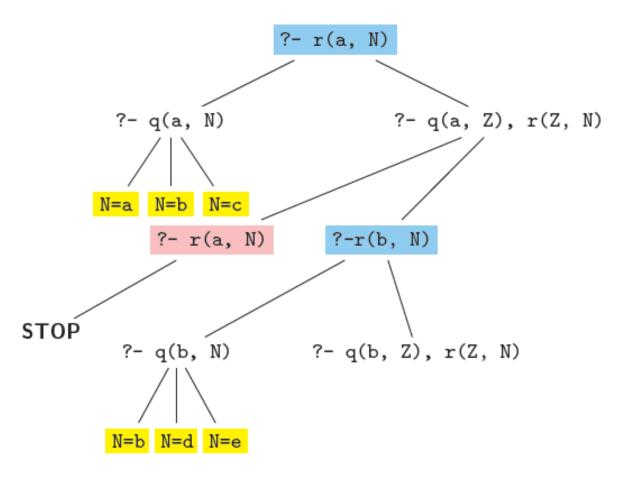


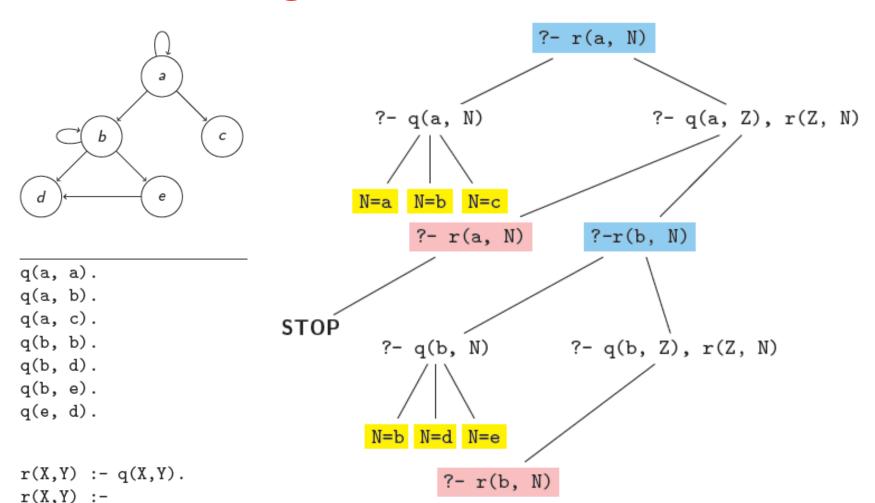


```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

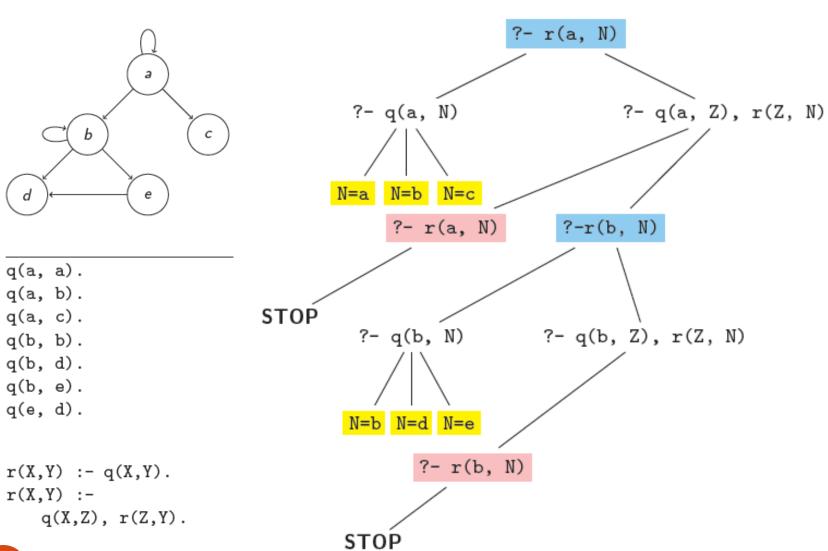
```
r(X,Y) := q(X,Y).

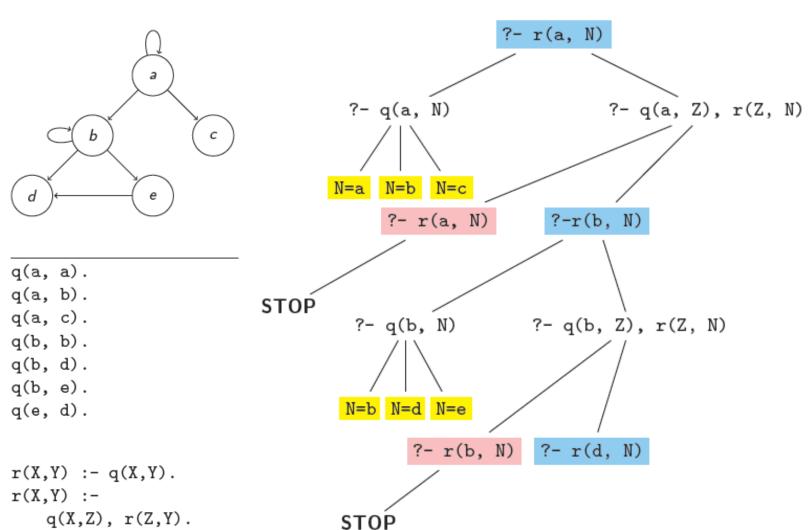
r(X,Y) := q(X,Z), r(Z,Y).
```

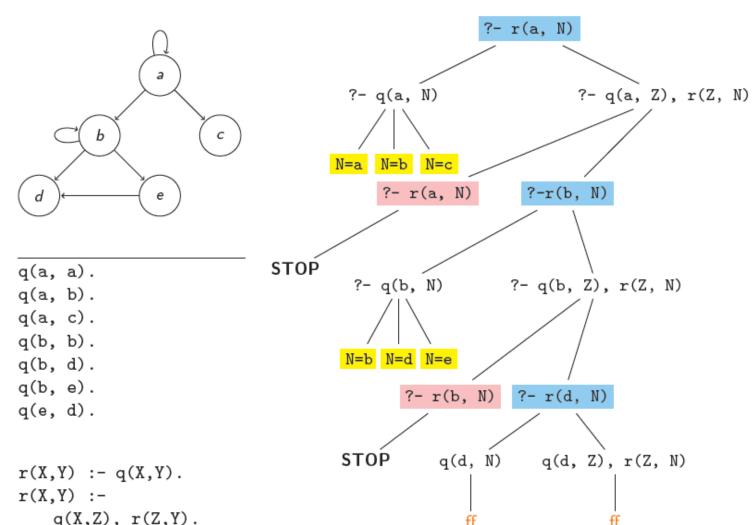


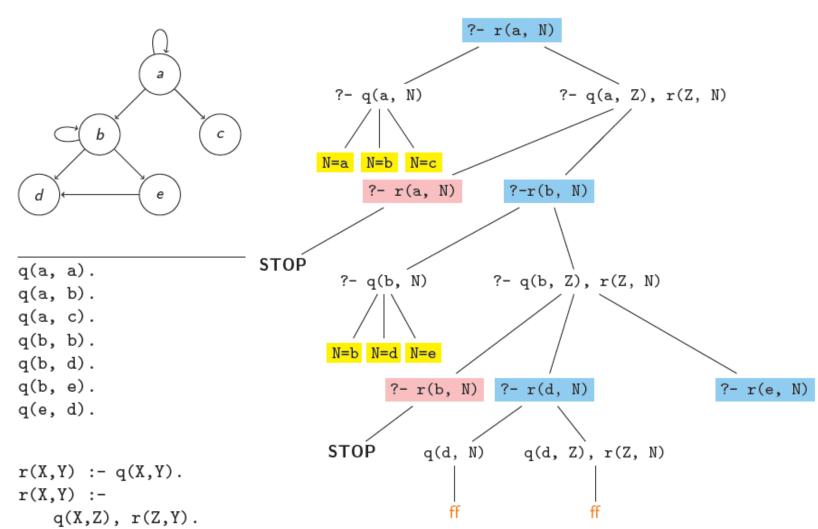


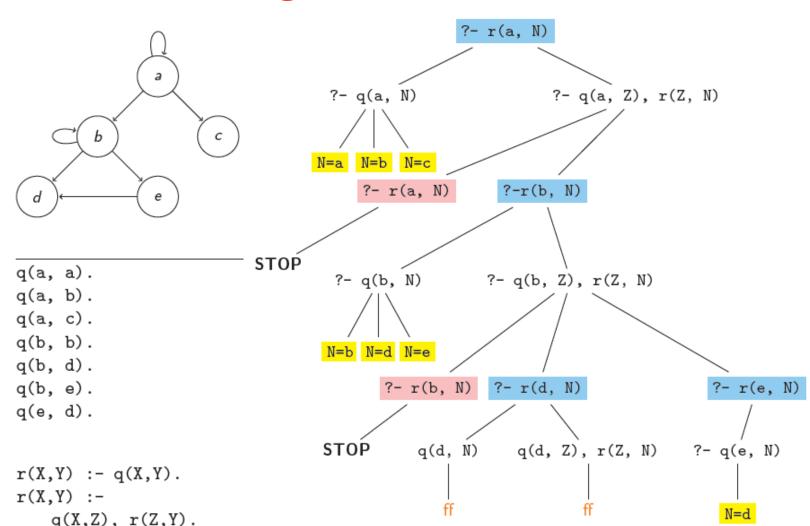
q(X,Z), r(Z,Y).

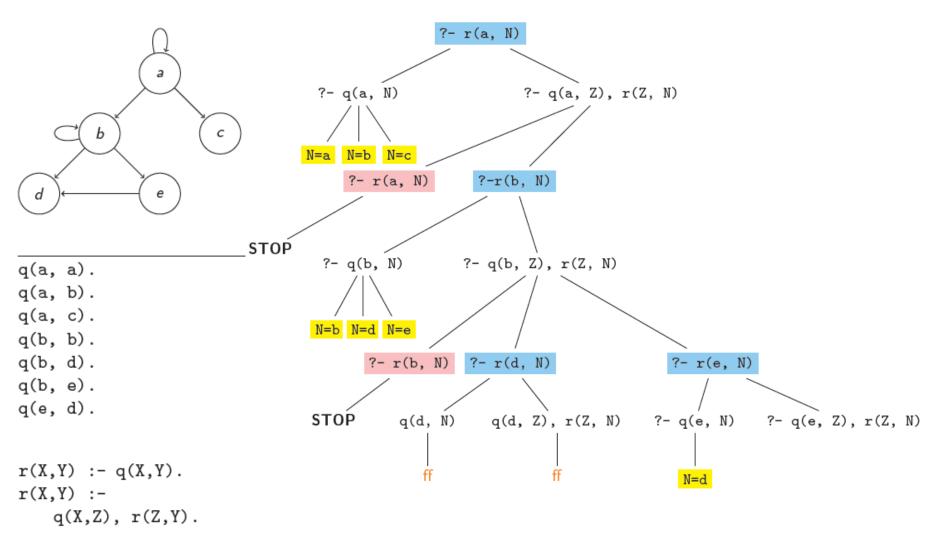


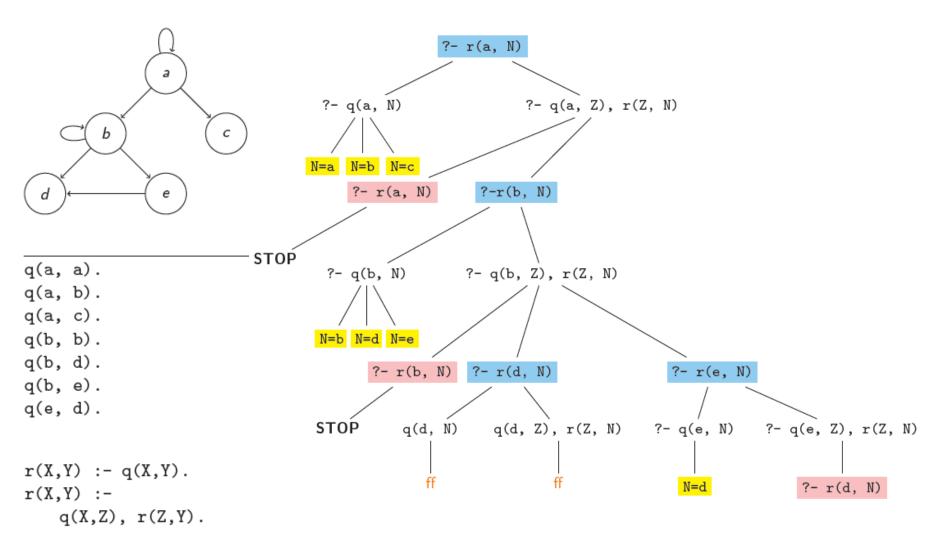


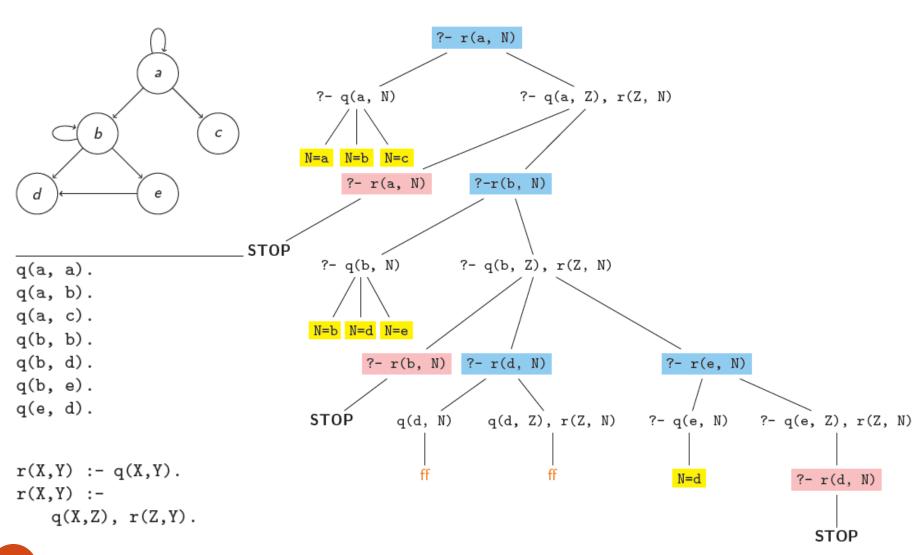


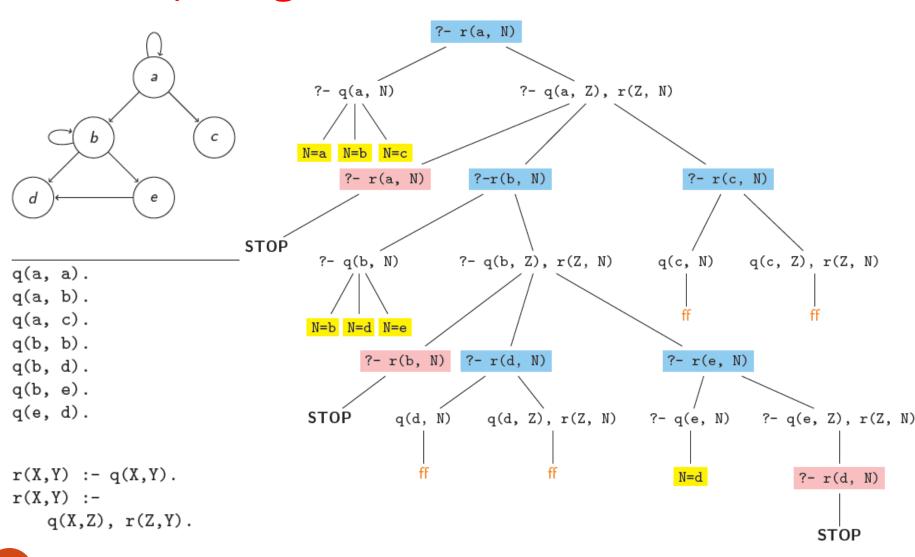






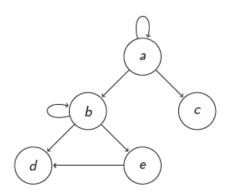






# Rationale for goal-based stopping

- The OLD tree is a representation of search for successful derivations
  - which are finite sequences of goals terminating in an empty goal.
- If there is a successful derivation, then there is an equivalent one that does not repeat the same goal (compare to reachability via loop-free paths in a graph).
- Hence ignoring paths with repeated goals is *sound*: the derivations pruned away by stopping have equivalent ones that will not be ignored.
- Unfortunately, this scheme still does not fix the problem of infinite derivations

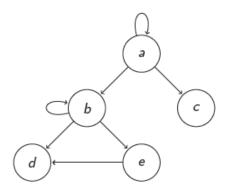


```
?- p(a, N)
```

```
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).

%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :-
p(X,Z), q(Z,Y).
```

q(a, a).

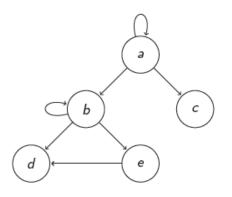


```
?- p(a, N)
?- q(a, N)
N=a N=b N=c
```

```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

Expand tree as usual

```
%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :-
p(X,Z), q(Z,Y).
```



```
q(a, a).
```

q(a, b).

q(a, c).

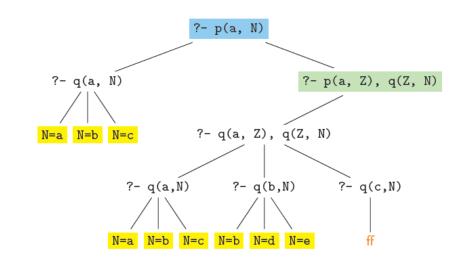
q(b, b).

q(b, d).

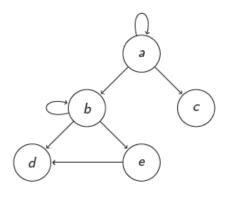
q(b, e).

q(e, d).

%Alternative formulation % of reachability: Note % use of LEFT recursion p(X,Y) :- q(X,Y). p(X,Y) :p(X,Z), q(Z,Y).



Expand tree as usual



```
q(a, a).
```

q(a, b).

q(a, c).

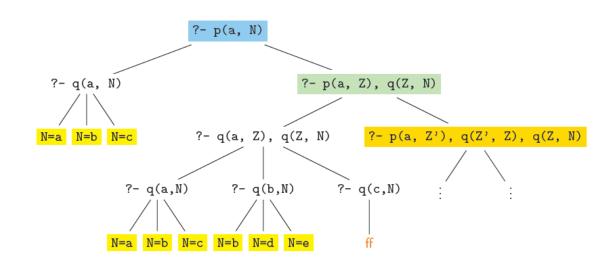
q(b, b).

q(b, d).

q(b, e).

q(e, d).

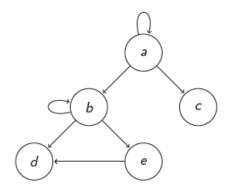
%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :p(X,Z), q(Z,Y).



Note that the right-most branch has ever-growing goals.

# OLD Resolution with Tabling (OLDT)

- The selected literal at a step in a derivation is known as a *call*.
- OLDT maintains a **table of calls** (initially empty).
- With each call, it maintains a table of computed answers (initially empty).
- Start resolution as in OLD.
- When a literal is selected, check the call table.
  - If the literal is in the table, <u>resolve it with its answers</u> in its answer table.
  - If the literal is not in the table, resolve with program clauses (as in OLD), and <u>add computed answers to its answer table</u>.



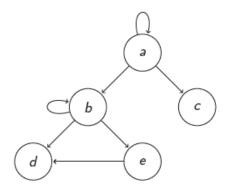
```
?- p(a, N)
```

```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :-
p(X,Z), q(Z,Y).
```

Calls	Answers

#### Start with empty tables



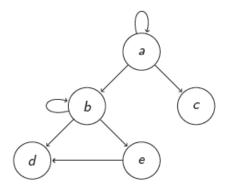
```
?- p(a, N)
```

```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :-
p(X,Z), q(Z,Y).
```

Calls	Answers

Pick selected literal. Is it in call table?



```
?- p(a, N)
```

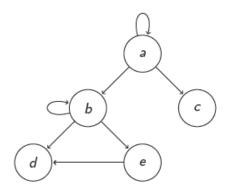
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :-
p(X,Z), q(Z,Y).
```

Calls	Answers
p(a, W)	

#### Add to call table

(c) Paul Fodor (CS Stony Brook) and Elsevier



q(a, a).

q(a, b).

q(a, c).

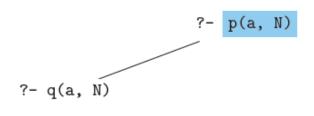
q(b, b).

q(b, d).

q(b, e).

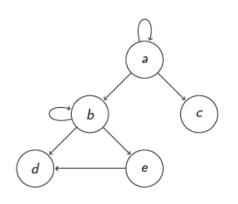
q(e, d).

%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :p(X,Z), q(Z,Y).



Calls	Answers
p(a, W)	

Do OLD resolution with program clauses



```
?- q(a, N)

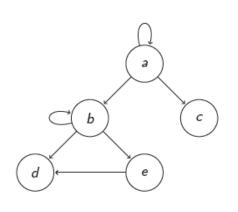
N=a
```

```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :-
p(X,Z), q(Z,Y).
```

Calls	Answers
p(a, W)	{p(a,a)

Add computed answer to table (if not already there)



```
?- q(a, N)

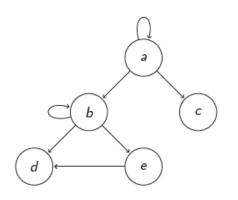
N=a N=b N=c
```

```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) := q(X,Y).
p(X,Y) :-
p(X,Z), $q(Z,Y)$ .

Calls	Answers
p(a, W)	{p(a,a), p(a,b), p(a,c)

Add computed answer to table (if not already there)



```
?- q(a, N)
?- q(a, N)
?- p(a, Z), q(Z, N)
N=a N=b N=c
```

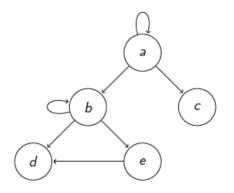
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
        Calls
        Answers

        p(a, W)
        {p(a,a), p(a,b), p(a,c)
```

%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :p(X,Z), q(Z,Y).

#### Continue with OLD resolution



```
q(a, a).
```

q(a, b).

q(a, c).

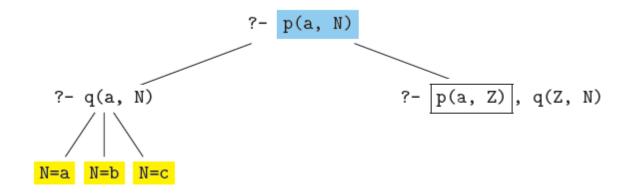
q(b, b).

q(b, d).

q(b, e).

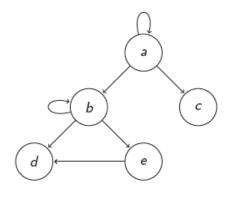
q(e, d).

%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :p(X,Z), q(Z,Y).



Calls	Answers
p(a, W)	{p(a,a), p(a,b), p(a,c)

Pick selected literal. Is it in call table?



```
q(a, a).
```

q(a, b).

q(a, c).

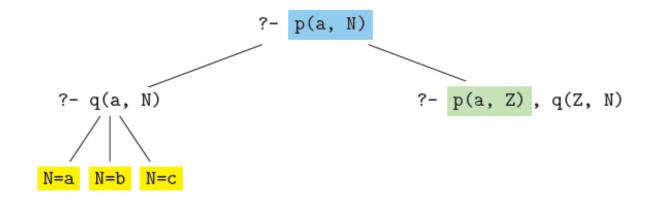
q(b, b).

q(b, d).

q(b, e).

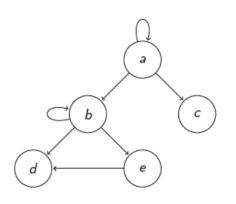
q(e, d).

```
%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :-
p(X,Z), q(Z,Y).
```



Calls	Answers
p(a, W)	{p(a,a), p(a,b), p(a,c)

Yes, resolve with answers in table



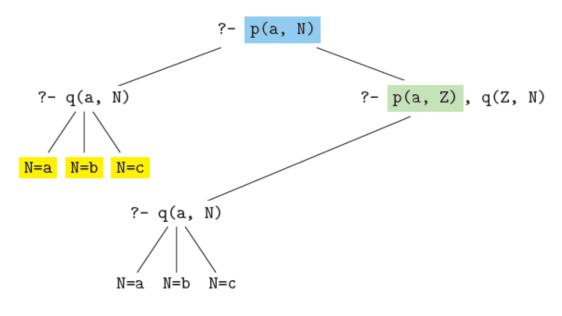
```
q(a, a).
q(a, b).
q(a, c).
```

q(b, b).

q(b, d). q(b, e).

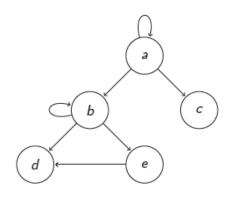
q(e, d).

%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :p(X,Z), q(Z,Y).



Calls	Answers
p(a, W)	{p(a,a), p(a,b), p(a,c)

Add computed answer to table (if not already there)



```
q(a, a).
```

q(a, b).

q(a, c).

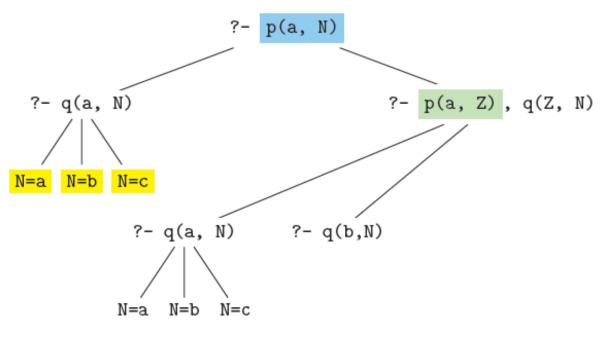
q(b, b).

q(b, d).

q(b, e).

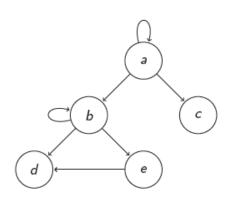
q(e, d).

%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :p(X,Z), q(Z,Y).



Calls	Answers
p(a, W)	{p(a,a), p(a,b), p(a,c)

Continue resolving with answers in table



```
q(a, a).
```

q(a, b).

q(a, c).

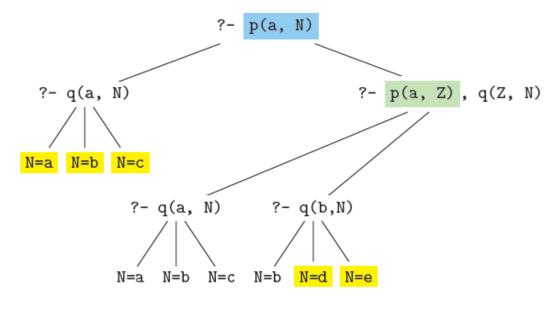
q(b, b).

q(b, d).

q(b, e).

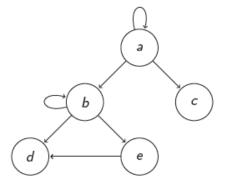
q(e, d).

%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :p(X,Z), q(Z,Y).



Calls	Answers
p(a, W)	${p(a,a), p(a,b), p(a,c)}$
	p(a,d), p(a,e)

Add computed answer to table (if not already there)



```
q(a, a).
```

q(a, b).

q(a, c).

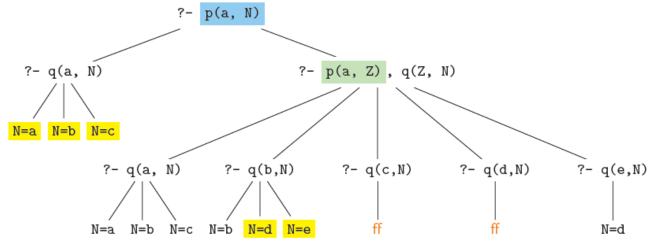
q(b, b).

q(b, d).

q(b, e).

q(e, d).

%Alternative formulation
% of reachability: Note
% use of LEFT recursion
p(X,Y) :- q(X,Y).
p(X,Y) :p(X,Z), q(Z,Y).



Calls	Answers
p(a, W)	${p(a,a), p(a,b), p(a,c)}$
	p(a,d), p(a,e)

Complete when all answers have been considered for resolution

#### **OLDT Forest**

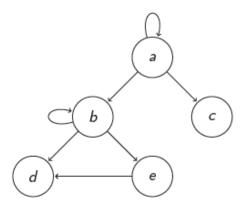
- When we get new answer, we will have to return to previous queries to continue their execution.
- When a literal is selected, mark it as a *consumer*.
- Check the call table.
  - If the literal is not in the table, **start a new tree for that literal** with its root marked as *generator*.
  - Resolve generator with program clauses (as in OLD), and add computed answers to its answer table.
- Resolve consumer with answers in its generator's table.

#### Calls and answers in tables

- Calls in table are standardized apart: i.e. their variables are renamed so that they are not identical to any other variable.
- Answers in a call's computed answer table share variables with their call. Any other variables are standardized apart.
- When checking if a literal is in call table:
  - We can check for *variance*: for a call **c** that is identical to the given literal **1**, modulo names of variables.
    - All answers to  $\mathbf{c}$  are answers to  $\mathbf{1}$ , and vice versa.
  - We can check for *subsumption*: for a call  ${\bf c}$  that is more general than a given literal  ${\bf l}$ , . i.e. if there is a substitution  ${\bf \theta}$  such that  ${\bf c}{\bf \theta}={\bf l}$ .
    - ullet Not all answers to ullet may be answers to ullet , but every answer to ullet is an answer to ullet

#### Notes on OLDT

- We can selectively mark which predicates we want to maintain tables for. (e.g. "p" in the previous example).
  - In general, no need to maintain tables for predicates defined solely by facts (i.e. clauses with empty bodies).
- For a Datalog program, there can be only finitely many distinct calls and answers.
  - So the size of tables is bounded.
- The number of literals in each goal is limited by the largest clause in the program (or original goal).
- Hence for Datalog, the OLDT forest as well as table sizes are bounded



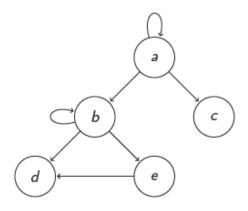
```
?- r(a, N)
```

```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

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- Construct a forest: one tree for each call
- Root of each tree (blue) is a generator
- Selected literal that matches a tabled call (green) is a consumer

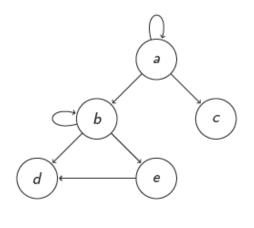
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



```
?- r(a, N)
?- q(a, N)
N=a N=b N=c
```

```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

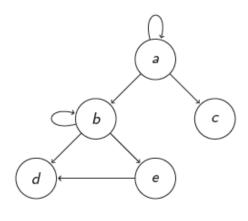
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
?- r(a, N)
?- q(a, N) ?- q(a, Z), r(Z, N)
N=a N=b N=c
?- r(a, N)
```

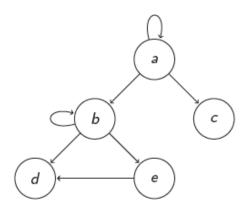
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

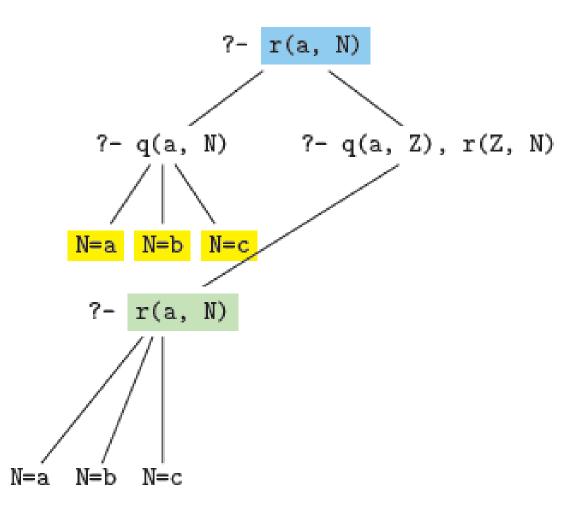
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```

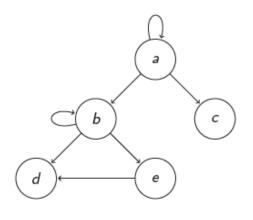
```
r(a, N)
      ?- q(a, N)
                     ?-q(a, Z), r(Z, N)
    N=a
         N=b
              N=c
         r(a, N)
N=a
```



```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).

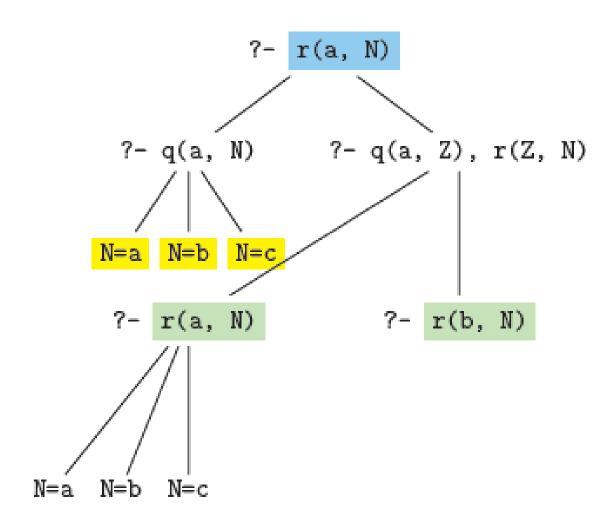


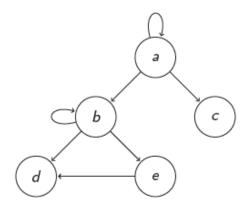


```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
```

q(e, d).

%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).





```
q(a, a).
```

q(a, b).

q(a, c).

q(b, b).

q(b, d).

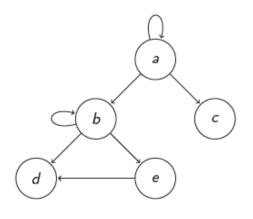
q(b, e).

q(e, d).

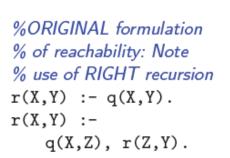
```
?- r(a, N)
?- q(a, N) ?- q(a, Z), r(Z, N)
N=a N=b N=c
?- r(a, N) ?- r(b, N)
```

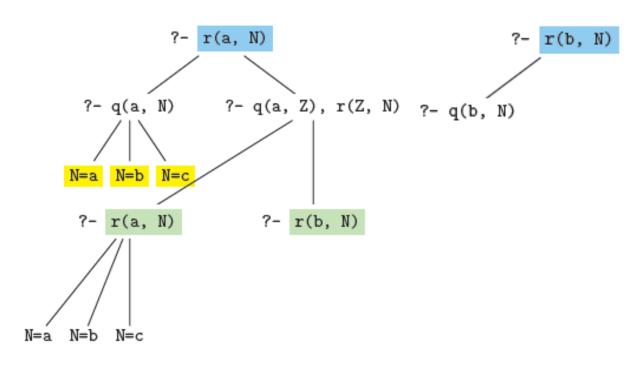
?- r(b, N)

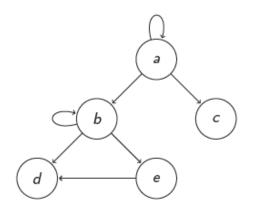
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :q(X,Z), r(Z,Y).



```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```







```
q(a, a).
q(a, b).
q(a, c).
```

q(b, b).

q(b, d).

q(b, e).

q(e, d).

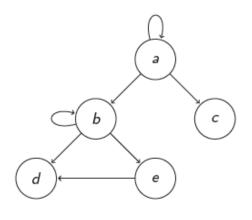
```
?- r(a, N)
?- q(a, Z), r(Z, N)

N=a N=b N=c
?- r(a, N)
?- r(b, N)
```

```
?- q(b, N)

N=b
```

```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
```

q(b, d). q(b, e).

q(e, d).

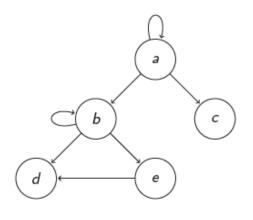
```
?- q(a, N) ?- q(a, Z), r(Z, N)

N=a N=b N=c ?- r(b, N)

N=a N=b N=c N=b
```

```
?- r(b, N)
?- q(b, N)
N=b
```

```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



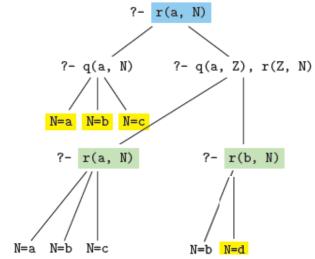
```
q(a, a).
q(a, b).
q(a, c).
```

q(b, b).

q(b, d).

q(b, e).

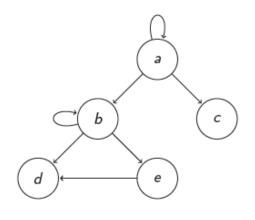
q(e, d).



```
?- q(b, N)
N=b N=d
```

r(b, N)

```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) := q(X,Y).
r(X,Y) :-
    q(X,Z), r(Z,Y).
```



```
q(a, a).
q(a, b).
```

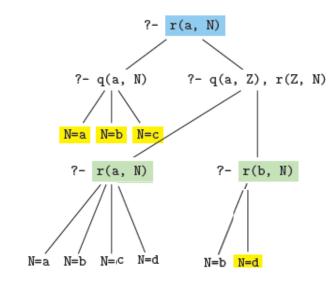
q(a, c).

q(b, b).

q(b, d).

q(b, e).

q(e, d).

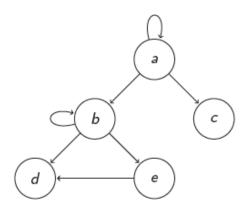


```
?- q(b, N)

?- q(b, N)

N=b N=d
```

```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



```
q(a, a).
```

q(a, b).

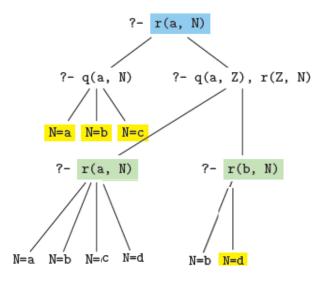
q(a, c).

q(b, b).

q(b, d).

q(b, e).

q(e, d).

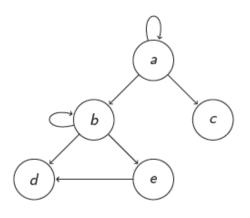


r(b, N)

?- q(b, N)

N=b N=d N=e

```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



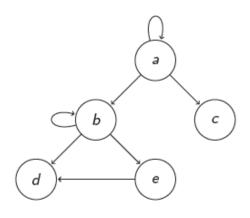
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
```

q(e, d).

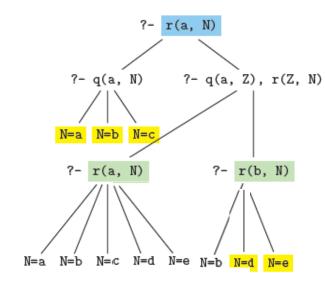
```
?- r(a, N)
?- q(a, N) ?- q(a, Z), r(Z, N)
N=a N=b N=c P- r(b, N)
N=a N=b N=c N=d N=b N=d N=e
```

```
?- r(b, N)
?- q(b, N)
N=b N=d N=e
```

```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

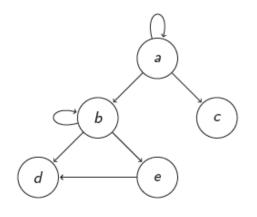


```
?- q(b, N)

?- q(b, N)

N=b N=d N=e
```

```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



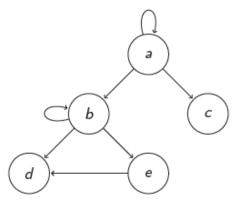
```
?- q(b, N)

?- q(b, N)

N=b N=d N=e
```

```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```

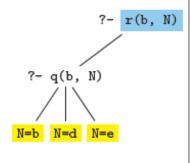


```
?- q(a, N) ?- q(a, Z), r(Z, N)

N=a N=b N=c

?- r(a, N) ?- r(b, N) ?- r(c, N)

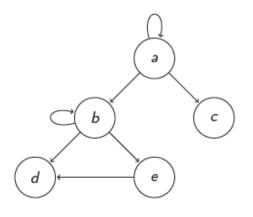
I=a N=b N=c N=d N=e N=b N=d N=e
```



```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

```
?- r(c, N)
```

```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



```
q(a, a).
q(a, b).
```

q(a, c).

q(b, b).

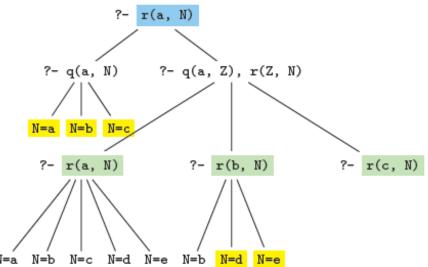
q(b, d).

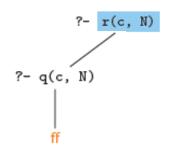
q(b, e).

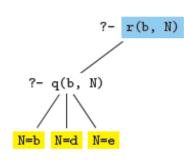
q(e, d).

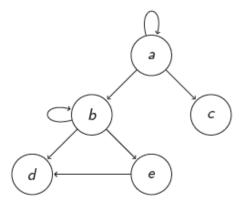
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-

q(X,Z), r(Z,Y).









```
q(a, a).
```

q(a, b).

q(a, c).

q(b, b).

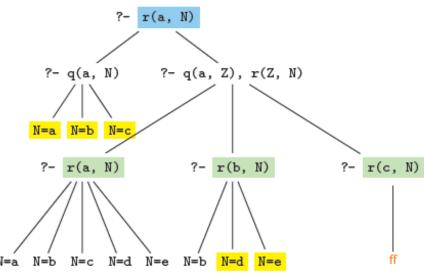
q(b, d).

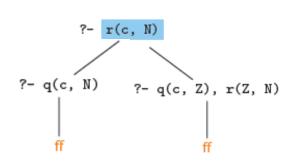
q(b, e).

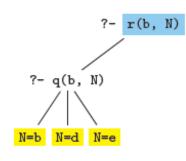
q(e, d).

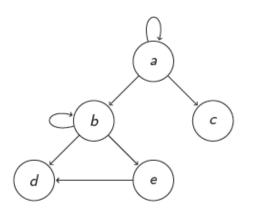
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-

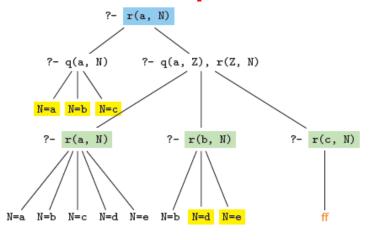
q(X,Z), r(Z,Y).

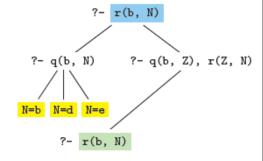




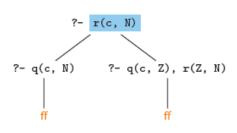




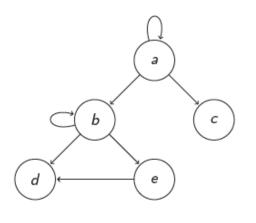




```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```



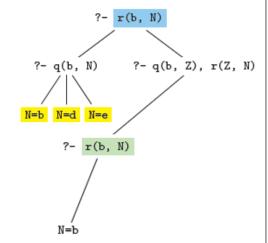
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



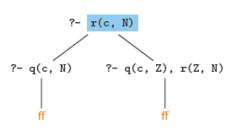
```
?- r(a, N)
?- q(a, N) ?- q(a, Z), r(Z, N)

N=a N=b N=c
?- r(a, N) ?- r(b, N) ?- r(c, N)

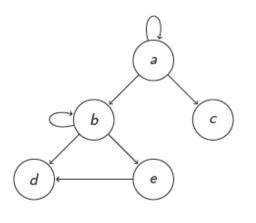
N=a N=b N=c N=d N=e N=b N=d N=e ff
```

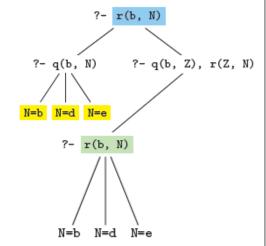


```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

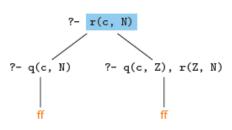


```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```

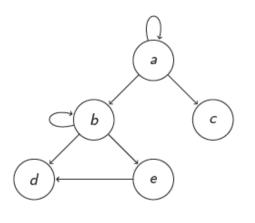


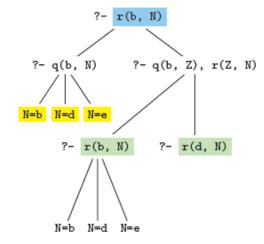


```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

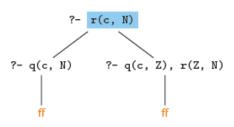


```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```

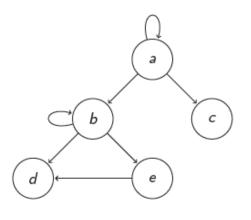


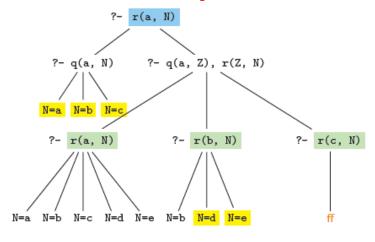


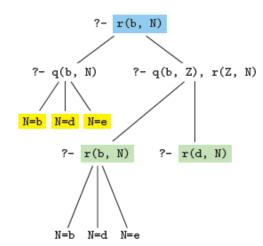
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```



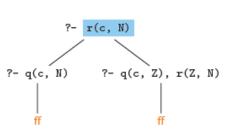
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```





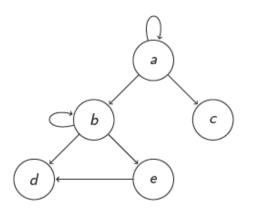


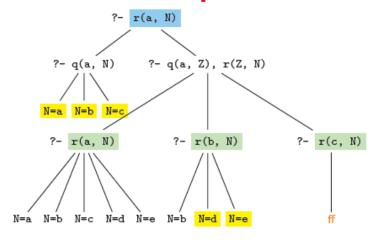
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

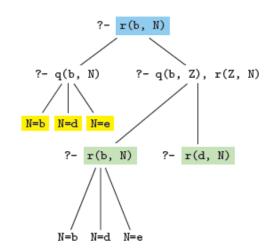


?- r(d, N)

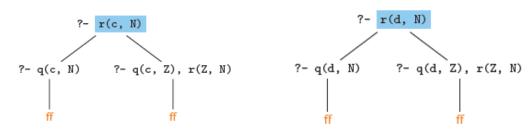
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



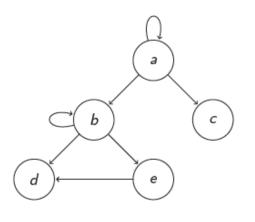


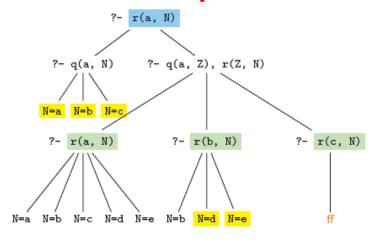


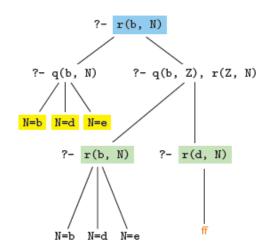
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```



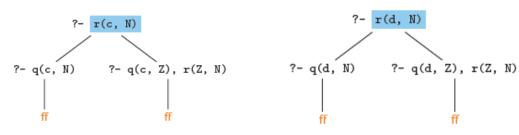
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



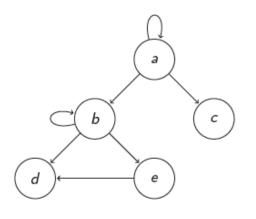


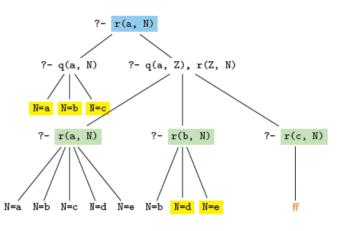


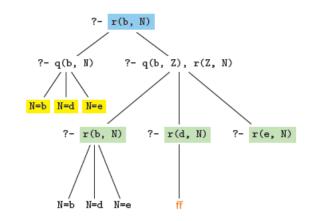
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```



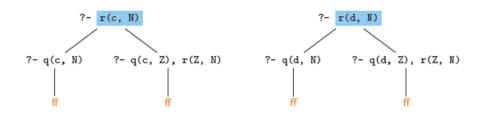
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



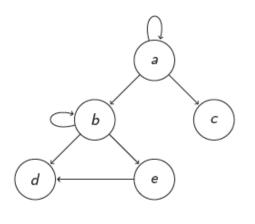


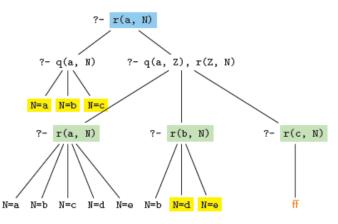


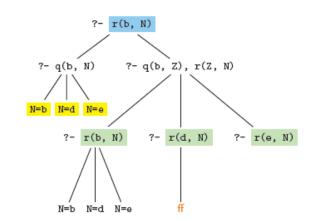
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```



```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```

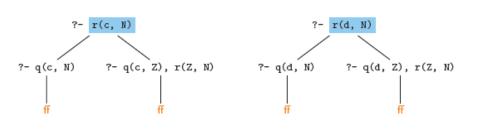




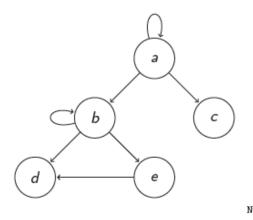


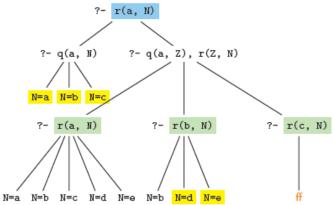
?- r(e, N)

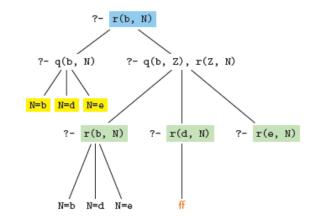
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```



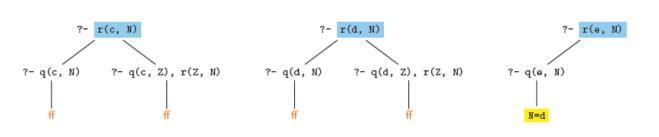
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



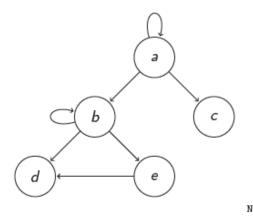


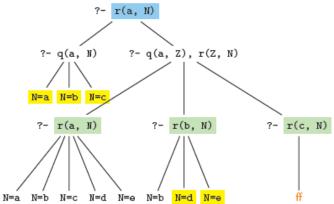


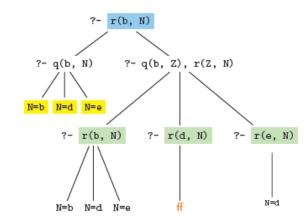
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```



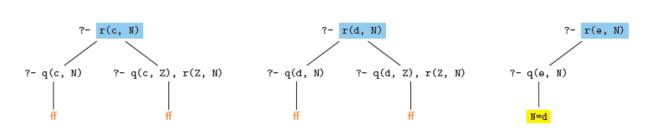
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



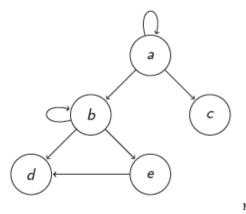


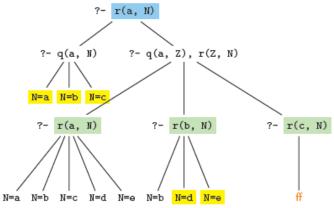


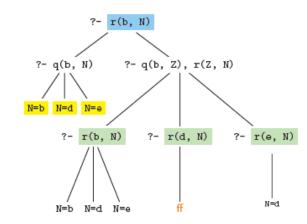
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```



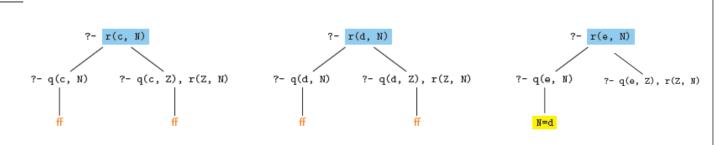
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```



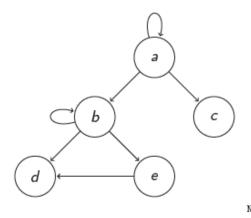


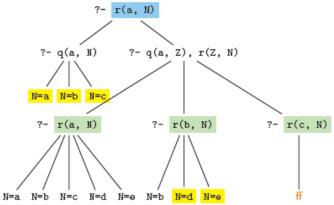


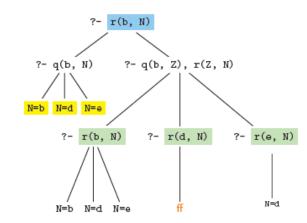
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```



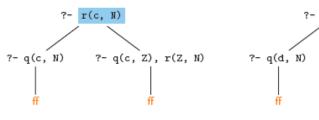
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```

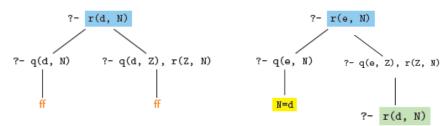




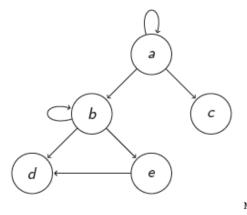


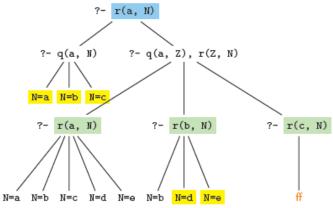
```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```

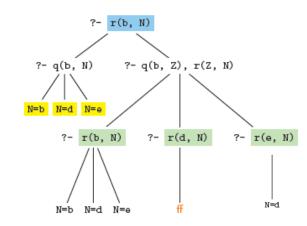




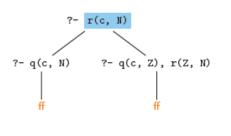
```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
q(X,Z), r(Z,Y).
```

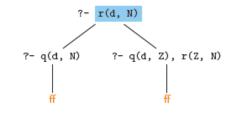


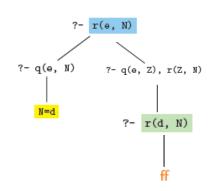




```
q(a, a).
q(a, b).
q(a, c).
q(b, b).
q(b, d).
q(b, e).
q(e, d).
```







```
%ORIGINAL formulation
% of reachability: Note
% use of RIGHT recursion
r(X,Y) :- q(X,Y).
r(X,Y) :-
```

q(X,Z), r(Z,Y).

## OLD Resolution with Tabling (OLDT)

- OLDT evaluation can be used to infer <u>negative</u> <u>answers</u>: e.g. a vertex is not reachable from another.
- Note that Breadth-First evaluation, or even the evaluation with goal-based stopping condition cannot do this.

# Tabled Resolution in XSB

```
reach(X,Y) := edge(X,Y).
reach(X,Y) := reach(X,Z), edge(Z,Y).
edge (a,a).
edge (a,b).
edge(b,c).
:- table(reach/2).
%OR :- auto table.
```

• Call:

?- reach(a,V).

```
 \begin{split} \text{reach} \left( X,Y \right) &:= \text{edge} \left( X,Y \right) \, . \\ \text{reach} \left( X,Y \right) &:= \text{reach} \left( X,Z \right) \, , \text{ edge} \left( Z,Y \right) \, . \\ \text{edge} \left( a,a \right) \, . \\ \text{edge} \left( a,b \right) \, . \\ \text{edge} \left( b,c \right) \, . \end{split}
```

• Calls
?- reach(a,V).
Answers

```
edge(a,Y)

reach(a,Y)

reach(a,Z), edge(Z,Y)
```

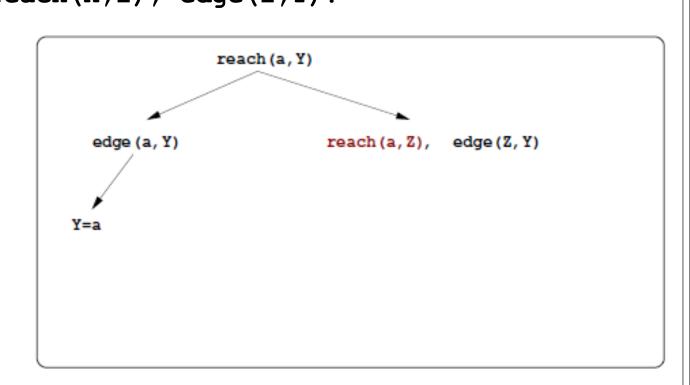
```
 reach(X,Y) := edge(X,Y). 
 reach(X,Y) := reach(X,Z), edge(Z,Y). 
 edge(a,a). 
 edge(a,b). 
 edge(b,c).
```

Calls

```
?- reach(a,V).
```

Answers

$$V = a$$



```
 \begin{split} \text{reach} \left( X,Y \right) &:= \text{edge} \left( X,Y \right) \, . \\ \text{reach} \left( X,Y \right) &:= \text{reach} \left( X,Z \right) \, , \text{ edge} \left( Z,Y \right) \, . \\ \text{edge} \left( a,a \right) \, . \\ \text{edge} \left( a,b \right) \, . \\ \text{edge} \left( b,c \right) \, . \end{split}
```

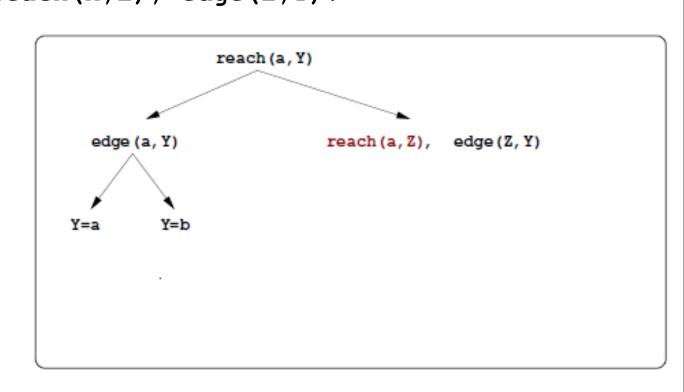
Calls

?- reach(a, V).

Answers

V = a

V = b



```
 reach(X,Y) := edge(X,Y). 
 reach(X,Y) := reach(X,Z), edge(Z,Y). 
 edge(a,a). 
 edge(a,b). 
 edge(b,c).
```

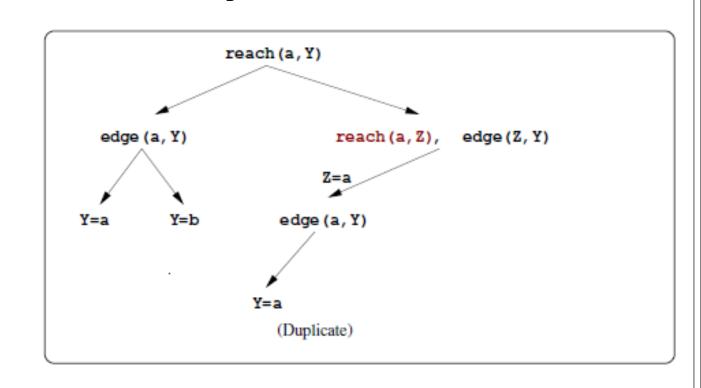
Calls

?- reach(a,V).

Answers

V = a

V = b



```
 reach(X,Y) := edge(X,Y). 
 reach(X,Y) := reach(X,Z), edge(Z,Y). 
 edge(a,a). 
 edge(a,b). 
 edge(b,c).
```

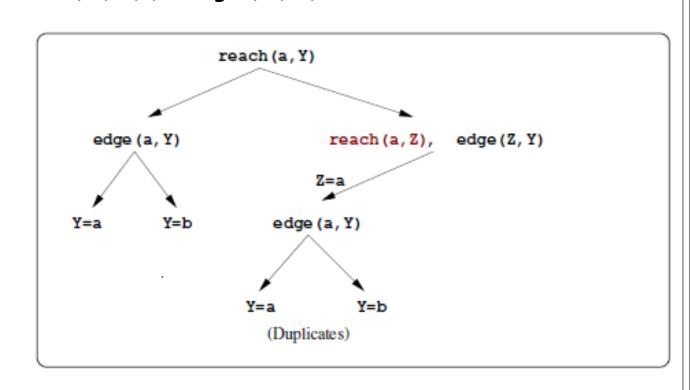
Calls

?- reach(a, V).

**Answers** 

V = a

V = b



```
 reach(X,Y) := edge(X,Y). 
 reach(X,Y) := reach(X,Z), edge(Z,Y). 
 edge(a,a). 
 edge(a,b). 
 edge(b,c).
```

Calls

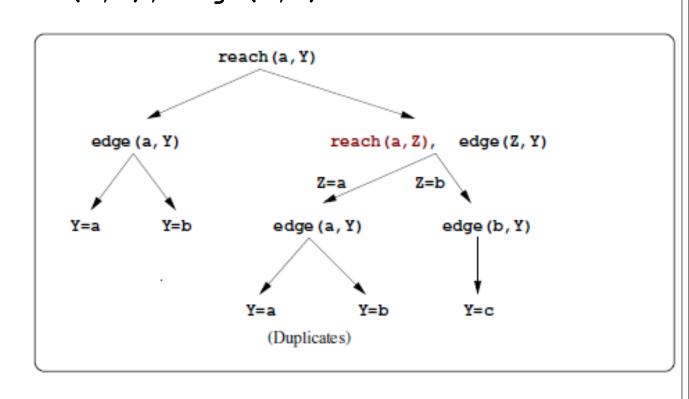
?- reach(a, V).

**Answers** 

V = a

V = b

V = c



```
 \begin{split} \text{reach} \left( X,Y \right) &:= \text{edge} \left( X,Y \right) \, . \\ \text{reach} \left( X,Y \right) &:= \text{reach} \left( X,Z \right) \, , \text{ edge} \left( Z,Y \right) \, . \\ \text{edge} \left( a,a \right) \, . \\ \text{edge} \left( a,b \right) \, . \\ \text{edge} \left( b,c \right) \, . \end{split}
```

Calls

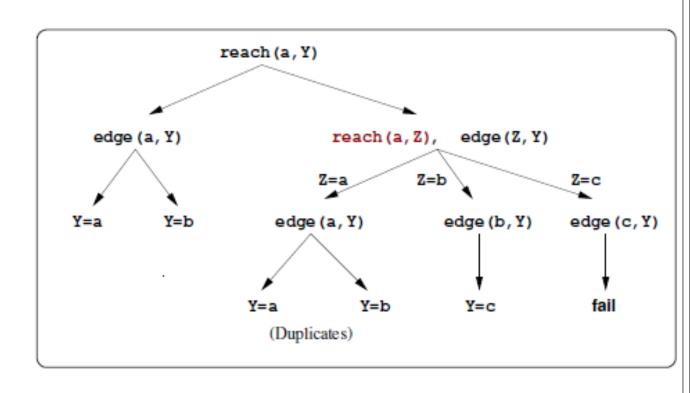
?- reach(a, V).

**Answers** 

V = a

V = b

V = c



**Answer completion!**