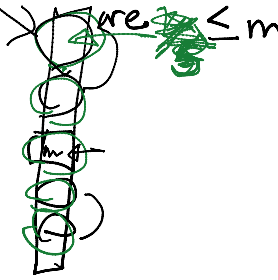


2. Choosing Partition element "near" the middle.

Rough analysis:



① at least  $n/4$  elements



② at least  $n/4$  elements  
 $n_{\geq m} \geq n/4$

Let  $T(n)$  be the time req. to select  $k$ th smallest element from  $S$

- Set  $M$  such that  $|M|$  at most  $n/5$  ( $r=5$ )  
 $\Rightarrow$  recursively call req. at most  $T(n/5)$  times
- Subset  $S_1^{<m}$  and  $S_3^{>m}$  are each of size at most  $3/4n$ .

$\Rightarrow S_1$ : at least  $\lfloor n/10 \rfloor$  elements of  $M_{\geq m}$

and there two distinct  $\overline{n \times 2}$  elements in  $S$  that are at least as large

$$\therefore |S_1| \leq n - 3 \lfloor n/10 \rfloor \leq 3/4n$$

Sub Analysis

$$1/4n \leq 3 \lfloor n/10 \rfloor$$

$$n/12 \leq \lfloor n/10 \rfloor$$

true for  $n \geq 50$

proof  $\searrow$

proof

$$\frac{n-9}{10} \leq \lfloor n/10 \rfloor$$

$$\frac{n}{12} \leq \frac{n-9}{10}$$

$$10n \leq 12n - 9 \times 12$$

$$\underline{54 \leq n.}$$

$$n = 53$$

$$n = 52$$

$$n = 51$$

$$\rightarrow \boxed{n=50} \leftarrow$$

$$h=49$$

$$\underline{n \neq 50}$$

$$\boxed{n > 50}$$

□

Recurrence relation of SELECT with PARTITION

$$T(n) < c \quad \text{when } n < 50$$

$$T(n) < T(n/5) + T(3n/4) + cn \quad \text{when } n \geq 50$$

Proof by induction on n that  $T(n) < an$  :