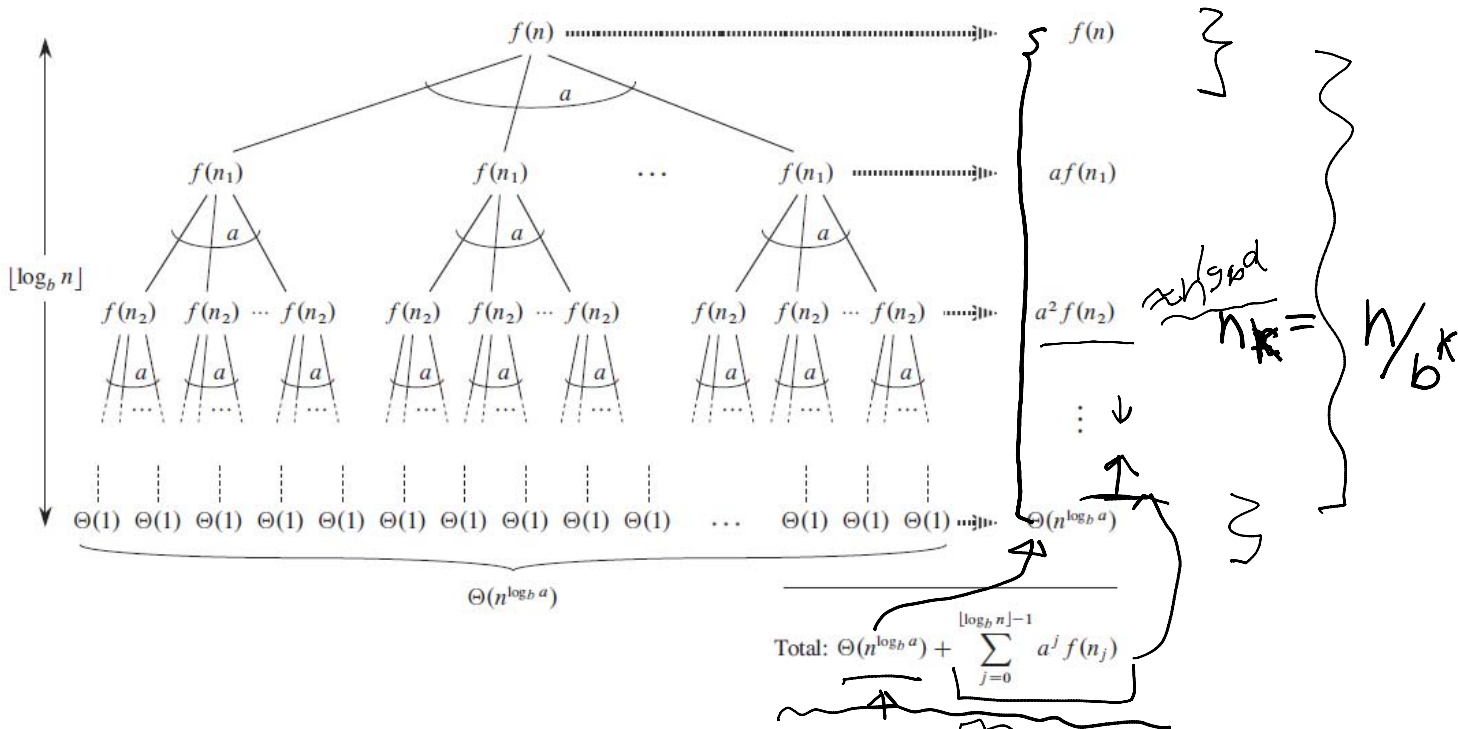


Master Theorem Cont.

2017년 9월 11일 월요일
오후 2:39

Recursion Tree for Master Theorem :
 $T(n) = aT(n/b) + f(n)$



Lemma 4.3

Let $a \geq 1$ and $b > 1$ be constants, and let $f(n)$ be a nonnegative function defined on exact powers of b . A function $g(n)$ defined over exact powers of b by

$$g(n) = \sum_{j=0}^{\log_b n - 1} a^j f(n/b^j) \tag{4.22}$$

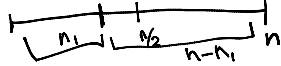
has the following asymptotic bounds for exact powers of b :

1. If $f(n) = O(n^{\log_b a - \epsilon})$ for some constant $\epsilon > 0$, then $g(n) = O(n^{\log_b a})$. 1st part done
2. If $f(n) = \Theta(n^{\log_b a})$, then $g(n) = \Theta(n^{\log_b a} \lg n)$.
3. If $af(n/b) \leq cf(n)$ for some constant $c < 1$ and for all sufficiently large n , then $g(n) = \Theta(f(n))$.

$$T(n) \leq C \quad n=1$$

$$T(n) \leq cn + T(n_1) + T(n-n_1)$$

$$\text{where } \lceil n/3 \rceil \leq n_1 \leq \lceil 2/3n \rceil \quad \text{if } n > 1$$



$$\text{Claim: } T(n) \leq an \lg n + bn \quad \begin{cases} a = \frac{3C}{\lg(\frac{27}{4})} \\ b = C \end{cases}$$

proof by Induction on n

Basis ($n=1$)

$$T(1) \leq a \cdot 1 \cdot \lg 1 + b \cdot 1 = b (=C)$$

Induction step ($n > 1$)

Assume by IH that claim holds for all $n' < n$

$$T(n) \leq cn + T(n_1) + T(n-n_1)$$

$$\leq cn + (a n_1 \lg n_1 + b n_1) + (a(n-n_1) \lg(n-n_1) + b(n-n_1))$$

$$= a [n_1 \lg n_1 + (n-n_1) \lg(n-n_1)] + \cancel{b n_1} + b n - \cancel{b n_1} + cn$$

$$\leq a \left[\frac{n}{3} \lg \frac{n}{3} + \left(\frac{2}{3}n \right) \lg \left(\frac{2}{3}n \right) \right] + (c+b)n$$

$$= \frac{an}{3} \lg n - \frac{an}{3} + \frac{2}{3}an \lg \frac{2}{3} - \frac{2}{3}an \lg \frac{3}{2} + (c+b)n$$

$$= \frac{an}{3} \lg n + n \left[-\frac{a}{3} \lg \frac{3 \times 9}{4} + b + c \right]$$

$$= an \lg n + bn$$

$$-\frac{a}{3} \lg \frac{27}{4} + b + c = b$$

$$a = \frac{3C}{\lg(\frac{27}{4})}$$

$$\therefore O(n \lg n)$$

assume

$$n_1 = \frac{1}{3}n$$

Theorem (Case 1)

let c and $0 < \alpha \leq \frac{1}{2}$ be const.

let $T(n) \leq c$ if $n=1$

$$T(n) \leq cn + T(n_1) + T(n-n_1)$$

$$\text{where } \lceil \alpha n \rceil \leq n_1 \leq \lceil (1-\alpha)n \rceil$$

then $T(n)$ is $O(n \lg n)$

Claim the previous proof to prove the theorem