

$$T(n) \leq c \quad \text{when } n < 50$$

$$T(n) \leq T(n/5) + T(3n/4) + cn \quad \text{when } n \geq 50$$

Proof by induction on n that $T(n) \leq an$:

* Randomized Algorithm (RA)

⇒ Algorithm is randomized if its behavior is determined by

- ① input
- ② values produced by random-number generator in the Algo.

* Analysis of RA takes expected run-time over the distribution of values returned by the random number generator.

* probabilistic Analysis

* Choosing PE, m at random.

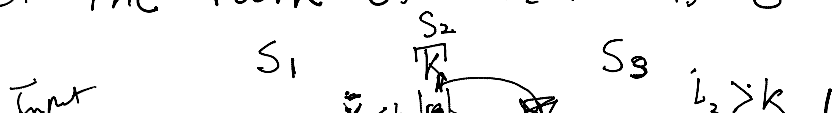
(assume equal prob of choosing $m = \frac{1}{n}$)

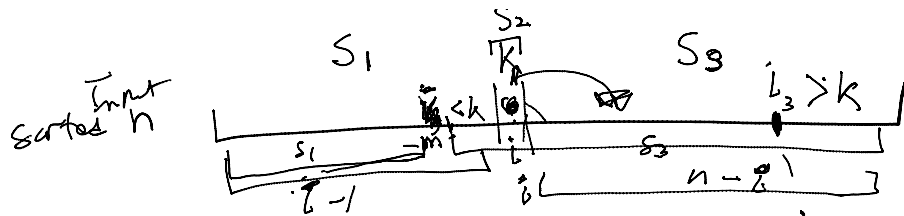
Base: if $|S| = 1$ $\therefore \text{prob}(m) = \frac{1}{n}$

then return that element (assume elements are distinct)

Choose index \bar{i} from any of $\{1, 2, \dots, n\}$

Let the rank of PE, m is \bar{i}





S_2 ① if $i_2 = k$ ($|S_1| < k$ & $|S_3| < k$) ← return m takes const time

S_3 ② if $i_2 < k$ ($|S_1| < k$) ← $|S_3| = n - i$
 \Rightarrow

S_1 ③ if $i_2 > k$ need to recurse on S_1
 $(S_1, i = i_2 - 1)$

∴ for given k , expected time of recurrence step:

$$\frac{1}{n} \left[\sum_{i=1}^{k-1} T(n-i) + \sum_{i=k+1}^n T(i-1) \right]$$

$i < k$ $i > k$

$n-1, n-2, \dots, n-k+1$ $k, k+1, \dots, n-1$

$$\approx \frac{1}{n} \left[\sum_{i=n-k+1}^{n-1} T(i) + \sum_{i=k}^{n-1} T(i) \right]$$

$n-k+1, \dots, n-1$ $k, k+1, \dots, n-1$

$$T(n) \leq \max_{1 \leq k \leq n} \left\{ \frac{1}{n} \left[\sum_{i=n-k+1}^{n-1} T(i) + \sum_{i=k}^{n-1} T(i) \right] \right\}$$

* Recurrence relationship of Randomized SELECT

$$T(n) \leq C \quad n=1$$

$$T(n) \leq Cn + \max_{1 \leq k \leq n} \left\{ \frac{1}{n} \left[\sum_{i=n-k+1}^{n-1} T(i) + \sum_{i=k}^{n-1} T(i) \right] \right\}$$

claim : $T(n) \leq an$ where $a = 4C$

proof : by induction n .

proof : by induction n.

Basiz (n=1)

$$T(1) \leq \boxed{C \leq a \cdot 1}$$

check letter
✓

Induction step. (n > 1)

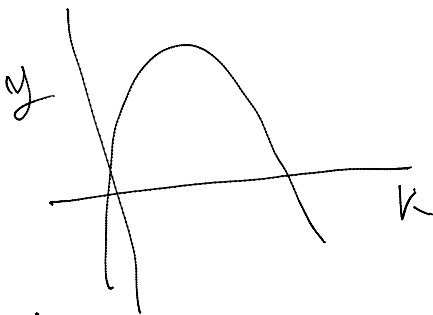
Assume as IH that claim holds for all n' < n

$$T(n) \leq Cn + \max_{1 \leq k \leq n} \left\{ \frac{a}{n} \left[\sum_{i=n-k+1}^{n-1} i + \frac{n-1}{2} \right] \right\}$$

$\frac{n-1}{2}$ $\overset{k \sim n+1}{\sim}$
 $\frac{n}{2} i$ $\overset{k+1 \sim n}{\sim}$

$$\sum_{i=1}^n i = \frac{n(n+1)}{2}$$

$$\begin{aligned} &\leq Cn + \max_{1 \leq k \leq n} \left\{ \frac{a}{n} \left[\frac{(n+1)(n) - (n-k)(n-k+1)}{2} \right] \right\} \\ &= Cn + \max_{1 \leq k \leq n} \left\{ \frac{a}{2n} \left(\frac{n(n+1)}{2} - \frac{k(k+1)}{2} \right) \right\} \\ &= Cn + \max_{1 \leq k \leq n} \left\{ \frac{a}{2n} (n^2 - n - n^2 + 2nk - n - k^2 + k + n^2 + n - k^2 - k) \right\} \\ &= Cn + \max_{1 \leq k \leq n} \left\{ \frac{a}{2n} [-2k^2 + 2nk + n^2 - n] \right\} \end{aligned}$$



$$\frac{dy}{dk} = -4k + 2n = 0 \quad \Rightarrow \quad \boxed{Cn + \frac{a}{4} [3n^2 - 2n]} \quad \Rightarrow \quad (C + \frac{3}{4}a)n = \frac{a}{2} \leq (C + \frac{3}{4}a)n = an$$

$$C + \frac{3}{4}a = a$$

$$C = \frac{1}{4}a$$

$$\boxed{a = 4C}$$

□

-5